

# ALLOY CONSTRAINTS

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# RELATIONAL OPERATORS

## ■ COMBINING RELATIONS

■  $\rightarrow$  – ARROW PRODUCT

■  $\cdot$  – DOT JOIN

■  $\square$  – BOX JOIN

## ■ REACHABILITY

■  $\wedge$  – TRANSITIVE  
CLOSURE

■  $*$  – REFLEXIVE-  
TRANSITIVE CLOSURE

## ■ “MODIFYING” RELATIONS

■  $\sim$  – TRANSPOSE

■  $\leftarrow :$  – DOMAIN  
RESTRICTION

■  $:\rightarrow$  – RANGE  
RESTRICTION

■  $++$  – OVERRIDE

# “MODIFYING” RELATIONS: RESTRICT DOMAIN/RANGE

- $s \prec: r$  – DOMAIN RESTRICTION
  - YIELDS A NEW RELATION BY THROWING OUT TUPLES IN  $r$  THAT DON'T START WITH AN ELEMENT OF  $s$
  - $s$  MUST BE A SET
  - USEFUL FOR FOCUSING ON “SUBCLASSES”
- $r \succ: s$  – RANGE RESTRICTION, LIKE DOMAIN RESTRICTION BUT MATCHES LAST ELEMENT OF  $r$ 'S TUPLES

# “MODIFYING” RELATIONS: OVERRIDE

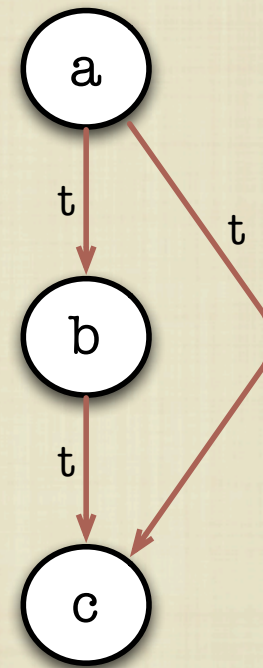
- $p \text{ ++ } q$  – OVERRIDE
  - LIKE UNION, BUT TUPLES IN Q REPLACE “MATCHING” TUPLES FROM P
  - “MATCHING” MEANS FIRST ELEMENTS ARE THE SAME
  - YIELDS A NEW RELATION
  - $p$  AND  $q$  MUST HAVE SAME ARITY
  - USEFUL FOR MODELING MUTATION



TRAN  
**TRANSVESTITE CLOTHIER**  
CLOSURE

# TRANSITIVE RELATION

- **TRANSITIVE RELATION:**
  - **BINARY RELATION,  $t$ ,**  
SUCH THAT IF  $a \rightarrow b$  AND  
 $b \rightarrow c$  ARE IN  $t$ , THEN  $a \rightarrow c$   
ALSO IS



# TRANSITIVE CLOSURE

- **TRANSITIVE CLOSURE OF BINARY RELATION  $r$ :**
  - **SMALLEST TRANSITIVE RELATION CONTAINING  $r$**
- **IN ALLOY:  $\hat{r}$**
- **EQUIVALENT TO  $r + r.r + r.r.r + \dots$**

# REACHABILITY

- TRANSITIVE CLOSURE IS USED TO EXPRESS “REACHABILITY”

- `bacon.^appearedWith`

- RELATED:

- REFLEXIVE TRANSITIVE CLOSURE,  $*r$

- $*r = ^r + \text{iden}$

KEVIN BACON IS ZERO DEGREES FROM KEVIN BACON

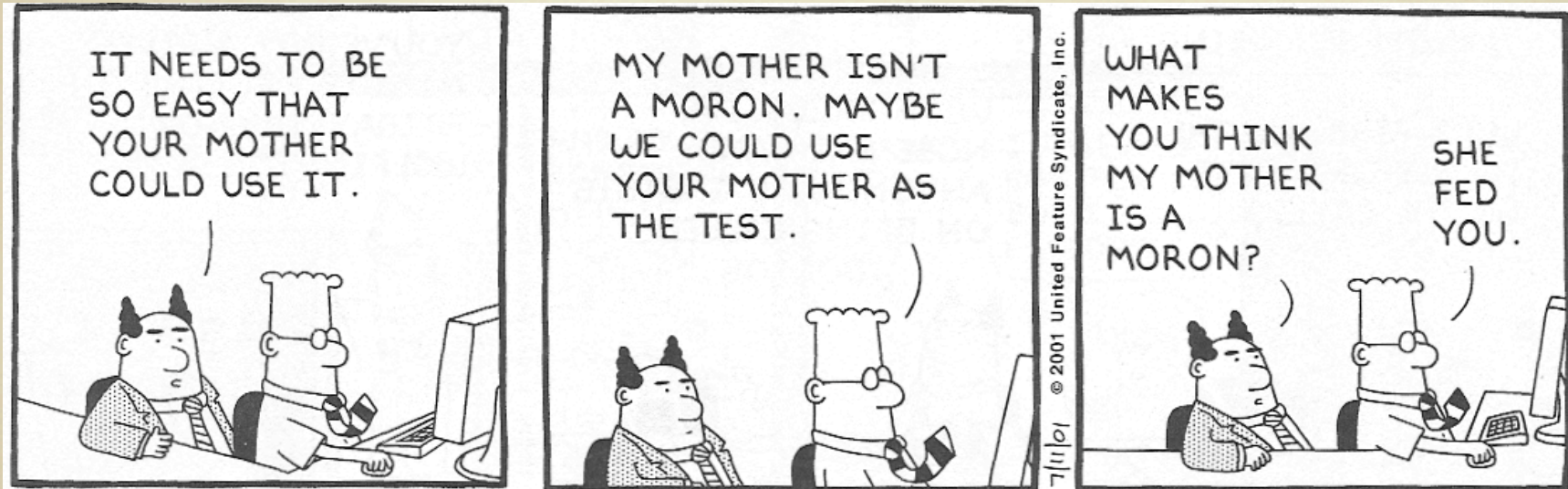


```
sig Actor {  
  appearedWith: set Actor  
}
```

```
pred degrees[bacon: Actor] {  
  all a: Actor |  
    a in bacon.^appearedWith  
}
```

```
run degrees for 6
```

