

CSSE 490 Model-Based Software Engineering: Automatic Programming Perspectives



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Learning Outcomes: MBE Discipline

Relate Model-Based Engineering as an engineering discipline.

- Discussion of Milestone 3
- Introduce Automatic Programming
- Look at Assistant approach (if time)



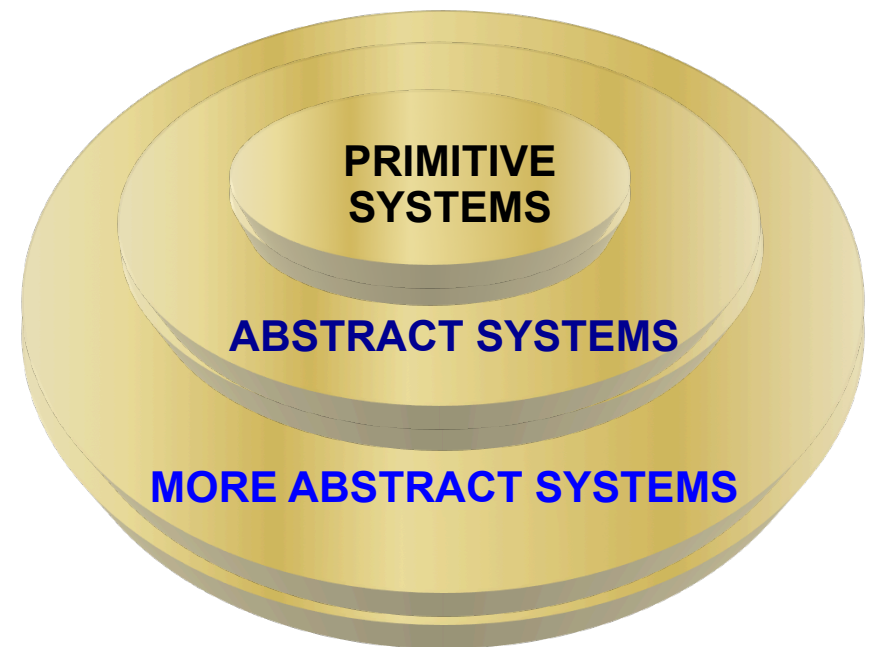
What would you say to the statement “Today’s specification language becomes tomorrow’s programming language?”

- Think for 15 seconds...
- Let’s talk...



Philosophy: Reliable Systems are Defined in Terms of Reliable Systems

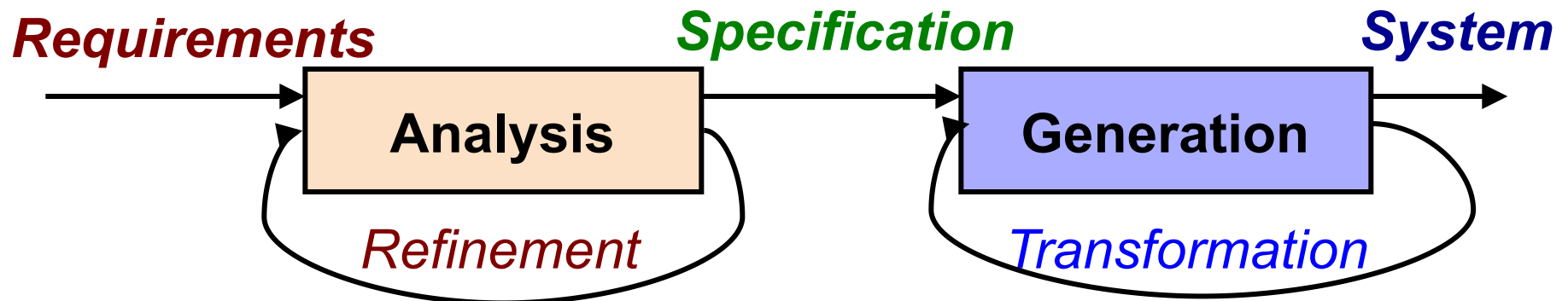
- Use only reliable systems
- Integrate these systems with reliable systems
- The result is a system(s) which is reliable
- Use resulting reliable system(s) along with more primitive ones to build new and larger reliable systems



A recursively reliable and reusable process

Automatic Programming

- Getting software to write software
- Great idea, but turns out to be hard
- Should be easier than other tasks
 - But programming requires some strategy (i.e., cunning and guile 😊)
 - Many human tasks difficult to automate



