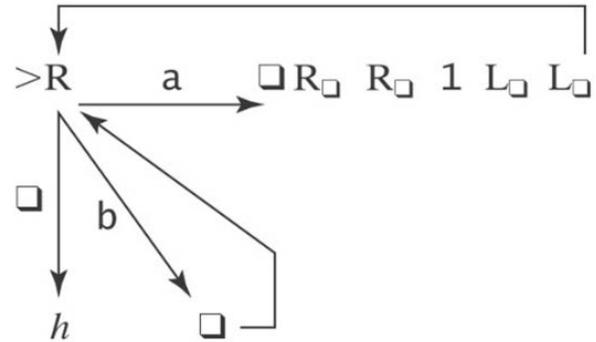


**MA/CSSE 474 Day 29 Summary**

- 1) A macro language for TMs. On a separate handout. You need to learn to read and write this language.  
Suggestion: Write notes about the macro language on your handout.
- 2) Exercise: What does this machine do?



- 3) Exercise: Initial input on the tape is an integer written in binary, most significant bit first (110 represents 6).  
Design a TM that replaces the binary representation of  $n$  by the binary representation of  $n+1$ .

- 4) **TMs as language recognizers.** Let  $M = (K, \Sigma, \Gamma, \delta, s, \{y, n\})$ .
  - a)  $M$  **accepts** a string  $w$  iff  $(s, \square w) \vdash_{-M}^* (y, w')$  for some string  $w'$ .
  - b)  $M$  **rejects** a string  $w$  iff  $(s, \square w) \vdash_{-M}^* (n, w')$  for some string  $w'$ .
  - c)  $M$  **decides** a language  $L \subseteq \Sigma^*$  iff for any string  $w \in \Sigma^*$  it is true that:
    - i) if  $w \in L$  then  $M$  accepts  $w$ , and
    - ii) if  $w \notin L$  then  $M$  rejects  $w$ .
  - d) A language  $L$  is **decidable** iff there is a Turing machine  $M$  that decides it.
  - e) We define the set **D** to be the set of all decidable languages.
  - f)  $M$  **semidecides**  $L$  iff, for any string  $w \in \Sigma_M^*$ :
    - i)  $w \in L \rightarrow M$  accepts  $w$
    - ii)  $w \notin L \rightarrow M$  does not accept  $w$ .  $M$  may either reject or fail to halt.
  - g) A language  $L$  is **semidecidable** iff there is a Turing machine that semidecides it.
  - h) We define the set **SD** to be the set of all semidecidable languages.