

## 474 Instructor Notes from Day 14 slides:

Slide 3: How many languages are there?

The answer is no, so  $L_d$  is not in the enumeration, and the set of languages is uncountable.

Write "Floyd-Warshall" on the board.

Slide 4: **How Many Regular Languages?**

How do we know that there are countably many regular expressions over a given alphabet?

We can easily enumerate them, based on the number of rules involved in the construction.

Slide 8: **Regular Does Not Always Mean Tractable**

How many states are there? Each disk is on one of the three poles. The order of the disks on a given pole is fixed. So there are  $3^{64}$  states.

Slide 10: **Closure Properties of Regular Languages**

### Reverse:

By construction. Let  $M = (K, \Sigma, \delta, s, A)$  be any FSM that accepts  $L$ .  $M$  must be written out completely, without an implied dead state. Then construct  $M' = (K', \Sigma', \delta', s', A')$  to accept  $reverse(L)$  from  $M$ :

Initially, let  $M'$  be  $M$ .

Reverse the direction of every transition in  $M'$ .

Construct a new state  $q$ . Make it the start state of  $M'$ . Create an  $\epsilon$ -transition from  $q$  to every state that was an accepting state in  $M$ .

$M'$  has a single accepting state, the start state of  $M$ .

Slide 11: Closure of Regular Languages under intersection

$$L(M1) \cap L(M2) = \neg(\neg L(M1) \cup \neg L(M2)).$$

In the homework, you will directly construct a machine to do the intersection

Slide 12: **Closure of Regular Languages Under Difference**

$$L(M1) - L(M2) = L(M1) \cap \neg L(M2).$$

Slide 16: **Don't try to use closure backwards**

Also intersection of  $\{a^n b^n\}$  and  $\{a^n b^{n+1}\}$  is empty, hence regular.