Automatic Programming

- Oversold early on and under-delivered on promises
- So people began to avoid this area
- “Automated Programming” became words of warning
- Since then, the limitations have eased
 - Memory space
 - Knowledge representation
 - Transformation systems



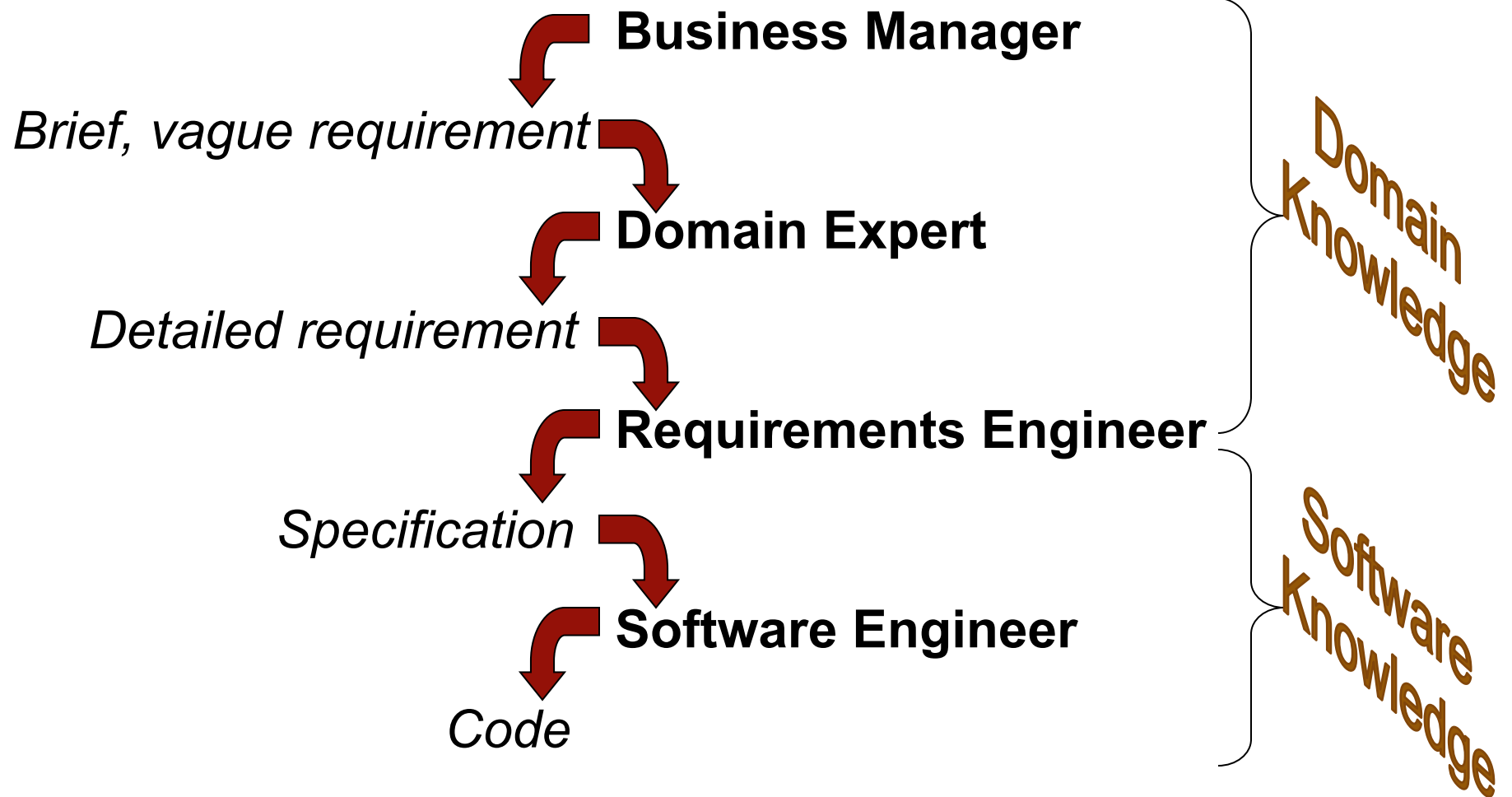
Complexity: Volume of Detail (Strauss)

31

1.2. gr. Fl.
3. gr. Fl.
1.2. Ob.
Engl. h.
1.2. Kl. (A)
1.2. Fag.
K: Fag.
1.2. H. (F)
3.4. H. (E)
Tb.
Pk.
Hfe.
Vi-Solo.
1. Vi.
2. Vi.
Br.
Vlc.
Kb.

