## 474 HW 08 problems (highlighted problems are the ones to turn in)

	1. For each of the following languages $L$ , state whether $L$ is regular or not and
9 loogaltintur	prove your answer:
8. Tacegriptuz	<b>a.</b> $\{a:b^{j}: i, j \ge 0 \text{ and } i + j = 5\}$ . <b>b.</b> $\{s^{j}b^{j}: i, j \ge 0 \text{ and } i - j = 5\}$ .
(#1)	<b>c.</b> $\{a^{j}b^{j}: i, j \ge 0 \text{ and }  i - j  = s\}$ .
	<b>d.</b> $\{w \in \{0, 1, \#\}^* : w = x \# y, \text{ where } x, y \in \{0, 1\}^* \text{ and }  x  \cdot  y  \equiv 0\}.$
	e. $\{\mathbf{a}^i \mathbf{b}^j : 0 \le i < j < 2000\}.$
<mark>8.1 e (</mark> #2)	<b>f.</b> $\{w \in \{Y, N\}^* : w \text{ contains at least two Y's and at most two N's}\}.$
	g. $\{w = xy : x, y \in \{a, b\}^* \text{ and }  x  =  y  \text{ and } \#_a(x) \ge \#_a(y)\}.$
<mark>8.1f</mark>	<b>h.</b> $\{w = xyzy^{R}x : x, y, z \in \{a, b\}^*\}.$
	i. $\{w = xyzy : x, y, z \in \{0, 1\}^+\}$ .
<mark>8.1g</mark>	j. $\{w \in \{0,1\}^* : \#_0(w) \neq \#_1(w)\}.$
	<b>k.</b> $\{w \in \{a, b\}^* : w = w^\kappa\}$ .
<mark>8.1h</mark>	1. $\{w \in \{a, b\}^* : \exists x \in \{a, b\}^+ (w = x x^n x)\}$ .
0.41	<b>m.</b> $\{w \in \{a, b\}^* : \text{the number of occurrences of the substring ab equals the num ber of occurrences of the substring ba}.$
<mark>8.1K</mark>	<b>n.</b> $\{w \in \{a, b\}^* : w \text{ contains exactly two more b's than a's}\}$ .
<mark>0.1 m</mark>	0. $\{w \in \{a, b\}^* : w = xyz,  x  =  y  =  z , \text{ and } z = x \text{ with every a replaced by}$
<mark>8.1m</mark>	b and every b replaced by a}. Example: abbbabbaa $\in L$ , with $x =$
	abb, y = bab, and z = baa.
(#2 - #7)	<b>p.</b> $\{w: w \in \{a - z\}^* \text{ and the letters of } w \text{ appear in reverse alphabetical order}\}.$
<mark>3 pts. each</mark>	For example, spoonteed $\in L$ .
	q. $\{w: w \in \{a - z\}^* \text{ every letter in } w \text{ appears at least twice}\}$ . For example,
	unprosperousness $\in L$ .
	<b>r.</b> $\{w : w \text{ is the decimal encoding of a natural number in which the digits appear in a non-decreasing order without leading zeros}.$
	s. {w of the form: $\langle integer_1 \rangle + \langle integer_2 \rangle = \langle integer_2 \rangle$ , where each of the
	substrings $\langle integer_1 \rangle$ , $\langle integer_2 \rangle$ , and $\langle integer_3 \rangle$ is an element of $\{0 -$
	9}* and integer <sub>3</sub> is the sum of integer <sub>1</sub> and integer <sub>2</sub> }. For example,
	$124+5=129 \in L.$
	t. $L_0^*$ , where $L_0 = \{ ba^i b^j a^k, j \ge 0, 0 \le i \le k \}.$
	<b>u.</b> $\{w : w \text{ is the encoding of a date that occurs in a year that is a prime number}\}$ .
	and v is drawn from $\{0-9\}$ .
	<b>v.</b> $\{w \in \{1\}^* : w \text{ is, for some } n \ge 1, \text{ the unary encoding of } 10^n\}$ . (So $L =$
	$\{1111111111, 1^{100}, 1^{1000}, \dots\}.\}$
	6. Prove by construction that the regular languages are closed under:
	a. intersection.
	b. set difference.
	7. Prove that the regular languages are closed under each of the following operations:
	<b>a.</b> $pref(L) = \{w : \exists x \in \Sigma^*(wx \in L)\}.$
<mark>8.7a</mark>	<b>b.</b> $suff(L) = \{w : \exists x \in \Sigma^*(xw \in L)\}.$
(#8) <mark>9</mark>	<b>c.</b> $reverse(L) = \{x \in \Sigma^* : x = w^* \text{ for some } w \in L\}.$
	<b>u.</b> letter substitution (as defined in Section 8.5).
	<b>8.7a</b> Do this by construction, i.e., produce an algorithm that takes as input a DFSM
	$M = (K \sum \delta c A)$ that accents L and produces a DESM $M' = (K' \sum \delta' c' A')$ that accents
	[1] $[1, 2]$ $[0, 3]$ $[1]$
	<i>prej(L)</i> . Describe now to get from M to M
	<b>HINT:</b> MI will have a lot of its elements in common with M, but it takes a somewhat
	complex calculation (based on M) to determine exactly what has to be changed.
	On the main HW8 assignment document. I posted the author's solutions to the other
	three nexts of problem 8.7 so that you will have more events is
	uree parts of problem 6.7, so that you will have more examples.

