19.1
(\#1) 6

1. Consider the language $L=\{\langle M\rangle$ : Turing machine $M$ accepts at least two strings $\}$.
a. Describe in clear English a Turing machine $M$ that semidecides $L$.
b. Now change the definition of $L$ just a bit. Consider:
$L^{\prime}=\langle<M\rangle$ :Turing machine $M$ accepts exactly 2 strings $\rangle$.
Can you tweak the Turing machine you described in part a to semidecide $L^{\prime}$ ?
2. Consider the language $L=\{\langle M\rangle$ : Turing machine $M$ accepts the binary encodings of the first three prime numbers $\}$.
a. Describe in clear English a Turing machine $M$ that semidecides $L$.
b. Suppose (contrary to fact, as established by Theorem 19.2) that there were a Turing machine Oracle that decided H. Using it, describe in clear English a Turing machine $M$ that decides $L$.
3. Show that the set D (the decidable languages) is closed under:
a. Union
b. Concatenation
c. Kleene star
d. Reverse
c. Intersection
4. Show that the set SD (the semidecidable languages) is closed under:
a. Union
b. Concatenation
c. Klene star
d. Reverse
e. Intersection
5. Let $L_{1}, L_{2}, \ldots, L_{k}$ be a collection of languages over some alphabet $\Sigma$ such that:

- For all $i \neq j, L_{i} \cap L_{j}=\varnothing$.
- $L_{1} \cup L_{2} \cup \ldots \cup L_{k}=\Sigma^{*}$.
- $\forall i\left(L_{i}\right.$ is in SD).

Prove that each of the languages $L_{1}$ through $L_{k}$ is in D.
4. If $L_{1}$ and $L_{3}$ are in D and $L_{1} \subseteq L_{2} \subseteq L_{3}$, what can we say about whether $L_{2}$ is in D ?
5. Let $L_{1}$ and $L_{2}$ be any two decidable languages. State and prove your answer to each of the following questions:
a. Is it necessarily true that $L_{1}-L_{2}$ is decidable?
b. Is it possible that $L_{1} \cup L_{2}$ is regular?
6. Let $L_{1}$ and $L_{2}$ be any two undecidable languages. State and prove your answer to each of the following questions:
a. Is it possible that $L_{1}-L_{2}$ is regular?
b. Is it possible that $L_{1} \cup L_{2}$ is in D?
7. Let $M$ be a Turing machine that lexicographically enumerates the language $L$. Prove that there exists a Turing machine $M^{\prime}$ that decides $L^{\mathrm{R}}$.
8. Construct a standard one-tape Turing machine $M$ to enumerate the language:
$\{w: w$ is the binary encoding of a positive integer that is divisible by 3$\}$.
Assume that $M$ starts with its tape equal to $\perp$. Also assume the existence of the printing subroutine $P$, defined in Section 20.5.1. As an example of how to use $P$, consider the following machine, which enumerates $L^{\prime}$, where $L^{\prime}=\{w: w$ is the unary encoding of an even number $\}$ :

13. Show that every infinite semidecidable language has a subset that is not decidable.