The image displays a page of a musical score, page 31, featuring complex notation for various instruments. The score is written in a key signature of three sharps (F#, C#, G#) and a 3/4 time signature. The instruments listed on the left include: 1.2. gr. Fl., 3. gr. Fl., 1.2. Ob., Engl. h., 1.2. Kl. (A), 1.2. Fag., K: Fag., 1.2. H. (F), 3.4. H. (E), Tb., Pk., Hfe., Vi-Solo., 1. Vi., 2. Vi., Br., Vlc., and Kb. The notation is dense, with many notes, rests, and dynamic markings such as *cresc.* and *p*. The page number 31 is located in the top right corner.

Automatable Programming Activities



Transformational Approaches

Clear \longrightarrow *Efficient*

Specification

$X^{**}2 \longrightarrow X * X$

$m \leftarrow \min(A)$

$m \leftarrow \infty$
for $i \leftarrow 1$ **to** $\text{size}(A)$ **do**
 if $A[i] < m$ **then**
 $m \leftarrow A[i]$

Implementation



Natural Language Specification

“The SystemX transmission times are entered into the schedule”



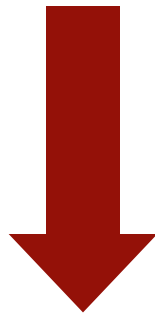
“Each SystemX clock transmission times and transmission length is made a component of a new transmission entry which is entered into the transmission schedule”

