474 HW 5 problems (highlighted problems are the ones to turn in)

Problems 1-3 from HW5 are not from the textbook, and they are described in much detail on the main HW5, so they are nopt shown here.

	1. Describe in English, as briefly as possible, the language defined by each of these	
	regular expressions:	6.2: For the "to-turn-in"
6.1	a. $(b \cup ba) (b \cup a)^* (ab \cup b).$	
4	b. $(((a^*b^*)^*ab) \cup ((a^*b^*)^*ba))(b \cup a)^*.$	parts of 6.2, aim for as
	2. Write a regular expressions to describe each of the following languages:	simple a regular expression
<mark>6.1a</mark>	a. $\{w \in \{a, b\}^* : \text{ every } a \text{ in } w \text{ is immediately preceded and followed by } b\}$.	as you can come up with.
	b. $\{w \in \{a, b\}^* : w \text{ does not end in ba}\}$.	Some of the credit may be
	c. $\{w \in \{0, 1\}^* : \exists y \in \{0, 1\}^* (xy \text{ is even})\}.$	for simplicity. If your
	d. {w ∈ {0, 1}* : w corresponds to the binary encoding, without leading 0s, of natural numbers that are evenly divisible by 4}.	expression is very
<mark>6.5</mark>	e. $\{w \in \{0, 1\}^* : w corresponds to the binary encoding, without leading 0s, of the binary encoding is the set of the binary encoding is the set of the binary encoding is the bin$	complicated, some
<mark>10(6-6)</mark>	natural numbers that are powers of 4}. Claude Anderson	annotation may help the
	f. $\{w \in \{0.9\}^* : w \text{ corresponds to the decimal encoding, without leading 0s, of } \}$	grader to know whether it
	an odd natural number}.	is correct. The "burden of
	g. $\{w \in \{0, 1\}^* : w \text{ has 001 as a substring} \}$.	correctness" is on you. I
<mark>6.2b</mark>	h. $\{w \in \{0, 1\}^* : w \text{ does not have 001 as a substring}\}$.	did not require some of the
<mark>7 (3)</mark>	i. $\{w \in \{a, b\}^* : w \text{ has bba as a substring}\}$.	more complex parts of this
	j. $\{w \in \{a, b\}^* : w \text{ has both } aa \text{ and } bb \text{ as substrings}\}.$	problem due to difficulties
	k. $\{w \in \{a, b\}^* : w \text{ has both aa and aba as substrings}\}.$	with grading, but you
<mark>6.2c</mark>	1. $\{w \in \{a, b\}^* : w \text{ contains at least two b's that are not followed by an } a\}$.	should try some of them.
<mark>8 (3)</mark>	m. $\{w \in \{0, 1\}^* : w has at most one pair of consecutive 0s and at most one pair$	should up some of them.
See	of consecutive 1s}.	Note for part c: The
"note for	n. $\{w \in \{0, 1\}^* : \text{none of the prefixes of } w \text{ ends in } 0\}.$	L
	1. $\{w \in \{0, 1\}^{*} : \#_{a}(w) \equiv_{3} 0\}.$	<i>x</i> is a typo, should be <i>w</i> .
part c" in	p. $\{w \in \{a, b\}^* : \#_a(w) = 30\}$.	wy means "the length of
<mark>the box</mark>	q. $\{w \in \{a, b\}^* : w \text{ contains exactly two occurrences of the substring aa}\}.$	string wy"
<mark>on the</mark>	r. $\{w \in \{a, b\}^* : w \text{ contains no more than two occurrences of the substring}\}$	
right.	aa}.	
1.8.1C.	s. $\{w \in \{a, b\}^* - L\}$, where $L = \{w \in \{a, b\}^* : w \text{ contains bba as a}$	
	substring}.	
	t. { $w \in \{0, 1\}^*$: every odd length string in <i>L</i> begins with 11}.	
	u. $\{w \in \{0.9\}^* : w represents the decimal encoding of an odd natural number without leading 0s.$	
	v. $L_1 - L_2$, where $L_1 = a^*b^*c^*$ and $L_2 = c^*b^*a^*$.	
	w. The set of legal United States zip codes \square .	
	x. The set of strings that correspond to domestic telephone numbers in your	
	country.	
	3. Simplify each of the following regular expressions:	
	a. $(\mathbf{a} \cup \mathbf{b})^* (\mathbf{a} \cup \varepsilon) \mathbf{b}^*$.	
	b. $(\emptyset^* \cup b) b^*$.	
	c. (a ∪ b)*a* ∪ b.	
	d. $((a \cup b)^*)^*$.	
	e. $((a \cup b)^+)^*$.	
	f. $a((a \cup b)(b \cup a))^* \cup a((a \cup b)a)^* \cup a((b \cup a)b)^*$.	
	4. For each of the following expressions E, answer the following three questions	
	and prove your answer:	
	i. Is E a regular expression?ii. If E is a regular expression, give a simpler regular expression.	
	iii. Does <i>E</i> describe a regular language?	
	a. $((a \cup b) \cup (ab))^*$.	
	b. $(a^+ a^n b^n)$.	
	c. ((a b)*∅).	
	d. $((ab) \cup c)^* \cap (b \cup c^*)$.	
	e. $(\emptyset^* \cup (bb^*)).$	
	- (~ - ()).	