

First-Order Structures

Satisfaction

Introduction

- A First-Order Structure is analogous to a truth assignment in Propositional Logic.
- For quantified sentences, we need a *domain of discourse*.
- It represents circumstances that determine the truth values of all the sentences of a language.
- It does it in a way that identity and FO Quantifiers are respected.

A Sample World

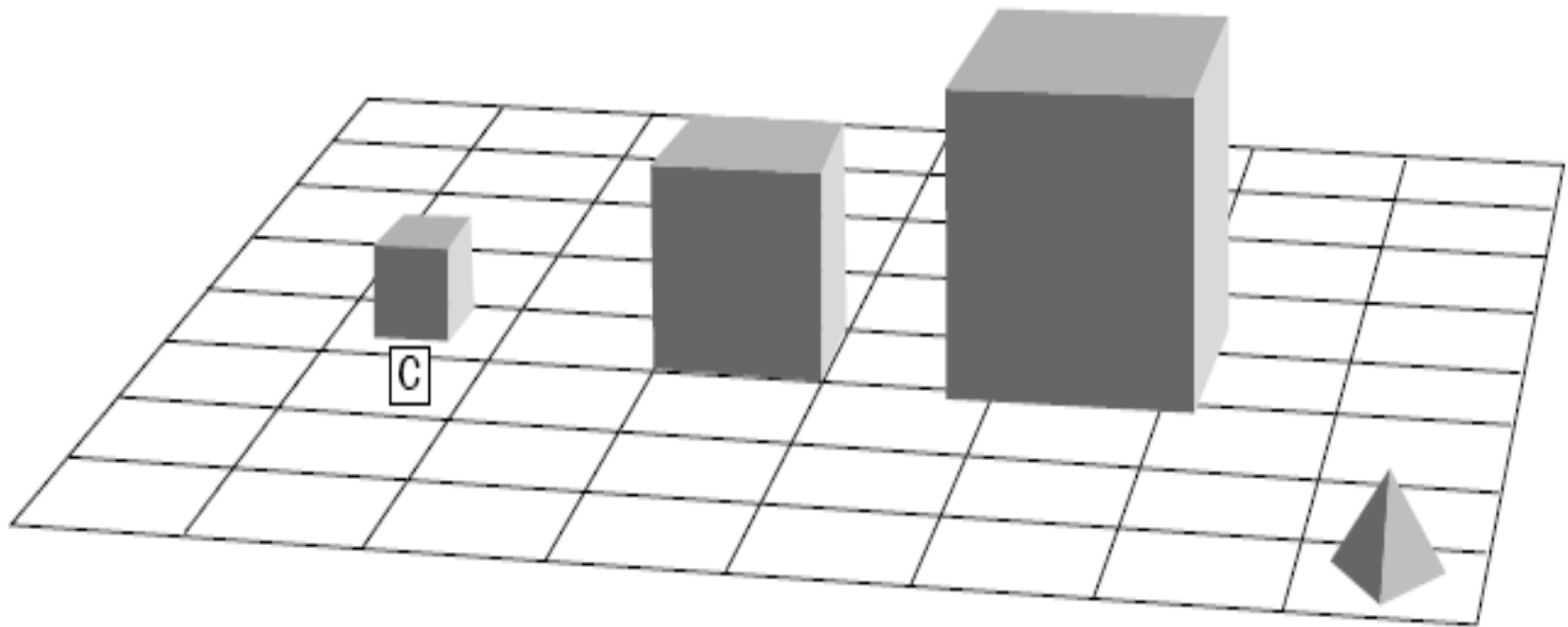


Figure 18.1: Mary Ellen's World.

Domain of Discourse

- *Domain of discourse* in this world is: $D = \{b_1, b_2, b_3, b_4\}$
- Let b_1 represent the left-most block, and go in order.
- Suppose that we have one name, **c**
- b_1 is the *referent* of the name **c**.

