

## ● Announcements:

- DES due now
- Chapter 3 Exam tomorrow
  - No cheat sheets allowed
- Term project groups and topics due end of next week
  - Use ch 10 – 19 as inspiration

## ● Today

- Using RSA: practical considerations

## ● Questions?

# RSA (Rivest – Shamir – Adelman)

For Alice to send a message to Bob.

- Bob chooses primes  $p, q$  (large,  $\sim 100$  digits each)
- He publishes his public key  $(n, e)$ :
  - $n = pq$
  - $e$ , a large number such that  $\gcd(e, (p-1)(q-1)) = 1$
- Alice has a message  $m < n$ .
  - Otherwise (if  $m > n$ ), break message into chunks  $< n$
- Alice sends  $c = m^e \pmod n$
- Bob computes  $c^d \pmod n = (m^e)^d = m \pmod n$ .
- What does he use for  $d$ ?

# The security of RSA lies in the difficulty of factoring products of large primes

Are there any shortcuts to decryption?

Consider:

1. Can we find the  $e^{\text{th}}$  root of  $c = m^e$  quickly?  
**No, since mod  $n$**
2. Is  $\phi(n)$  as hard to find as factors  $p$  and  $q$ ?  
Yesterday: **yes**, using  $n - \phi(n) + 1$  and the quadratic formula
3. Is finding  $d$  directly is as hard to do as finding  $p$  and  $q$ ?  
Next week: **yes!**

# Toy example

- Alice – (m) → Bob
- Bob's key:
  - $n = pq = (13)(17) = 221$
  - $e = 35: \gcd(e, (p-1)(q-1)) = 1$
  - $d = e^{-1} \pmod{192}$  exists:  
 $d = \underline{\quad 11 \quad}$
- $m = 20$  (letter t)
  - 1-based, so leading 'a' = 1 not ignored
- $c = m^e \pmod{n} = \underline{\quad 197 \quad}$
- $c^d \pmod{n} = \underline{\quad 20 \quad}$

Issues:

How to compute  
 $20^{35} \pmod{221}$ ?

Efficiency is  $O(\log e)$

How to compute  $d$ ?

Extended Euclidean alg.

- And why is this secure?
  - Why can't Eve calculate  $d$  herself?

# Example with larger numbers

- Maple's worksheet mode
- For some reason, inert power (&^) only works for me when entering in the red (single-line exponents) entry-mode; press F5 (or |> button) to toggle.
- myConcatenator is a lambda expression.

```
> msg := convert("hello", bytes);
                                     msg := [104, 101, 108, 108, 111]
> myConcatenator := arr → sum(arr[i] · 1000(i-1), i=1..nops(arr));
                                     myConcatenator := arr → ∑i=1nops(arr) arr[i] · 1000(i-1)
> m := myConcatenator(msg);
                                     m := 111108108101104
> p := nextprime(1020);
                                     p := 1000000000000000000039
> q := nextprime(1021);
                                     q := 100000000000000000000117
>
> n := p · q;
                                     n := 1000000000000000000050700000000000000004563
> e := 65537;
                                     e := 65537
> φ := (p-1) · (q-1);
                                     φ := 1000000000000000000049600000000000000004408
> c := m &^ e mod n;
                                     c := 41172530747560554631603662398453570506594
>
> d := e&^(-1) mod phi;
                                     d := 45366739399118055472095262218288905505761
> dec := c&^d mod n;
                                     dec := 111108108101104
> convert(dec, base, 1000);
                                     [104, 101, 108, 108, 111]
> convert(%, bytes);
                                     "hello"
> {
```