

ALLOY OPERATORS

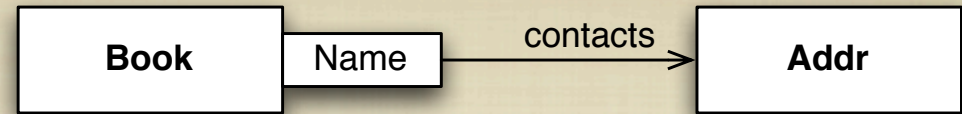
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GOALS FOR TODAY

- RECALL MORE SET THEORY IDEAS AND SEE HOW THEY'RE USED IN ALLOY
- LEARN SOME ALLOY OPERATORS THAT WE'LL USE TO DESCRIBE DESIGNS
- GET MORE COMFORTABLE EXPERIMENTING WITH ALLOY TO SEE HOW IT BEHAVES WITH AN EXAMPLE

CONSTANTS



- **none** – THE EMPTY SET
- **univ** – THE UNIVERSAL SET
 - CONTAINS ONE OF EVERY ATOM IN THE INSTANCE
- **iden** – THE IDENTITY RELATION
 - MAPS EVERY ATOM TO ITSELF
 - INCLUDING Int ATOMS -8, -7, ..., -1, 0, 1, ..., 6, 7

IN EVERY INSTANCE, UNLESS WE CHANGE “BIT WIDTH”

Q1,2

SET OPERATORS

■ SET VALUED:

■ + – UNION

■ & – INTERSECTION

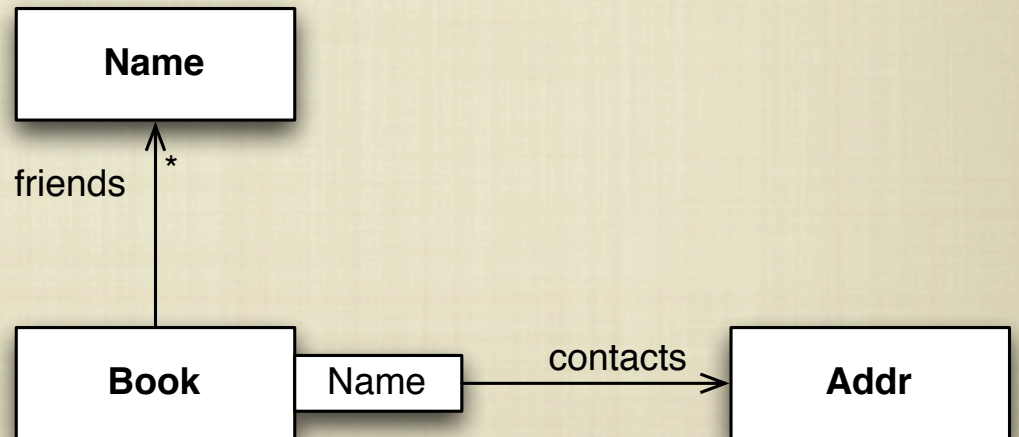
■ - – DIFFERENCE

■ BOOLEAN VALUED:

■ **in** – SUBSET

■ = – EQUALITY

ALL OF THESE
APPLY TO ANY PAIR
OF RELATIONS
WITH THE SAME
ARITY



RELATIONAL OPERATORS

■ COMBINING RELATIONS

■ \rightarrow – ARROW PRODUCT

■ \cdot – DOT JOIN

■ $[\]$ – BOX JOIN

■ REACHABILITY

■ \wedge – TRANSITIVE
CLOSURE

■ $*$ – REFLEXIVE-
TRANSITIVE CLOSURE

■ “MODIFYING” RELATIONS

■ \sim – TRANSPOSE

■ $\leftarrow :$ – DOMAIN
RESTRICTION

■ $:\rightarrow$ – RANGE
RESTRICTION

■ $++$ – OVERRIDE

COMBINING RELATIONS: ARROW PRODUCT

- $p \rightarrow q$ – EVERY POSSIBLE WAY TO CONCATENATE A TUPLE FROM P AND A TUPLE FROM Q
- EXAMPLES IN ALLOY...

