

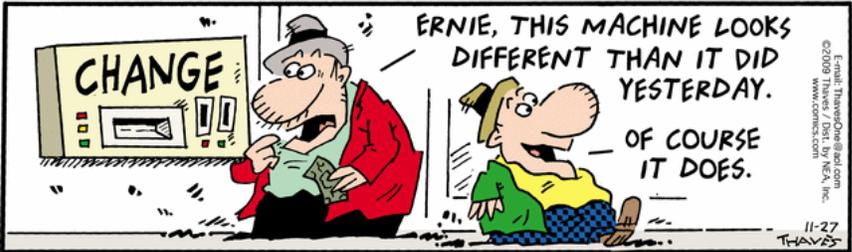
MA/CSSE 474

Theory of Computation

Languages, prefixes, sets,
cardinality, functions

Your Questions?

- Syllabus
- Tuesday's discussion
- Reading Assignments
- HW1 or HW2
- Anything else



Must not be a FSM

ERNIE, THIS MACHINE LOOKS DIFFERENT THAN IT DID YESTERDAY.

OF COURSE IT DOES.

11-27 THAVES

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More about Languages and Strings

Mostly very quick.

Some should be review of previous courses, and some others you should have gotten for Reading Quiz 2.

Ask questions if there are things I list here that you are not sure about.



Relations on Strings:

Substring, proper substring

Every string is a substring of itself.

ε is a substring of every string.

prefix, proper prefix

Every string is a prefix of itself.

ε is a prefix of every string.

s is a *suffix, proper suffix*, self, ε

Defining a Language

A **language** is a (finite or infinite) set of strings over a finite alphabet Σ . Examples for $\Sigma = \{a, b\}$

1. $L = \{x \in \{a, b\}^* : \text{all } a\text{'s precede all } b\text{'s}\}$
 $\epsilon, a, aa, aabbb, \text{ and } bb \text{ are in } L. \text{ } aba, ba, \text{ and } abc \text{ are not in } L.$
2. $L = \{x : \exists u \in \{a, b\}^* : x = ua\}$
 Simple English description:
3. $L = \{x\#y : x, y \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}^* \text{ and, when } x \text{ and } y \text{ are viewed as the decimal representations of natural numbers, } \text{square}(x) = y\}$.
 Examples (in L or not?):
 $3\#9, 12\#144, 3\#8, 12, 12\#12\#12, \#$
4. $L = \{a^n : n \geq 0\}$ simpler description of L?
5. $A^n B^n = \{a^k b^k : k \geq 0\}$
6. $L = \emptyset = \{ \}$
7. $L = \{\epsilon\}$

You saw in Reading Quiz 2 that the last two examples are different languages

Natural Languages are Tricky

$L = \{w : w \text{ is a sentence in English}\}$.

Examples:

Kerry hit the ball.

Colorless green ideas sleep furiously.

The window needs fixed.

Ball the Stacy hit blue.



A Halting Problem Language

$L = \{w: w \text{ is a Java program that, when given any finite input string, is guaranteed to halt}\}.$

- Is this language well specified?
- Can we decide which strings L contains?



Languages and Prefixes

What are the following languages?

$L = \{w \in \{a, b\}^*: \text{no prefix of } w \text{ contains } b\}$

$L = \{w \in \{a, b\}^*: \text{no prefix of } w \text{ starts with } a\}$

$L = \{w \in \{a, b\}^*: \text{every prefix of } w \text{ starts with } a\}$

Concatenation of Languages

If L_1 and L_2 are languages over Σ :

$$L_1 L_2 = \{w \in \Sigma^* : \exists s \in L_1 (\exists t \in L_2 (w = st))\}$$

Example:

$$L_1 = \{a, aa\}$$

$$L_2 = \{a, c, \varepsilon\}$$

$$L_1 L_2 =$$

Alternate definition:

$$L_1 L_2 = \{st : s \in L_1 \wedge t \in L_2\}$$

Simpler than the first definition,
but the first one conveys the idea
more precisely.

