

474 HW 10 problems (highlighted problems are the ones to turn in)

11.1, 2, 3
(#1, 2, 3)

11.4

(#4) (3)

11.6b

(#5) (6)

11.6c

(#6) (6)

11.6d

(#7) (6)

11.6e (#8)

(#9) not

from

textbook

(6)

11.6h

(#10) (6)

11.6i

(#11) (6)

11.6k

(#12) (6)

11.8b

(#13) (9)

11.9 (#14)

1. Let $\Sigma = \{a, b\}$. For the languages that are defined by each of the following grammars, do each of the following:

- List five strings that are in L .
- List five strings that are not in L (or as many as there are, whichever is greater).
- Describe L concisely. You can use regular expressions, expressions using variables (e.g., $a^m b^n$), or set theoretic expressions (e.g., $\{x: \dots\}$).
- Indicate whether or not L is regular. Prove your answer.

- $S \rightarrow aS \mid Sb \mid \epsilon$
- $S \rightarrow aSa \mid bSb \mid a \mid b$
- $S \rightarrow aS \mid bS \mid \epsilon$
- $S \rightarrow aS \mid aSbS \mid \epsilon$

2. Let G be the grammar of Example 11.12. Show a third parse tree that G can produce for the string $(())()$.

3. Consider the following grammar G :

$$S \rightarrow 0S1 \mid SS \mid 10$$

Show a parse tree produced by G for each of the following strings:

- 010110.
- 00101101.

4. Consider the following context free grammar G :

$$S \rightarrow aSa$$

$$\begin{aligned} S &\rightarrow T \\ S &\rightarrow \epsilon \\ T &\rightarrow bT \\ T &\rightarrow cT \\ T &\rightarrow \epsilon \end{aligned}$$

One of these rules is redundant and could be removed without altering $L(G)$. Which one?

6. Show a context-free grammar for each of the following languages L :

- BalDelim = $\{w : \text{where } w \text{ is a string of delimiters: } (,), [,], \{, \}, \text{ that are properly balanced}\}$.
- $\{a^i b^j : 2i = 3j + 1\}$.
- $\{a^i b^j : 2i \neq 3j + 1\}$.
- $\{w \in \{a, b\}^* : \#_a(w) = 2 \cdot \#_b(w)\}$.
- $L = \{w \in \{a, b\}^* : w = w^R\}$.
- $\{a^i b^j c^k : i, j, k \geq 0 \text{ and } (i \neq j \text{ or } j \neq k)\}$.
- $\{a^i b^j c^k : i, j, k \geq 0 \text{ and } (k \leq i \text{ or } k \leq j)\}$.
- $\{w \in \{a, b\}^* : \text{every prefix of } w \text{ has at least as many a's as b's}\}$.
- $\{a^n b^m : m \geq n, m-n \text{ is even}\}$.
- $\{a^m b^n c^p d^q : m, n, p, q \geq 0 \text{ and } m + n = p + q\}$.
- $\{x c^n : x \in \{a, b\}^* \text{ and } (\#_a(x) = n \text{ or } \#_b(x) = n)\}$.
- $\{b \# b_i \# 1^R : b_i \text{ is the binary representation of some integer } i, i \geq 0, \text{ without leading zeros}\}$. (For example $101 \# 011 \in L$.)
- $\{x^R \# y : x, y \in \{0, 1\}^* \text{ and } x \text{ is a substring of } y\}$.

9. (t-6) Show the details of the definition of a CFG that generates $\{a^i b^j c^k : i, j, k \geq 0 \text{ and } (i + j = k)\}$

8. Consider the unambiguous expression grammar G' of Example 11.19.

- Trace a derivation of the string $id + id * id * id$ in G' .
- Add exponentiation ($**$) and unary minus ($-$) to G' , assigning the highest precedence to unary minus, followed by exponentiation, multiplication, and addition, in that order.

9. Let $L = \{w \in \{a, b, \cup, \epsilon, (,), *, +\}^* : w \text{ is a syntactically legal regular expression}\}$.

- Write an unambiguous context-free grammar that generates L . Your grammar should have a structure similar to the arithmetic expression grammar G' that we presented in Example 11.19. It should create parse trees that:

- Associate left given operators of equal precedence, and
- Correspond to assigning the following precedence levels to the operators (from highest to lowest):
 - $*$ and $+$
 - concatenation
 - \cup

- Show the parse tree that your grammar will produce for the string $(a \cup b) ba^*$.