Problem is informality

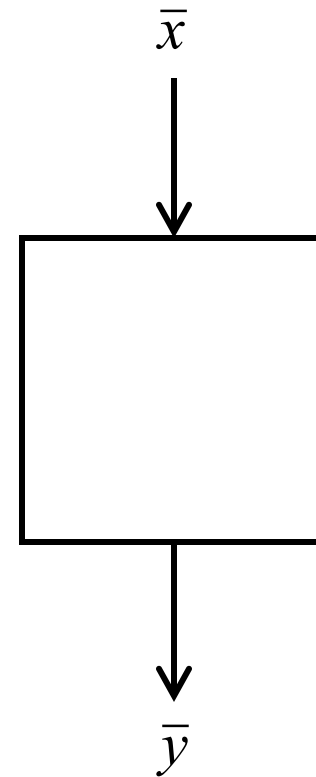
Deductive Synthesis

Precondition: $P(\bar{x})$

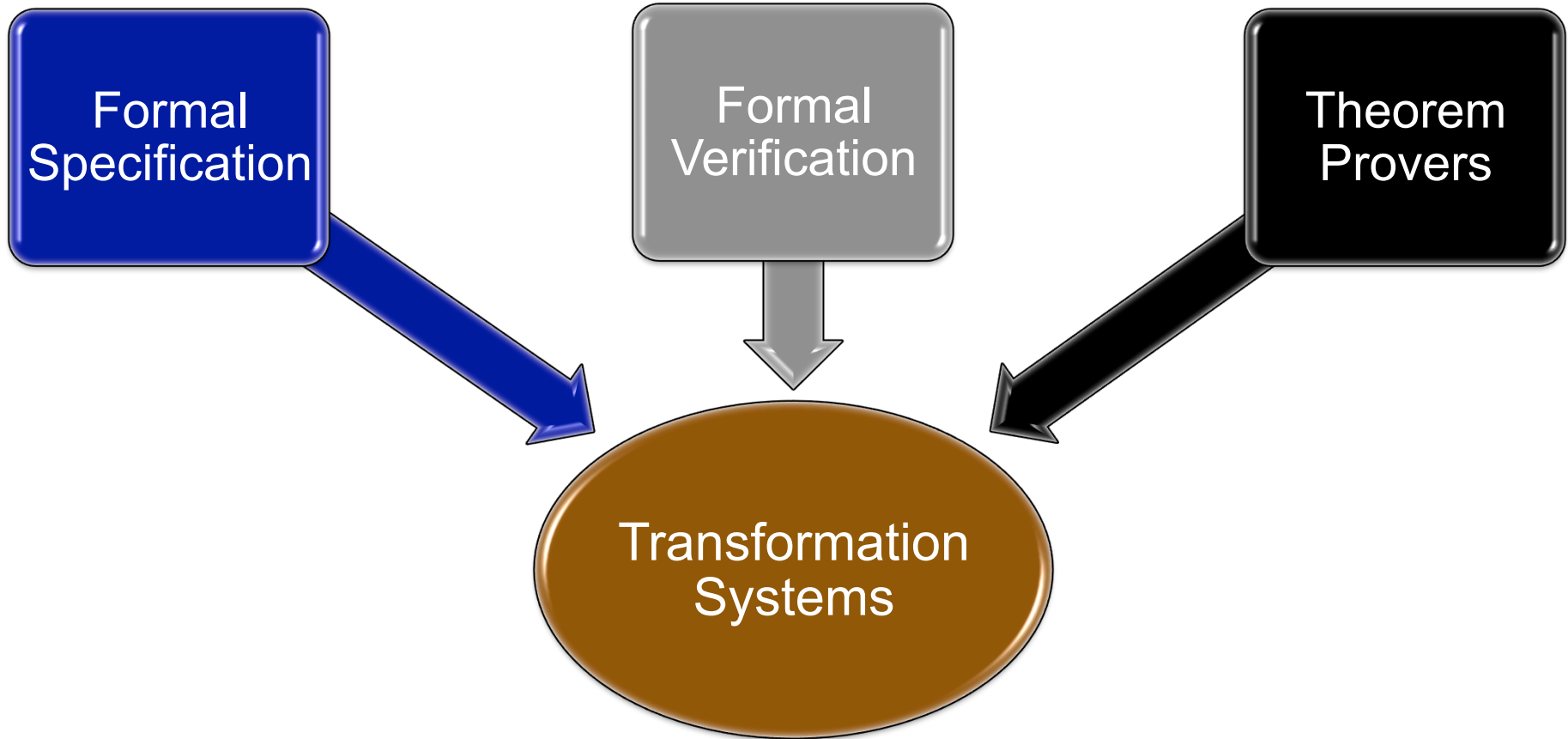
Postcondition: $Q(\bar{x}, \bar{y})$



$\forall \bar{x} \exists \bar{y} P(\bar{x}) \rightarrow Q(\bar{x}, \bar{y})$



Recall: Levels of Formality





Very High Level Formal Languages

prev := {}; *val* := {};

val(*x*) := {*x*};

(while *newnodes* ≠ {})

n from *newnodes*;

($\forall m \in \text{graph } \{n\}$)

newval := *val*(*n*) + *cost*(*n*, *m*);

if *val*(*m*) = *om* or *val*(*m*) > *newval* then

val(*m*) := *newval*;

prev(*m*) := *n*;

if *m* ≠ *y* then *newnodes* with := *m*; end if;

end if;

end \forall ;

end while;

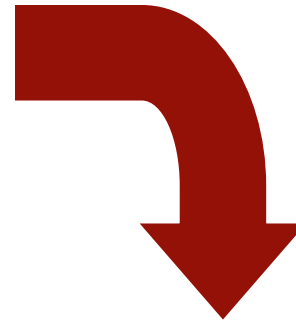
Programming by Example

{() → ()

(A B) → (A)

(A B C D) → (A B)