Operations on Languages

- L^R

- L^3 Is this the same as $\{w^3 : w \in L\}$

- L^0

- L^k

- L^*

- L^+

Concatenation and Reverse of Languages

Theorem: $(L_1 L_2)^R = L_2^R L_1^R$.

Proof:

$\forall x (\forall y ((xy)^R = y^R x^R))$ Theorem 2.1 we proved last time

$(L_1 L_2)^R = \{(xy)^R : x \in L_1 \text{ and } y \in L_2\}$ Definition of concatenation of languages

$= \{y^R x^R : x \in L_1 \text{ and } y \in L_2\}$ Thm 2.1

$= L_2^R L_1^R$ Definition of concatenation of languages

Sets and Relations

Sets of Sets

- The **power set** of S is the set of all subsets of S .

Let $S = \{1, 2, 3\}$. Then:

$$\mathcal{P}(S) = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}.$$

- $\Pi \subseteq \mathcal{P}(S)$ is a **partition** of a set S iff:
 - Every element of Π is nonempty,
 - Every pair of elements of Π is disjoint, and
 - the union of all the elements of Π equals S .

Some partitions of $S = \{1, 2, 3\}$:

$$\{\{1\}, \{2, 3\}\} \text{ or } \{\{1, 3\}, \{2\}\} \text{ or } \{\{1, 2, 3\}\}.$$

How many *different* partitions of S ?

Closure

- A set S is **closed** under binary operation op iff

$$\forall x, y \in S (x \text{ op } y \in S),$$

closed under unary

function f iff

$$\forall x \in S (f(x) \in S)$$

If S is not closed under unary function f , a **closure** of S is a set S' such that

- S is a subset of S'
- S' is closed under f
- No proper subset of S' contains S and is closed under f

- Examples**
- \mathbb{N}^+ (the set of all positive integers) is closed under addition and multiplication but not negation, subtraction, or division.
- What is the closure of \mathbb{N}^+ under subtraction? Under division?
- The set of all finite sets is closed under union and intersection. **Closed under infinite union?**

Equivalence Relations

A relation on a set A is any set of ordered pairs of elements of A .

A relation $R \subseteq A \times A$ is an **equivalence relation** iff it is:

- reflexive,
- symmetric, and
- transitive.

Examples of equivalence relations:

- Equality
- Lives-at-Same-Address-As
- Same-Length-As
- Contains the same number of a's as

Show that \equiv_3
is an
equivalence
relation

Cardinality of a set.

The cardinality of every set we will consider is one of the following :

- a specific natural number (if S is finite),
- “countably infinite” (if S has the same number of elements as there are integers), or
- “uncountably infinite” (if S has more elements than there are integers).

Functions on Languages

Functions whose domains and ranges are languages

$$\text{maxstring}(L) = \{w \in L : \forall z \in \Sigma^* (z \neq \varepsilon \rightarrow wz \notin L)\}.$$

Examples:

- $\text{maxstring}(A^n B^n)$
- $\text{maxstring}(\{a\}^*)$

Exercise for later:

What language is
 $\text{maxstring}(\{b^n a : n \geq 0\})$?

Let INF be the set of all infinite languages.

Let FIN be the set of all finite languages.

Are the language classes FIN and INF closed under maxstring ?

Functions on Languages

$\text{chop}(L) =$

$$\{w : \exists x \in L (x = x_1 c x_2, x_1 \in \Sigma_L^*, x_2 \in \Sigma_L^*, c \in \Sigma_L, |x_1| = |x_2|, \text{ and } w = x_1 x_2)\}.$$

What is $\text{chop}(A^n B^n)$?

What is $\text{chop}(A^n B^n C^n)$?

Are FIN and INF closed under chop ?

Functions on Languages

$$\text{firstchars}(L) = \{w : \exists y \in L (y = cx \wedge c \in \Sigma_L \wedge x \in \Sigma_L^* \wedge w \in \{c\}^*)\}.$$

What is $\text{firstchars}(A^n B^n)$?

What is $\text{firstchars}(\{a, b\}^*)$?

Are FIN and INF closed under firstchars ?

- Look at the Reading Quiz 1 comments from the Day 2 slides.