8.2ac	2. For each of the following languages L state whether L is regular or not and prove
(#9)	your answer:
	<b>a.</b> $\{w \in \{a, b, c\}^* : \text{ in each prefix } x \text{ of } w, \#_a(x) = \#_b(x) = \#_c(x))\}.$
	<b>b.</b> $\{w \in \{a, b, c\}^* : \exists \text{ some prefix } x \text{ of } w(\#_a(x) = \#_b(x) = \#_c(x))\}.$
	<b>c.</b> $\{w \in \{a, b, c\}^{\infty}: \exists$ some prenx x of $w (x \neq \varepsilon \text{ and } \#_a(x) = \#_b(x) = \#_c(x))\}$ . <b>3.</b> Define the following two languages:
8.3	$L_{i} = \{ w \in \{a, b\}^* : \text{ in each prefix } x \text{ of } w, \#_i(x) \ge \#_i(x) \}.$
(#10)	$L_{a} = \{w \in \{a, b\}^{*} : \text{ in each prefix } x \text{ of } w, \#_{a}(x) \ge \#_{a}(x)\}.$
	<b>a.</b> Let $I_{4} = I_{4} \cap I_{4}$ . Is $I_{4}$ regular? Prove your answer
	<b>b.</b> Let $L_2 = L_a \cup L_b$ . Is $L_2$ regular? Prove your answer.
<mark>8.4a</mark>	
(#11) <mark>6</mark>	4. For each of the following languages $L$ , state whether $L$ is regular or not and prove
	your answer: <b>a</b> $\int u u u r R^{R} = \int u u u r R^{-1} \int u r R^{-1} dr$
8.4b	<b>h.</b> $\{xvzv^{R}x: x, v, z \in \{a, b\}^{+}\}$ .
(#12)	
-	7. Prove that the regular languages are closed under each of the following operations:
8.7	<b>a.</b> $pref(L) = \{w : \exists x \in \Sigma^*(wx \in L)\}.$
(#13)	<b>b.</b> suff $(L) = \{ w : \exists x \in \Sigma^* (xw \in L) \}.$
	<b>c.</b> $reverse(L) = \{x \in \Sigma^* : x = w^{\mathbb{R}} \text{ for some } w \in L\}.$
	d. letter substitution (as defined in Section 8.3).
	8. Using the definitions of maxstring and mix given in Section 8.6, give a precise def-
	inition of each of the following languages:
	<b>a.</b> maxstring( $A^n B^n$ ).
<b>a a</b>	<b>b.</b> maxstring ( $a^i b^j c^k, 1 \le k \le i \le i$ ).
<mark>6.18</mark>	<b>c.</b> maximum $(L_1 L_2)$ , where $L_4 = \{w \in \{a, b\}^* : w \text{ contains exactly one } a\}$ and
(#14) <mark>9</mark>	$L_2 = \{a\}.$
	<b>d.</b> $mix((aba)^*)$ .
	<b>e.</b> $mix(a*b*)$ .
	9. Prove that the regular languages are not closed under <i>mix</i> .
6.20	J. Trove that the regular tangaages are not crosed ander man
0.20	18. Let $\Sigma$ = {a, b}. Let L = { $\varepsilon$ , a, b}. Let R be a relation defined on $\Sigma^*$ as follows: $\forall xy$ (xRy iff y = xb). Let R' be
(#15)	the reflexive, transitive closure of R. Let $L' = \{x : \exists y \in L (yR'x)\}$ . Write a regular expression for L'.
Good	
nractica	Note on 6.18 Transitive and reflexive closures are introduced
practice	in Section A 5 Closures under various operations are also
problems	mentioned on pages 17, 57, and 72
for exams	mentioned on pages 17, 37, and 72.
(no proof	
necessary)	
	20. For each of the following statements state whether it is <i>True</i> or <i>False</i> . Prove your
	answer.
	<b>a.</b> $(ab)^*a = a(ba)^*$ .
	<b>b.</b> $(a \cup b)^* b (a \cup b)^* = a^* b (a \cup b)^*$ .
	c. $(a \cup b)^* b (a \cup b)^* \cup (a \cup b)^* a (a \cup b)^* = (a \cup b)^*$ .
	d. $(a \cup b)^* b (a \cup b)^* \cup (a \cup b)^* a (a \cup b)^* = (a \cup b)^+$ .
	e. $(a \cup b)^{\infty} ba(a \cup b)^{\infty} \cup a^{\infty} b^{\infty} = (a \cup b)^{\alpha}$ .

- **f.**  $a^* b (a \cup b)^* = (a \cup b)^* b (a \cup b)^*$ .
- **g.** If  $\alpha$  and  $\beta$  are any two regular expressions, then  $(\alpha \cup \beta)^* = \alpha (\beta \alpha \cup \alpha)$ . **h.** If  $\alpha$  and  $\beta$  are any two regular expressions, then  $(\alpha \beta)^* \alpha = \alpha (\beta \alpha)^*$ .