# CARTOON OF THE DAY



# CONSTRAINTS

- LIKE USUAL BOOLEAN OPERATORS
- else GOES WITH implies:
  - $C1 \Rightarrow F1$  **else**  $F2$
  - $C1 \Rightarrow F1$   
**else**  $C2 \Rightarrow F2$   
**else**  $C3 \Rightarrow F3$

LONG FORM	SHORTHAND
<b>not</b>	!
<b>and</b>	&&
<b>or</b>	
<b>implies</b>	=>
<b>else</b>	
<b>iff</b>	<=>

# QUANTIFICATION

CAN DECLARE  
MULTIPLE VARS HERE

<b>all</b> $x: e \mid F$	$F$ HOLDS FOR EVERY $x$ IN $e$
<b>some</b> $x: e \mid F$	$F$ HOLDS FOR SOME $x$ IN $e$
<b>no</b> $x: e \mid F$	$F$ HOLDS FOR NO $x$ IN $e$
<b>lone</b> $x: e \mid F$	$F$ HOLDS FOR AT MOST ONE $x$ IN $e$
<b>one</b> $x: e \mid F$	$F$ HOLDS FOR EXACTLY ONE $x$ IN $e$

THINK LESS THAN  
OR EQUAL TO **ONE**

# EXAMPLES

■ **some** n: Name, a: Address |  
a **in** n.address

■ **no** n: Name |  
n **in** n.^address

■ **all** n: Name |  
**lone** d: Address |  
d **in** n.address

■ **all** n: Name |  
**no disj** d, d': Address |  
d + d' **in** n.address

“DISJOINT”

```
sig Address {}  
sig Name {  
  // multi-level address book  
  address: set (Name + Address)  
}
```

# QUANTIFIED EXPRESSIONS

- **some** Name
- **some** address
- **no** (address.Addr - Name)
- **all** n: Name | **lone** n.address

```
sig Address {}  
sig Name {  
  // multi-level address book  
  address: set (Name + Address)  
}
```

# LET EXPRESSIONS

- **let** x = e | A
- JUST A SHORTHAND TO AVOID WRITING OUT E MULTIPLE TIMES
- **all** a: Alias |  
    **let** w = a.workAddress |  
    a.address = (**some** w => w **else** a.homeAddress)

# VARIABLE AND FORMAL DECLARATIONS

- name: expression
  - name IS A **SUBSET** OF THE RELATION GIVEN BY expression
- **EXAMPLES:**
  - address: Name->Addr
  - addr: Book->Name->Addr
  - address: Name->(Name + Addr)
  - workAddress, homeAddress: Alias->Addr  
prefAddress: workAddress + homeAddress

# SET MULTIPLICITIES

- USED TO CONSTRAIN THE POSSIBLE SUBSETS THAT A VARIABLE CAN BE
- $x: \mathbf{set} e - x$  CAN BE ANY SUBSET OF  $e$
- $x: \mathbf{one} e - x$  IS A SINGLETON SUBSET OF  $e$  (I.E., AN ALLOY SCALAR)
- $x: \mathbf{lone} e - x$  IS AN OPTION, EITHER EMPTY SET OR A SCALAR
- $x: \mathbf{some} e - x$  IS A NON-EMPTY SUBSET OF  $e$

CAREFUL: IF  $e$  IS A UNARY RELATION (I.E., A SET),  
THEN  $x: e$  IS EQUIVALENT TO  $x: \mathbf{one} e$

# RELATION MULTIPLICITIES

- TOO BIZARRE FOR WORDS

- ALMOST

- $r: A \ m \rightarrow n \ B$  MEANS:

## EXAMPLES:

$r: A \rightarrow$  **one**  $B$

$r: A$  **one**  $\rightarrow B$

$r: A \rightarrow$  **lone**  $B$

$r: A$  **one**  $\rightarrow$  **one**  $B$

$r: A$  **some**  $\rightarrow$  **some**  $B$

- **EACH** MEMBER OF  $A$  MAPS TO  $n$  MEMBERS OF  $B$

- AND FOR **EACH** MEMBER OF  $B$ ,  $m$  MEMBERS OF  $A$  MAP TO IT

```
sig Thing, OtherThing {}
```

```
pred relMult[r: Thing some  $\rightarrow$  some OtherThing] {}
```

```
run relMult for 3
```

# CARDINALITY CONSTRAINTS

- #e GIVES THE SIZE (NUMBER OF TUPLES) IN THE RELATION GIVEN BY e

- CAN USE ALL REGULAR INTEGER OPERATIONS ON THE RESULT

- CAN USE 1, 2, 3, ... AS CONSTANTS

- **sum** x: e | ie MEANS  $\sum_{x \in e} ie$



# NEXT TIME

- MORE EXAMPLES
- BUILDING OUR OWN ALLOY MODELS