(A B C D E F) → (A B C)}



```
(DEFUN HALF (X)
  (H X X))
```

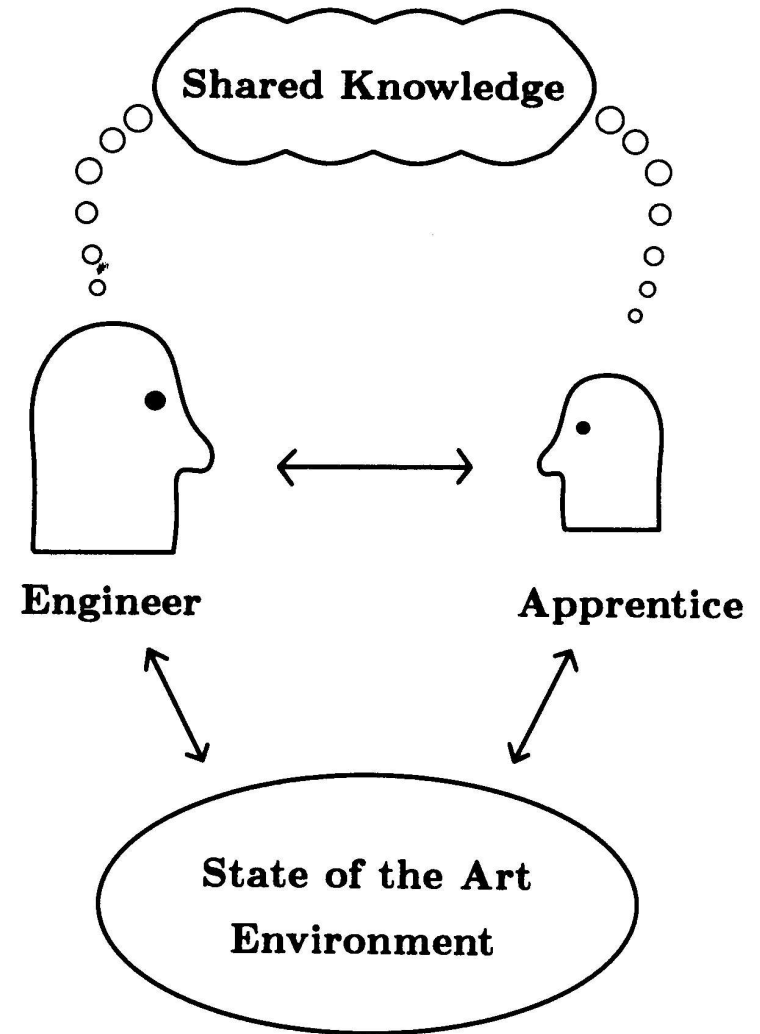
```
(DEFUN H (X Y)
  (COND ((ATOM Y) NIL)
        (T (CONS (CAR X)
                   (H (CDR X) (CDDR X)))))))
```

Computing Profession Choice... 😊



The Assistant Approach

- **Productivity: Delegate routine details**
- **Reliability: Standardization of common practices**





Example: Table-Lookup

function *table-lookup*(*table*, *key*)

bucket ← *table*[*hash*(*key*)]

loop

if *bucket* = *nil* then return *nil*

entry ← *head*(*bucket*)

if *key*(*entry*) = *key* then return *entry*

bucket ← *tail*(*bucket*)



Table-Insert

```
procedure table-insert(table, entry)  
push(entry, table[hash(key(entry))])
```



Table-Delete

```
procedure table-delete(table, key)  
index ← hash(key)  
bucket ← table[index]  
if key(head(bucket)) = key  
    then table[index] ← tail(bucket)  
    else bucket-delete(bucket, key)  
return table
```



Bucket-Delete

```
procedure bucket-delete(bucket, key)  
previous ← bucket  
loop  
    bucket ← tail(previous)  
    if bucket = nil then return nil  
    if key(head(bucket)) = key then  
        tail(previous) ← tail(tail(previous))  
        return nil  
    previous ← bucket
```



Analysis: Table-Lookup

function *table-lookup*(*table*, *key*)

bucket ← *table*[*hash*(*key*)]

loop

if *bucket* = *nil* then return *nil*

Linear
Search

entry ← *head*(*bucket*)

if *key*(*entry*) = *key* then return *entry*

bucket ← *tail*(*bucket*)



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List
Enumeration

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bucket ← *tail(previous)*

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if *bucket* = *nil* then return *nil*

if *key(head(bucket))* = *key* then

tail(previous) ← *tail(tail(previous))*

return *nil*

previous ← *bucket*



Analysis: Bucket-Delete

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previous ← *bucket*

loop

bucket ← *tail(previous)*

if *bucket* = *nil* then return *nil*

if *key*(*head(bucket)*) = *key* then

tail(previous) ← *tail(tail(previous))*

return *nil*

previous ← *bucket*

Trailing-
Pointer
List
Enumeration



Analysis: Bucket-Delete

procedure *bucket-delete*(*bucket*, *key*)

previous ← *bucket*

loop

bucket ← *tail(previous)*

if *bucket* = *nil* then return *nil*

if *key(head(bucket))* = *key* then

tail(previous) ← *tail(tail(previous))*

return nil

previous ← *bucket*

Splice
Out

Analysis: Bucket-Delete

procedure *bucket-delete*(*bucket*, *key*)

previous ← *bucket*

loop

bucket ← *tail(previous)*

if *bucket* = *nil* then return *nil*

if *key(head(bucket))* = *key* then

tail(previous) ← *tail(tail(previous))*

return *nil*

previous ← *bucket*

Linear
Search

Trailing-
Pointer
List
Enumeration

Splice
Out



Homework and Milestone Reminders

- Read Chapter 12 in text
- Milestone 3: Light-Weight Transformation Environment (see Milestone 3 assignment)
 - Due by 11:55pm, Thursday, May 5th, 2011.