

# MA/CSSE 473

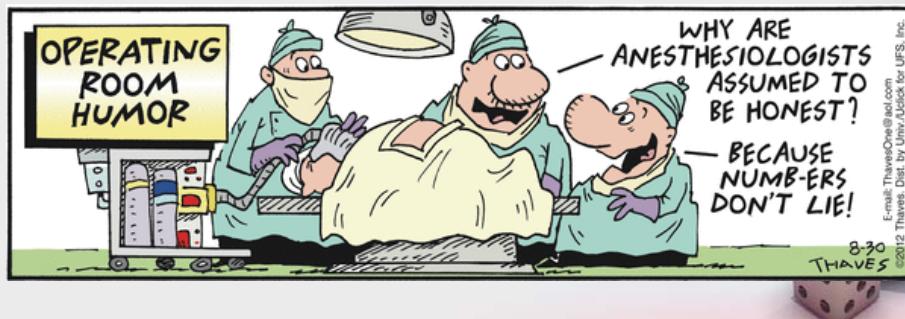
## Day 13

Permutation  
Generation



### MA/CSSE 473 Day 13

- HW 6 due Monday , HW 7 next Thursday,
- Student Questions
- Tuesday's exam
- Permutation generation



## Exam 1

- If you want additional practice problems for Friday's exam:
  - The "not to turn in" problems from various assignments
  - Feel free to post your solutions in a Piazza discussion forum and ask your classmates if they think it is correct
- Allowed for exam:  
Calculator, one piece of paper (1 sided, handwritten)
- See the exam specification document, linked from the exam day on the schedule page.



## About the exam

- Mostly it will test your understanding of things in the textbook and things we have discussed in class.
- Will not require a lot of creativity (it's hard to do much of that in 50 minutes).
- Many short questions, a few calculations.
  - Perhaps some T/F/IDK questions (example: 5/0/3)
- You may bring a calculator.
- And a piece of paper (handwritten on one side).
- I will give you the Master Theorem if you need it.
- Time will be a factor!
- First do the questions you can do quickly



## Possible Topics for Exam

- Formal definitions of O,  $\Theta, \Omega$ .
- Master Theorem
- Fibonacci algorithms and their analysis
- Efficient numeric multiplication
- Proofs by induction (ordinary, strong)
- Trominoes
- Extended Binary Trees
- Modular multiplication, exponentiation
- Extended Euclid algorithm
- Modular inverse
- Fermat's little theorem
- Rabin-Miller test
- Random Prime generation
- RSA encryption
- What would Donald (Knuth) say?



## Possible Topics for Exam

- Brute Force algorithms
- Selection sort
- Insertion Sort
- Amortized efficiency analysis
- Analysis of growable array algorithms
- Binary Search
- Binary Tree Traversals
- Basic Data Structures (Section 1.4)
- Graph representations
- BFS, DFS,
- DAGs & topological sort



Permutations

Subsets

## COMBINATORIAL OBJECT GENERATION



## Combinatorial Object Generation

- Generation of permutations, combinations, subsets.
- This is a big topic in CS
- We will just scratch the surface of this subject.
  - Permutations of a list of elements (no duplicates)
  - Subsets of a set



## Permutations

- We generate all permutations of the numbers 1..n.
  - Permutations of any other collection of n distinct objects can be obtained from these by a simple mapping.
- How would a "decrease by 1" approach work?
  - Find all permutations of 1.. n-1
  - Insert n into each position of each such permutation
  - We'd like to do it in a way that minimizes the change from one permutation to the next.
  - It turns out we can do it so that we always get the next permutation by swapping two adjacent elements.



## First approach we might think of

- for each permutation of 1..n-1
  - for i=0..n-1
    - insert n in position i
- That is, we do the insertion of n into each smaller permutation from left to right each time
- However, to get "minimal change", we alternate:
  - Insert n L-to-R in one permutation of 1..n-1
  - Insert n R-to-L in the next permutation of 1..n-1
  - Etc.



