

# Day 34

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POWER

## Tackling NP Problems

Two principal approaches to tackling problems in NP:

- Use a strategy that guarantees solving the problem exactly but doesn't guarantee to find a solution in polynomial time
- Use an *approximation algorithm* that can find an approximate (sub-optimal) solution in polynomial time

## Exact Solution Strategies

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*Exhaustive search* (brute force).

- useful only for small instances

*Dynamic programming.*

- applicable to some problems (e.g., the knapsack problem)

*Backtracking*

- eliminates some unnecessary cases from consideration,
- yields solutions in reasonable time for many instances, but worst case is still exponential

*branch-and-bound.*

- further refines the backtracking idea for optimization problems

## Backtracking

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When applicable, backtracking is often much faster than brute force enumeration, since it eliminates many candidates with a single test

Process:

- Construct the state-space tree
  - nodes: partial solutions
  - edges: choices in extending partial solutions
- Explore the state space tree using depth-first search
- “Prune” nonpromising nodes
  - stop exploring subtrees rooted at nodes that cannot lead to a solution
  - backtrack to such a node’s parent to continue the search

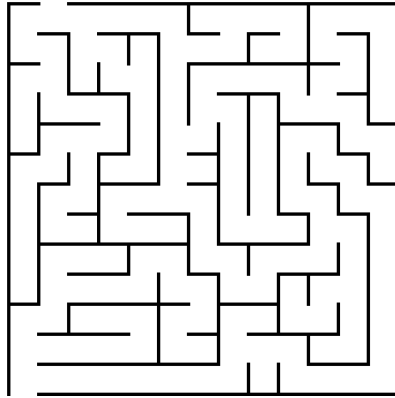
# State-Space Tree of the 4-Queens Problem

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The diagram illustrates the state-space tree for the 4-Queens problem. The root node is labeled 0 and shows an empty 4x4 grid. It branches into two nodes: 1 (left) and 5 (right). Node 1 has a queen in the first row, first column. It branches into three nodes: 2 (left), 3 (middle), and 4 (right). Node 2 has a queen in the second row, second column. Node 3 has a queen in the second row, third column. Node 4 has a queen in the second row, fourth column. Node 5 has a queen in the first row, second column. It branches into two nodes: 6 (left) and 7 (right). Node 6 has a queen in the second row, first column. Node 7 has a queen in the second row, third column. Node 2 branches into three nodes: 1 (left), 2 (middle), and 3 (right). Node 3 branches into two nodes: 4 (left) and 5 (right). Node 4 branches into one node: 1 (left). Node 5 branches into one node: 2 (left). Node 6 branches into one node: 7 (left). Node 7 branches into one node: 8 (left). Node 8 is labeled "solution" and shows a 4x4 grid with queens in the first row, second column, third row, first column, and fourth row, fourth column.