A Sample World

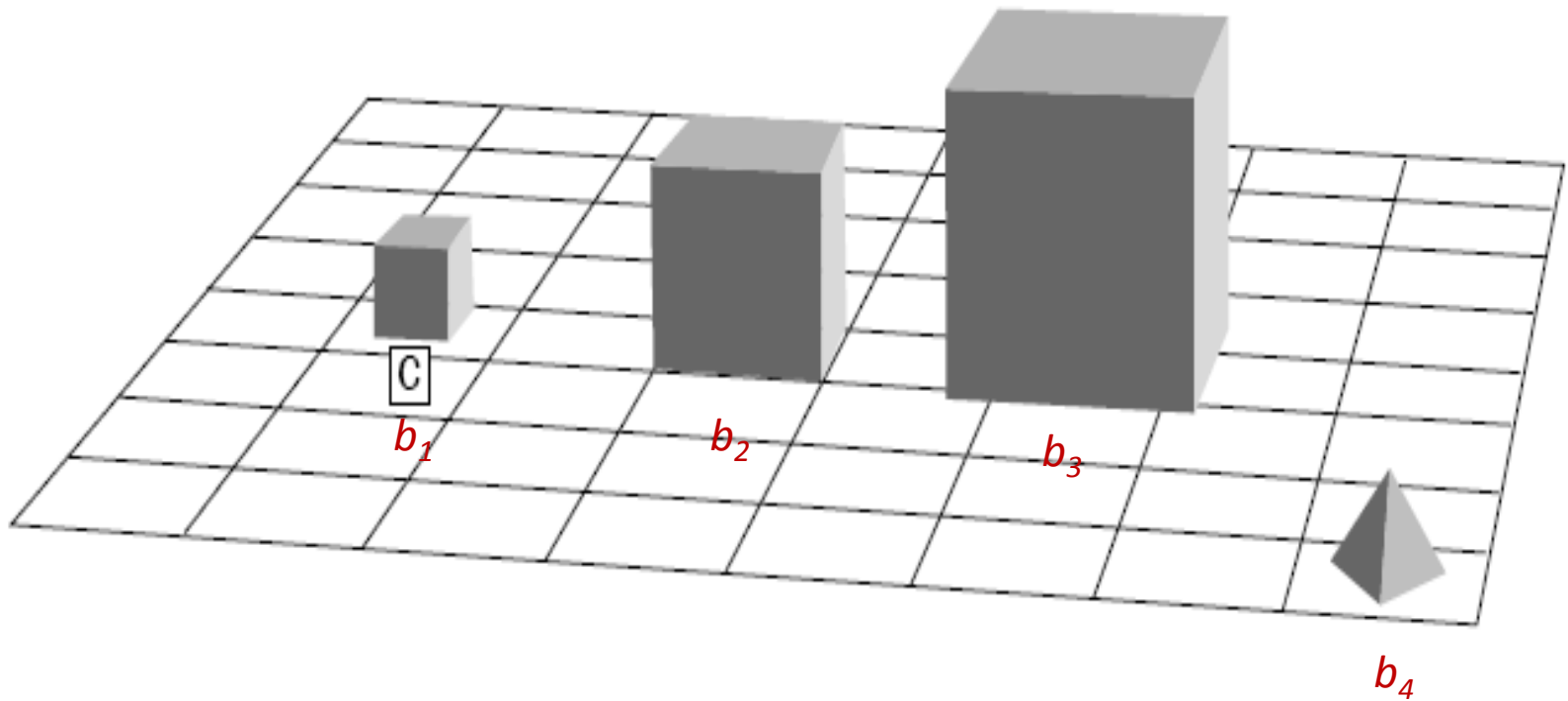


Figure 18.1: Mary Ellen's World.

Extensions

- Suppose in our language we only have the predicates: **Cube**, **Larger**, and =
- Assign the predicate Cube a subset of D :
 $Cu = \{b_1, b_2, b_3\}$
- Cu is said to be an *extension* of the predicate **Cube** in our structure.
- Assign the following pairs to the predicate **Larger**:
- $La = \{ \langle b_2, b_1 \rangle, \langle b_3, b_1 \rangle, \langle b_3, b_2 \rangle, \langle b_2, b_4 \rangle, \langle b_3, b_4 \rangle \}$
- La is said to be an *extension* of Larger.

Extension for Identity

- Assign the following pairs to the predicate =:
 $= \{ \langle b_1, b_1 \rangle, \langle b_2, b_2 \rangle, \langle b_3, b_3 \rangle, \langle b_4, b_4 \rangle \}$

Models

- In order to make truth assignments, we need:
 - A domain of discourse
 - Extensions of predicates
 - Referents of constants
- We will package them into one mathematical structure, called a “model.”
- A model is a function that is defined on the:
 - predicates of the language,
 - the names of the language,
 - and the quantifier symbol \forall .

Models

- The function \mathfrak{M} is called a *first-order structure*, provided the following conditions are satisfied:
 1. $\mathfrak{M}(\forall)$ is a nonempty set D , called the *domain of discourse*.
 2. If \mathbf{P} is an n -ary predicate symbol of the language, then $\mathfrak{M}(\mathbf{P})$ is a set of n -tuples $\langle x_1, \dots, x_n \rangle$ of elements of D . This set is called the extension of \mathbf{P} in \mathfrak{M} . It is required that the extension of the identity symbol consists of all pairs $\langle x, x \rangle$, for $x \in D$.
 3. If \mathbf{c} is any name of the language, then $\mathfrak{M}(\mathbf{c})$ is an element of D , and is called the *referent* of \mathbf{c} in \mathfrak{M} .

Modified Variable Assignments

- Consider wffs that start with a quantifier
- Let assignment g be defined on variable x .
- We want to say what it means for g to satisfy $\forall z \text{ Likes}(x, z)$.
- We need to be able to take any object b in the domain of discourse and consider the variable assignment which is just like g except that it assigns the value b to the variable z .
- We will say that g satisfies our wff $\forall z \text{ Likes}(x, z)$ if and only if every such modified assignment g' satisfies $\text{Likes}(x, z)$.
- To make this a bit easier to say, we introduce the notation " $g[z/b]$ " for the modified variable assignment.
- In general, $g[v/b]$ is the variable assignment whose:
 - domain is that of g plus the variable v
 - which assigns the same values as g ,
 - except that the new assignment assigns b to the variable v .

Denotation

- If a variable assignment g is appropriate for a wff \mathbf{P} , then between \mathfrak{M} and g , all the terms (constants and variables) in \mathbf{P} have a denotation.
- For any term \mathbf{t} , we write $[[\mathbf{t}]]_g^{\mathfrak{M}}$ for the denotation of \mathbf{t} .
- Thus $[[\mathbf{t}]]_g^{\mathfrak{M}}$ is:
 - $\mathbf{t}^{\mathfrak{M}}$ if \mathbf{t} is an individual constant
 - $g(\mathbf{t})$ if \mathbf{t} is a variable.