

## MA/CSSE 473 Day 02 Announcements and Summary

### Announcements:

1. Background survey on Moodle. Complete it by 5 PM TODAY.
2. HW1 Due Monday. A bit smaller than HW 2, due Thursday, Sept. 11. I suggest that you begin HW 2 before HW 1 is due.
3. In this course, I value giving you experience with many algorithms and design techniques more than I value any particular algorithms or design techniques.

### Main ideas from today:

1. Master Theorem:
  - a. Consider the recurrence  $T(n) = aT(n/b) + f(n)$ ,  $T(1)=c$ , where  $f(n) \in \Theta(n^k)$  and  $k \geq 0$ .
  - b. The solution is in
    - i.  $\Theta(n^k)$  if  $a < b^k$
    - ii.  $\Theta(n^k \log n)$  if  $a = b^k$
    - iii.  $\Theta(n^{\log_b a})$  if  $a > b^k$
2. For what kind of algorithms does the Master Theorem provide a tool for analysis?
3. What are the  $a$ ,  $b$ , and  $k$  in the Master Theorem for:  
merge sort       $a=$        $b=$        $k=$       asymptotic solution:  
  
binary search       $a=$        $b=$        $k=$       asymptotic solution:
4. Fibonacci numbers:  $F(0) = 0$ ,  $F(1) = 1$ .  $F(n) = F(n-1) + F(n-2)$  if  $n \geq 2$ .
5. How do we know that the recursive computation of the  $n^{\text{th}}$  Fibonacci number,  $F(n)$ , needs at least  $F(n)$  steps? (proof by induction)
6. How many additions\_\_\_\_\_ and multiplications\_\_\_\_\_ are required in order to multiply two  $2 \times 2$  integer matrices?
7. With your group, write an efficient algorithm for finding the  $n^{\text{th}}$  power of a  $2 \times 2$  matrix, assuming that you already have a constant-time algorithm for multiplying two  $2 \times 2$  matrices at a time.

