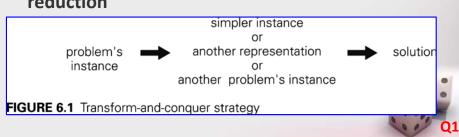


Transform and Conquer Algorithms But first: Answer student questions

- Transform a problem to a simpler instance of the same problem **instance simplification**
- Transformation to a different representation of the same instance **representation change**
- Transformation to an instance of a different problem that we know how to solve – problem reduction



Recap: Representation change: AVL Trees (what you should remember...)

- Named for authors of original paper, Adelson-Velskii and Landis (1962).
- An AVL tree is a height-balanced Binary Search Tree.
- A BST T is height balanced if T is empty, or if
 - | height(T_I) height(T_R) | ≤ 1, and
 - T_L and T_R are both height-balanced.
- Show: Maximum height of an AVL tree with N nodes is
 Θ(log N) Let's review that together
- How do we maintain balance after insertion?
- Exercise: Given a pointer to the root of an AVL tree with N nodes, find the height of the tree in log N time
- Details on balance codes and various rotations are in the CSSE 230 slides that are linked from the schedule page.

Representation change: 2-3 trees

- Another approach to balanced trees
- Keeps all leaves on the same level
- Some non-leaf nodes have 2 keys and 3 subtrees
- Others are regular binary nodes.

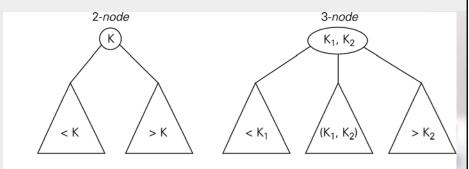
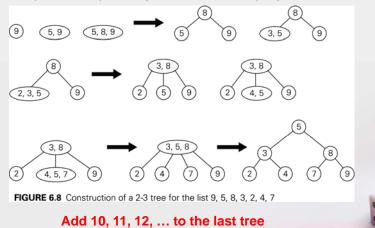


FIGURE 6.7 Two kinds of nodes of a 2-3 tree



http://slady.net/java/bt/view.php?w=450&h=300



Efficiency of 2-3 tree insertion

- Upper and lower bounds on height of a tree with n elements?
- Worst case insertion and lookup times is proportional to the height of the tree.