COMBINING RELATIONS: DOT JOIN

- ALSO CALLED “COMPOSITION”
- WRITTEN: $p \cdot q$



DOT JOIN ON TUPLES

- $\{(N0, A0)\} \cdot \{(A0, D0)\} = \{(N0, D0)\}$
- $\{(N0, D0)\} \cdot \{(N0, D0)\} = \{\}$
- $\{(N0, D0)\} \cdot \{(D1)\} = \{\}$
- $\{(N0, D0)\} \cdot \{(D0)\} = \{(N0)\}$
- $\{(N0)\} \cdot \{(N0, D0)\} = \{(D0)\}$
- $\{(B0)\} \cdot \{(B0, N0, D0)\} = \{(N0, D0)\}$

Recall: this is
the “math” font

DOT JOINS ON RELATIONS

- SUPPOSE p AND q ARE RELATIONS
- $p.q$ IS THE RESULT OF TAKING EVERY COMBINATION OF A TUPLE FROM p AND A TUPLE FROM q AND INCLUDING THEIR JOIN

COMBINING RELATIONS: BOX JOIN

- EXACTLY THE SAME MEANING AS DOT JOIN
- SYNTACTIC SUGAR—LOOKS LIKE ARRAY ACCESS
- WRITTEN: $q[p]$, MEANS $p.q$
 - TYPICALLY p WILL BE A SCALAR

“MODIFYING” RELATIONS: TRANSPOSE

- $\sim p$ – TRANSPOSE
 - YIELDS A NEW RELATION BY FLIPPING ORDER OF TUPLES IN p
 - p MUST BE BINARY
 - USEFUL FOR GETTING ORDER RIGHT BEFORE A JOIN

EXAMPLE: OF WHOM IS NAME\$2 A FRIEND?

“MODIFYING” RELATIONS: RESTRICT DOMAIN/RANGE

■ $S <: \mathcal{P}$, DOMAIN RESTRICTION

■ $\mathcal{P} :> S$, RANGE RESTRICTION

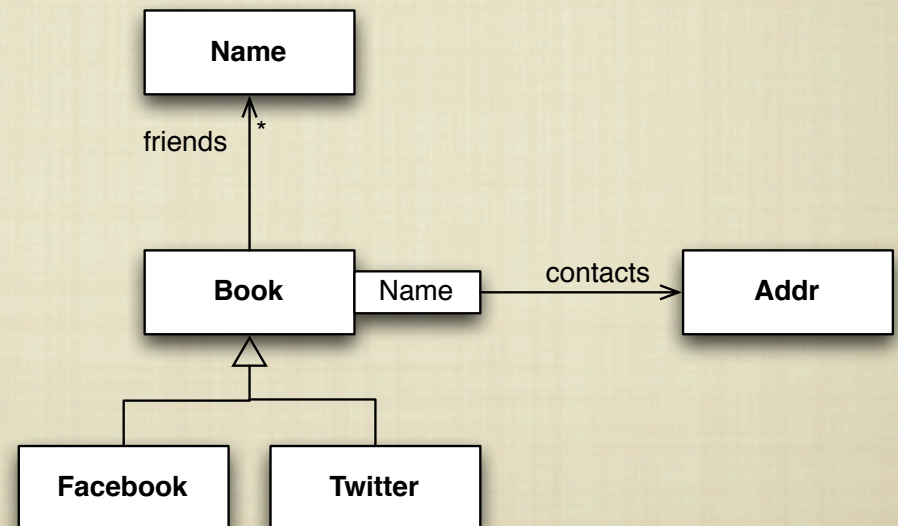
■ YIELDS A NEW RELATION BY THROWING OUT TUPLES IN \mathcal{P} THAT DON'T START WITH AN ELEMENT OF S

■ LIKE $<$: BUT MATCHES LAST ELEMENT OF \mathcal{P} 'S TUPLES

■ S MUST BE A SET

■ USEFUL FOR FOCUSING ON “SUBCLASSES”

■ UNLIKE DOT JOIN, IT KEEPS ALL COLUMNS



NEXT TIME

- **TECHNIQUES FOR TELLING ALLOY WHAT MUST BE TRUE ABOUT A DESIGN:**
 - **OVERRIDE OPERATOR**
 - **REACHABILITY OPERATORS**
 - **CONSTRAINTS**

ALLOY MODELS FOR
REFERENCE

MODEL FROM CLASS

```
abstract sig Book {  
  contacts: Name -> lone Addr,  
  friends: set Name,  
  owner: lone Name  
}  
sig Name, Addr {}
```

```
sig Twitter, Facebook extends Book {}
```

```
pred show[b1, b2:Book] {  
  b1 != b2  
  #b1.contacts > 1  
  #b2.contacts > 1  
}
```

```
run show for 3 but 2 Book
```

```
pred addContact[b, b': Book, n: Name, a: Addr] {  
  b != b'  
  some Facebook & b iff some Facebook & b'  
  b'.contacts = b.contacts ++ n->a  
}
```

```
pred addContacts[b, b': Book, c: Name -> lone Addr] {  
  b != b'  
  some Facebook & b iff some Facebook & b'  
  b'.contacts = b.contacts ++ c  
}
```

```
run addContacts for 4 but 2 Book
```