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graph TD; 0[0] --> 1[1]; 0 --> 5[5]; 1 --> 2[2]; 1 --> 3[3]; 1 --> 4[4]; 5 --> 6[6]; 5 --> 7[7]; 2 --> 1[1]; 2 --> 2[2]; 2 --> 3[3]; 3 --> 4[4]; 3 --> 5[5]; 4 --> 1[1]; 6 --> 7[7]; 7 --> 8[8]; 8 --> solution[solution];
```

## Familiar Problems



5	3			7				
6			1	9	5			
	9	8					6	
8				6				3
4			8		3			1
7				2				6
	6					2	8	
			4	1	9			5
				8			7	9

## Backtracking Notes

Generally just marginally better than exhaustive search (still asymptotically inefficient in the worst case)

Can exploit symmetries to further reduce cases

Rearrange data (e.g. presorting max remaining sums for Subset Sum)

Can estimate size of relevant state-space by generating random path from root to a leaf and counting the number of choices at each step

## Branch-and-Bound

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An enhancement of backtracking

For each node (partial solution) of a state-space tree, compute a bound on the value of the objective function for all descendants of the node (extensions of the partial solution)

Use the bound for:

- ruling out certain nodes as “nonpromising” to prune the tree – if a node’s bound is **not better than the best solution** seen so far
- guiding the search heuristically through state-space

## Branch-and-Bound in Playing Games

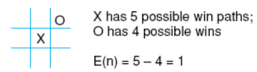
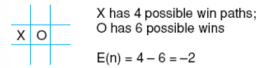
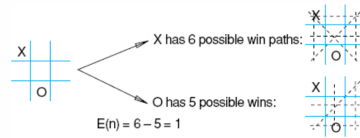
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Chess, Go, Tic-Tac-Toe variants, etc.

Minimax rule: **minimize** the possible loss for a worst-case (**max** loss) scenario (i.e., max the min gain)

- For each state we compute a “score”
  - Ideally,  $\infty$  for a winning position and  $-\infty$  for a losing position. But, usually too hard to compute this.
  - More practically, stop at base cases where we use a heuristic measure of “quality of position”.
- Then, ancestors are
  - If my move (max): max of children’s scores
  - If opponent’s move (min): min of children’s scores

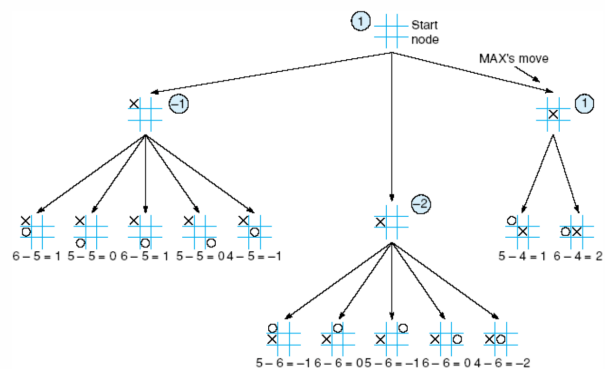
# MinMax for TicTacToe



Heuristic is  $E(n) = M(n) - O(n)$   
 where  $M(n)$  is the total of My possible winning lines  
 $O(n)$  is total of Opponent's possible winning lines  
 $E(n)$  is the total Evaluation for state  $n$

Luger: Artificial Intelligence, 6th edition. © Pearson Education Limited, 2009

# MinMax for TicTacToe



Luger: Artificial Intelligence, 6th edition. © Pearson Education Limited, 2009

# Alpha-Beta Pruning

Branch-and-bound for minimax algorithm

- $\alpha$ : (maximum) lower bound on score of possible moves so far
- $\beta$ : (minimum) upper bound on score of possible moves so far

Min nodes: opponent's moves. Take  $\beta = \min$  of children's  $\alpha$ 's.

Max nodes: my moves. Take  $\alpha = \max$  of children's  $\beta$ 's.

Stop evaluating a move when at least one possibility has been found that proves the move to be worse than a previously examined move

E.g. if current  $\alpha$  for max node is  $\geq$  the  $\beta$  at a child min node, then nonpromising and can prune.

# Alpha-Beta Pruning

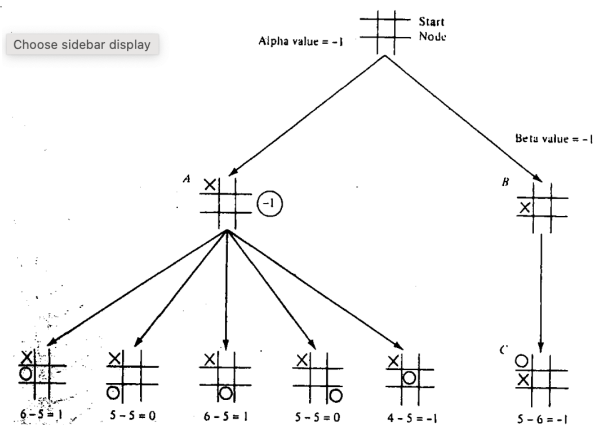


Fig. 3.11 Part of the first stage tic-tac-toe tree.

Image source: Nilson: Principles of AI, Tioga

## Approximation Approach

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Apply a fast (i.e., polynomial-time) approximation algorithm to get a solution that is not necessarily optimal but hopefully close to it

Accuracy measures:

- Change goes beyond a certain delta
- Alternatively, limit by CPU cycles and take and run

## Approximation Approach

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Genetic Algorithms for TSP

[https://www.heatonresearch.com/aifh/vol2/tsp\\_genetic.html](https://www.heatonresearch.com/aifh/vol2/tsp_genetic.html)

Ant-colony optimization:

<https://baobabsoluciones.es/en/blog/2020/10/01/travelling-salesman-problem-methods/>