## Example

- Bottom-up generation of permutations of 123

start	1		
insert 2 into 1 right to left	12	21	
insert 3 into 12 right to left	123	132	312
insert 3 into 21 left to right	321	231	213

- Example: Do the first few permutations for n=4



## Johnson-Trotter Approach

- integrates the insertion of n with the generation of permutations of 1..n-1
- Does it by keeping track of which direction each number is currently moving

→ ← → ←

3 2 4 1

The number k is **mobile** if its arrow points to an adjacent element that is smaller than itself

- In this example, 4 and 3 are mobile



## Johnson-Trotter Approach

→ ← → ←  
3 2 4 1

- The number k is **mobile** if its arrow points to an adjacent element that is smaller than itself.
- In this example, 4 and 3 are mobile
- To get the next permutation, exchange the largest mobile number (call it k) with its neighbor
- Then reverse directions of all numbers that are larger than k.
- Initialize: All arrows point left

Work with  
a partner  
on Q1



## Johnson-Trotter Driver

```
def main():
    p = Permutation(4)
    list = []
    next = p.next()
    while next:
        list += [next]
        next = p.next()
    print list
```



## Johnson-Trotter background code

```
left = -1 # equivalent to the left- and
right = 1 # right-pointing arrows in the book

def swap(list1, list2, i, j):
    "Swap positions i and j in both lists"
    list1[i], list1[j] = list1[j], list1[i]
    list2[i], list2[j] = list2[j], list2[i]

class Permutation:
    "Set current to the unpermuted list, and all directions pointing left"
    def __init__(self, n):
        self.current = range(1, n + 1)
        self.direction = [left] * n
        self.n = n
        self.more = True # This is not the last permutation.
```



## Johnson-Trotter major methods

```
def isMobile(self, k):
    ''' An element of a permutation is mobile if its direction "arrow"
        points to an element with a smaller value.'''
    return k + self.direction[k] in range(self.n) and \
           self.current[k + self.direction[k]] < self.current[k]

def next(self):
    "return current permutation and calculate next one"
    if not self.more:
        return False
    returnValue = [self.current[i] for i in range(self.n)]

    largestMobile = 0
    for i in range(self.n):
        if self.isMobile(i) and self.current[i] > largestMobile:
            largestMobile = self.current[i]
            largePos = i

    if largestMobile == 0:
        self.more = False # This is the last permutation
    else:
        swap(self.current, self.direction,
              largePos, largePos + self.direction[largePos])
        for i in range(self.n):
            if self.current[i] > largestMobile:
                self.direction[i] *= -1

    return ''.join([str(v) for v in returnValue])
```



## Lexicographic Permutation Generation

- Generate the permutations of 1..n in "natural" order.
- Let's do it recursively.



## Lexicographic Permutation Code

```
def permuterecursive(prefix, remaining):
    """ Generate all lists that begin with prefix and
        end with a permutation of remaining"""
    if remaining == []: # this is where the recursion ends
        return [prefix]
    result = [] # accumulate the list of generated prefixes
    for n in remaining:
        copy = [e for e in remaining] # need to remove a different
        copy.remove(n)             # number for each suffix we generate.
        result += permuterecursive(prefix + [n], copy)
    return result

def permute(n):
    return permuterecursive([], range(1, n+1))

print (permute(4))
```



## Permutations and order

number	permutation	number	permutation
0	0123	12	2013
1	0132	13	2031
2	0213	14	2103
3	0231	15	2130
4	0312	16	2301
5	0321	17	2310
6	1023	18	3012
7	1032	19	3021
8	1203	20	3102
9	1230	21	3120
10	1302	22	3201
11	1320	23	3210

- Given a permutation of 0, 1, ..., n-1, can we directly find the next permutation in the lexicographic sequence?
- Given a permutation of 0..n-1, can we determine its permutation sequence number?
- Given n and i, can we directly generate the  $i^{\text{th}}$  permutation of 0, ..., n-1?



## Discovery time (with a partner)

- Which permutation follows each of these in lexicographic order?
  - 183647520      471638520
  - Try to write an algorithm for generating the next permutation, with only the current permutation as input.
- If the lexicographic permutations of the numbers [0, 1, 2, 3, 4, 5] are numbered starting with 0, what is the number of the permutation 14032?
  - General form? How to calculate efficiency?
- In the lexicographic ordering of permutations of [0, 1, 2, 3, 4, 5], which permutation is number 541?
  - How to calculate efficiently?



## Side road: Polynomial Evaluation

- Given a polynomial  
 $p(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$
- How can we efficiently evaluate  $p(c)$  for some number  $c$ ?
- Apply this to evaluation of "31427894" or any other string that represents a positive integer.
- Write and analyze (pseudo)code

