474 HW 2 problems (highlighted problems are the ones to turn in)

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|------------------------|--|---|--|
| 3.1 | Consider the following problem: Given a digital circuit C, does C output 1 on all inputs? Describe this problem as a language to be decided. | | |
| <mark>3.2</mark> | Using the technique we used in Example 3.8 to describe addition, describe square root as a language recognition problem. | | |
| 3.2 | 3. Consider the problem of encrypting a password, given an encryption key. Formulate this problem as a language recognition problem. | | |
| | 4. Consider the optical character recognition (OCR) problem: Given an array of | | |
| 2.4 | black and white pixels and a set of characters, determine which character best | | |
| 3.4 | matches the pixel array. Formulate this problem as a language recognition problem. | | |
| <mark>3.5</mark> | 5. Consider the language $A^nB^nC^n = \{a^nb^nc^n : n \ge 0\}$, discussed in Section 3.3.3. We might consider the following design for a PDA to accept $A^nB^nC^n$: As each a | | |
| | | | |
| | is read, push two a's onto the stack. Then pop one a for each b and one a for each c. If the input and the stack come out even, accept. Otherwise reject. Why doesn't this work? | | |
| | 6. Define a PDA-2 to be a PDA with two stacks (instead of one). Assume that the | | |
| <mark>3.6</mark> | stacks can be manipulated independently and that the machine accepts iff it is | | |
| | in an accepting state and both stacks are empty when it runs out of input. De- | | |
| | scribe the operation of a PDA-2 that accepts $A^nB^nC^n = \{a^nb^nc^n : n \ge 0\}$. (<i>Note</i> : We will see, in Section 17.5.2, that the PDA-2 is equivalent to the Turing | | |
| | | machine in the sense that any language that can be accepted by one can be ac- | |
| | cepted by the other.) | | |
| | 1. Describe in clear English or pseudocode a decision procedure to answer the | | |
| 4.1 | question, "Given a list of integers N and an individual integer n , is there any element of N that is a factor of n ?" | | |
| <mark>4.2</mark> | 2. Given a Java program p and the input 0, consider the question, "Does p ever output anything?" | | |
| | a. Describe a semidecision procedure that answers this question. | | |
| | b. Is there an obvious way to turn your answer to part a into a decision procedure? | | |
| <mark>4.3</mark> , 4.4 | 3. Recall the function $chop(L)$, defined in Example 4.10. Let $L = \{w \in \{a, b\}^* : w = w^R\}$. What is $chop(L)$? | | |
| | 4. Are the following sets closed under the following operations? Prove your answer. If a set is not closed under the operation, what is its closure under the operation? | One more problem, | |
| 4.4c | a. $L = \{w \in \{a, b\}^* : w \text{ ends in } a\}$ under the function $odds$, defined on strings as follows: $odds(s) = \text{the string that is formed by concatenating together all of the odd numbered characters of s. (Start numbering the characters at 1.) For example, odds(\text{ababbbb}) = \text{aabb}.$ | not from the textbook: | |
| | b. FIN (the set of finite languages) under the function <i>oddsL</i> , defined on lan- | #14 is described in detail on | |
| | guages as follows: | the assignment document, so I | |
| | $oddsL(L) = \{w : \exists x \in L (w = odds(x))\}.$ | | |
| | c. INF (the set of infinite languages) under the function oddsL. | do not repeat that description | |
| | d. FIN under the function <i>maxstring</i> , defined in Example 8.22. | here. | |
| | e. INF under the function maxstring. | | |
| | 2. Show a DFSM to accept each of the following languages: | | |
| | a. {w∈ {a,b}*: every a in w is immediately preceded and followed by b}. b. {w∈ {a,b}*: w does not end in ba}. | | |
| 5.2 | c. $\{w \in \{0,1\}^* : w \text{ corresponds to the binary encoding, without leading 0's, of natural numbers that are evenly divisible by 4}.$ | | |
| <mark>5.2a</mark> | d. $\{w \in \{0,1\}^* : w \text{ corresponds to the binary encoding, without leading 0's, of natural numbers that are powers of 4}.$ | | |
| 5.2b | e. {w ∈ {0-9}*: w corresponds to the decimal encoding, without leading 0's, of an odd natural number}. f. {w ∈ {0,1}*: w has 001 as a substring}. | | |
| | g. $\{w \in \{0,1\}^* : w \text{ does not have 001 as a substring}\}$. | | |
| | h. $\{w \in \{a,b\}^* : w \text{ has bbab as a substring}\}.$ | | |
| | i. $\{w \in \{a, b\}^* : w \text{ has neither ab nor bb as a substring}\}.$ | | |
| | j. $\{w \in \{a,b\}^* : w \text{ has both aa and bb as a substrings}\}.$ | | |
| | k. $\{w \in \{a, b\}^* : w \text{ contains at least two b's that are not immediately followed by an a}\}$. | | |
| | 1. $\{w \in \{0,1\}^* : w \text{ has no more than one pair of consecutive 0's and no more than one pair of consecutive 1's}.$ | | |
| | m (an = [0, 1]* ; none of the profives of an ends in 0] | | |

 $\mathbf{m} \cdot \{ w \in \{0,1\}^* : \text{ none of the prefixes of } w \text{ ends in } 0 \}.$

 $\mathbf{n.} \ \{ w \in \{ \mathtt{a}, \mathtt{b} \}^* : (\#_{\mathtt{a}}(w) \ + \ 2 \cdot \#_{\mathtt{b}}(w)) \equiv_{5} 0 \}. \ (\#_{\mathtt{a}}(w) \text{ is the number of a's in } w).$