# Query Processing and Optimization

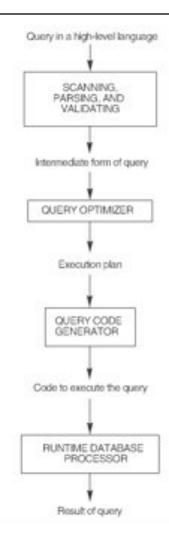
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#### Outline

- □ Basic Optimization Approach
- □ Algorithms for Processing Queries
- Pipelining
- □ Techniques for Automatic Query Optimization

## Introduction to Query Processing

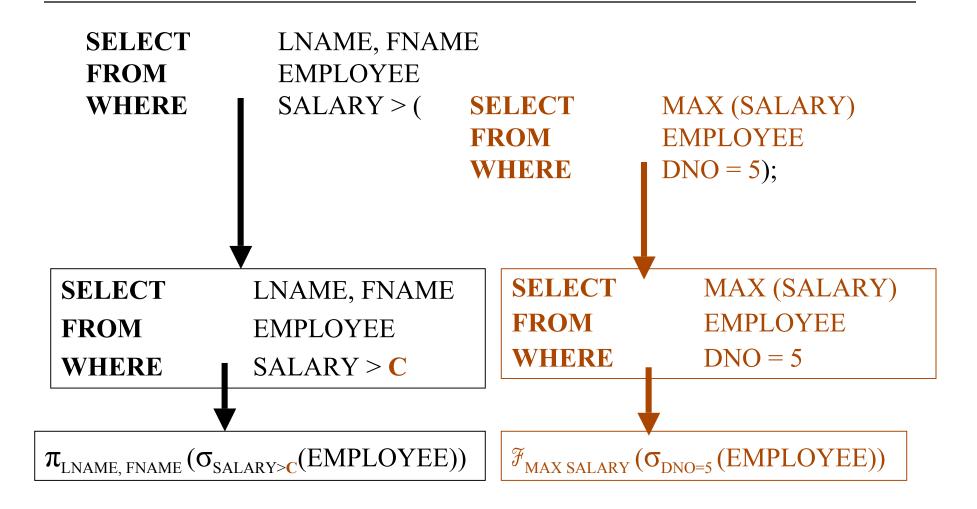
- What is query optimization?
- □ Typically intermediate form is a *query tree*



## From SQL to Relational Algebra

- □ Query block: the basic unit that can be translated into the algebraic operators and optimized
- Nested queries become separate query blocks
- □ Aggregate operators in SQL require extended algebra
- □ Example...

#### **Example Translation**



#### What Next?

- Queries reduced to query trees in relational algebra
- □ DBMS considers various algorithms for processing query
- □ Rewrites tree to use "best" algorithms
- □ Variety of algorithms exist to solve various query problems

## Problem: Sorting Huge Datasets

- □ Use external sorting
- □ Phase 1:
  - Load *n* pages into memory, as many as fit (a "run")
  - Sort them and save back to disk
  - Repeat until all runs are sorted
- □ Phase 2:
  - Perform an (n-1)-way merge
    - □ One page for "top" of each of n-1 runs
    - □ One page for "bottom" of merge results
  - Repeat until done

#### Problem: Selecting Subset of Rows

- □ Linear search:
  - Last resort, unless file is small
- □ Binary search:
  - For ordered data without an index
- □ Using an index for equality comparisons:
  - Just look up the record

#### Problem: Selecting Subset of Rows

- □ Using a primary index for order comparisons:
  - Find edge of range using index
  - Scan from there
- □ Using a secondary index for order comparisons:
  - Find edge of range using index
  - Scan leaf nodes of index from there, loading data based on pointers

#### Select With Complex Condition

- □ Simple conjunctive selection:
  - Pick one condition for which some previous method would work
  - Use brute force to filter those results based on other conditions
- □ Conjunctive selection with a composite index:
  - Works if index covers all attributes in the complex condition

#### Select With Complex Condition

- □ Conjunctive selection by intersection of record pointers:
  - Suppose:
    - □ Secondary indexes are several fields in condition
    - Indexes include record pointers
  - Then:
    - Use indexes to get sets of the record pointers for conjuncts
    - □ Take intersection of pointer sets
    - □ Then retrieve actual records

#### Problem: Joining Two Tables

- □ Nested-loop join (brute force):
  - Last resort unless tables are small
- □ Single-loop join when one table has index
  - Loop over one table
  - Use index to find matches in other table

#### Problem: Joining Two Tables

- □ Sort-merge join when both tables sorted by join attributes
  - Scan both files matching the records that have the same values for join attributes

## Problem: Combining Multiple Ops.

- Generating and saving temporary files is time expensive
- □ So, avoid constructing temporary results
- □ Pipeline the data through multiple operations:
  - Pass the result of a previous operator to the next
  - Page-by-page instead of operation-by-operation
- □ Example...

#### Pipelining Example

- □ SELECT (FName + ' ' + LName) AS Name FROM Employee e JOIN Department d ON e.DNo = d.DNumber WHERE e.Salary < 50000 AND d.Location <> 'Houston'
- □ What are the individual operations for this?
- □ How many ways could this be pipelined?

## Picking Algorithms and Plans

- Heuristics
- □ Cost estimation

#### Using Heuristics

- □ Uses pattern matching to transform parts of query tree to a "best" shape
- □ Patterns based on transformations that are likely to be more efficient:
  - E.g., Apply selection before applying join
  - Why is that likely (naively) to be more efficient?

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#### Cost-based Optimization

- □ Estimate the costs of a variety of different versions of the query based on:
  - Available indexes
  - Specificity of conditions
  - Statistics on data
  - Disk speed
  - Memory available
  - Block and record sizes
  - