

## 474 Notes on Day 7 slides:

Some of the early slides are repeats of Day 4; I do not duplicate those notes here.

Slide 3: Recap: Definition of a DFSA

Don't spend much time on these recap slides, except the yields slides.

These slides are mainly here for reference as we do other things.

Slide 85: **Recap: Accepting and Rejecting**

Sketch of proof.

Base case: If  $w$  is  $\epsilon$ , it halts in 0 steps.

Assume true for strings of length  $n$  and show for strings of length  $n+1$ .

Let  $w \in \Sigma^*$ ,  $w \neq \epsilon$   $|w| = n+1$  for some  $n \in \mathbb{N}$ .

Then  $w$  is  $ax$  for some  $a \in \Sigma$ ,  $x \in \Sigma^*$ ,  $|x| = n$ .

Let  $q'$  be  $\delta(q, a)$ . Then  $(q, w) \xrightarrow{-M} (q', x)$

By induction, from configuration  $(q', x)$ ,  $M$  halts in  $n$  steps.

So, starting from the original configuration,  $M$  halts in  $n+1$  steps.

**ASK: IF  $M$  is a DFSA, is membership in  $L(M)$  decidable?**

Slide 10: Example

Sometimes instead of cluttering up a transition diagram by showing a dead state and lots of transitions into it, we will simply not show a transition out of some state on a given input symbol.

This is a "shorthand" for a dead state.

Redraw the diagram!

Slide 15: (Hidden) Solution

Note that I fixed this slide; the author's had the dead state as an accepting state.

If a new version of PowerPoint comes out, I should make sure that my "coverup" still works

Slide 20: Vowels in Alphabetical Order

I fixed an author's error by covering up an  $a$  with a  $u$ , then grouping.

Slide 33: Pattern Matching: Multiple Keywords

*Note that states  $q4$  and  $q8$  could be combined.*