

ALLOY CONSTRAINTS

CURT CLIFTON

ROSE-HULMAN INSTITUTE OF TECHNOLOGY

HOMWORK TEAMS

TEAM **01**: MCGINNDA,WELLSKA1

TEAM **02**: THEISJE,WATTSBN

TEAM **03**: GAOT,SIEGLEAL,FREEMACC

TEAM **04**: GLOWSKST,DASHJB,SPIEGEKR

TEAM **05**: KEHMBV,MANNDJ,FULLERRA

TEAM **06**: STAMPTD,SHERMABJ,HOLLINTL

TEAM **07**: MOSTTW,ANDERSWC,ORLOWSAP

TEAM **08**: MENDELNT,COVERTCJ,MLYNARCS

TEAM **09**: BANKSDA,HINESEN,MANKEAP

TEAM **10**: PURVIATT,FRANKMP,LINT

TEAM **11**: HAFFNEDM,SCHEREPN,JONESJG

<http://svn.csse.rose-hulman.edu/repos/csse373-201130-hwteamXX>

PROJECT TEAMS

- PROJECT 1 IS POSTED
 - DUE TUESDAY, MARCH 29, 11:55PM
 - TEAMS OF 2-3 PEOPLE
 - CAN FORM YOUR OWN TEAMS
 - PROJECT TEAMS DISJOINT FROM HOMEWORK ONES
- COMPLETE TEAM SURVEY ON ANGEL BY 6AM, MONDAY, MARCH 21.

GOALS FOR TODAY

- LEARN SOME MORE ALLOY OPERATORS THAT WE'LL USE TO DESCRIBE DESIGNS
- LEARN HOW TO WRITE BASIC CONSTRAINTS ABOUT WHAT “MUST BE” TRUE ABOUT A DESIGN

RELATIONAL OPERATORS

■ COMBINING RELATIONS

■ \rightarrow – ARROW PRODUCT

■ \cdot – DOT JOIN

■ \square – BOX JOIN

■ REACHABILITY

■ \wedge – TRANSITIVE
CLOSURE

■ $*$ – REFLEXIVE-
TRANSITIVE CLOSURE

■ “MODIFYING” RELATIONS

■ \sim – TRANSPOSE

■ $\leftarrow :$ – DOMAIN
RESTRICTION

■ $:\rightarrow$ – RANGE
RESTRICTION

■ $++$ – OVERRIDE

“MODIFYING” RELATIONS: OVERRIDE

■ $p \mathbin{++} q$ – OVERRIDE

■ LIKE UNION, BUT TUPLES IN Q REPLACE
“MATCHING” TUPLES FROM P

■ “MATCHING” MEANS FIRST ELEMENTS ARE THE
SAME

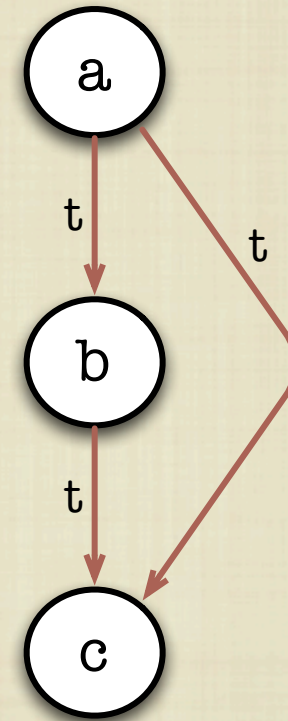
■ YIELDS A NEW RELATION

■ p AND q MUST HAVE SAME ARITY

■ USEFUL FOR MODELING MUTATION

TRANSITIVE RELATION

- **TRANSITIVE RELATION:**
- **BINARY RELATION, t , SUCH THAT IF $a \rightarrow b$ AND $b \rightarrow c$ ARE IN t , THEN $a \rightarrow c$ ALSO IS**



TRANSITIVE CLOSURE

- **TRANSITIVE CLOSURE OF BINARY RELATION r :**
 - **SMALLEST TRANSITIVE RELATION CONTAINING r**
- **IN ALLOY: \hat{r}**
- **EQUIVALENT TO $r + r.r + r.r.r + \dots$**

REACHABILITY

- TRANSITIVE CLOSURE IS USED TO EXPRESS “REACHABILITY”

- `bacon.^appearedWith`

- RELATED:

- REFLEXIVE TRANSITIVE CLOSURE, `*r`

- `*r = ^r + iden`

KEVIN BACON IS ZERO DEGREES FROM KEVIN BACON

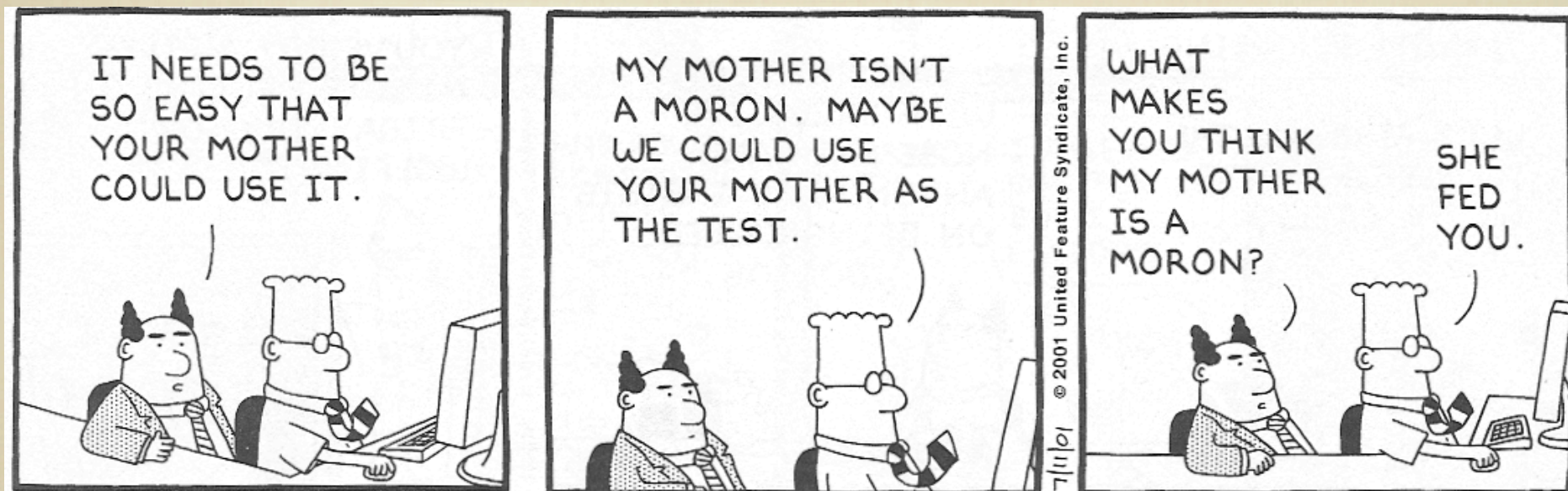


```
sig Actor {  
  appearedWith: set Actor  
}
```

```
pred degrees[bacon: Actor] {  
  all a: Actor |  
    a in bacon.^appearedWith  
}
```

```
run degrees for 6
```

CARTOON OF THE DAY



CONSTRAINTS

- LIKE USUAL BOOLEAN OPERATORS
- `else` GOES WITH `implies`:
 - `C1 => F1 else F2`
 - `C1 => F1`
`else C2 => F2`
`else C3 => F3`

| LONG FORM | SHORTHAND |
|----------------|-----------|
| not | ! |
| and | && |
| or | |
| implies | => |
| else | |
| iff | <=> |

QUANTIFICATION

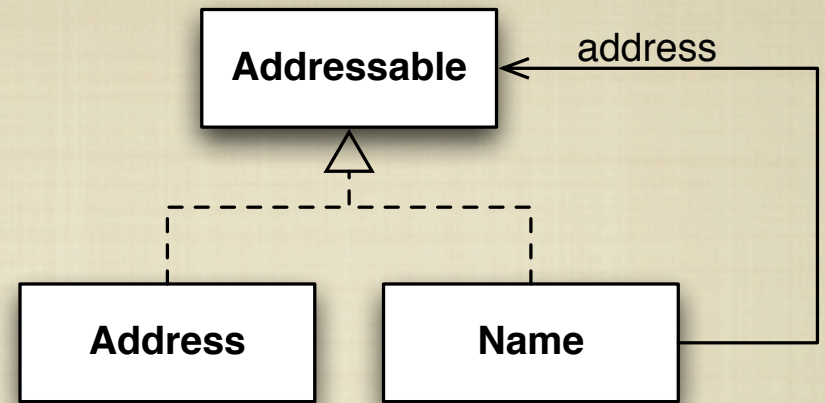
CAN DECLARE
MULTIPLE VARS HERE

| | |
|---------------------------|--------------------------------------|
| all $x: e \mid F$ | F HOLDS FOR EVERY x IN e |
| some $x: e \mid F$ | F HOLDS FOR SOME x IN e |
| no $x: e \mid F$ | F HOLDS FOR NO x IN e |
| lone $x: e \mid F$ | F HOLDS FOR AT MOST ONE x IN e |
| one $x: e \mid F$ | F HOLDS FOR EXACTLY ONE x IN e |

THINK LESS THAN
OR EQUAL TO ONE

EXAMPLES

- **some** n: Name, a: Address |
a **in** n.address
- **no** n: Name |
n **in** n.^address
- **all** n: Name |
lone d: Address |
d **in** n.address
- **all** n: Name |
no disj d, d': Address |
d + d' **in** n.address



```
sig Address {}
sig Name {
    // multi-level address book
    address: set (Name + Address)
}
```

“DISJOINT”

QUANTIFIED EXPRESSIONS

- **some** Name
- **some** address
- **no** (address.Addr - Name)
- **all** n: Name | **lone** n.address

```
sig Address {}  
sig Name {  
  // multi-level address book  
  address: set (Name + Address)  
}
```

LET EXPRESSIONS

- **let** $x = e \mid A$
- JUST A SHORTHAND TO AVOID WRITING OUT E MULTIPLE TIMES
- **all** $a: \text{Alias} \mid$
 let $w = a.\text{workAddress} \mid$
 $a.\text{address} = (\text{some } w \Rightarrow w \text{ else } a.\text{homeAddress})$

VARIABLE AND FORMAL DECLARATIONS

- name: expression
 - name IS A **SUBSET** OF THE RELATION GIVEN BY expression
- **EXAMPLES:**
 - address: Name->Addr
 - addr: Book->Name->Addr
 - address: Name->(Name + Addr)
 - workAddress, homeAddress: Alias->Addr
prefAddress: workAddress + homeAddress

SET MULTIPLICITIES

- USED TO CONSTRAIN THE POSSIBLE SUBSETS THAT A VARIABLE CAN BE
- $x: \text{set } e - x$ CAN BE ANY SUBSET OF e
- $x: \text{one } e - x$ IS A SINGLETON SUBSET OF e (I.E., AN ALLOY SCALAR)
- $x: \text{lone } e - x$ IS AN OPTION, EITHER EMPTY SET OR A SCALAR
- $x: \text{some } e - x$ IS A NON-EMPTY SUBSET OF e

CAREFUL: IF e IS A UNARY RELATION (I.E., A SET), THEN $x: e$ IS EQUIVALENT TO $x: \text{one } e$

RELATION MULTIPLICITIES

■ TOO BIZARRE FOR WORDS

■ ALMOST

■ $r: A \ m \rightarrow n \ B$ MEANS:

EXAMPLES:

$r: A \rightarrow$ **one** B

$r: A$ **one** $\rightarrow B$

$r: A \rightarrow$ **lone** B

$r: A$ **one** \rightarrow **one** B

$r: A$ **some** \rightarrow **some** B

■ **EACH** MEMBER OF A MAPS TO **n** MEMBERS OF B

■ **AND FOR EACH** MEMBER OF B , **m** MEMBERS OF A
MAP TO IT

```
sig Thing, OtherThing {}
```

```
pred relMult[r: Thing some -> some OtherThing] {}
```

```
run relMult for 3
```

CARDINALITY CONSTRAINTS

- #e GIVES THE SIZE (NUMBER OF TUPLES) IN THE RELATION GIVEN BY e

- CAN USE ALL REGULAR INTEGER OPERATIONS ON THE RESULT

- CAN USE 1, 2, 3, ... AS CONSTANTS

- **sum** x: e | ie MEANS $\sum_{x \in e} ie$



NEXT TIME

- MORE EXAMPLES
- BUILDING OUR OWN ALLOY MODELS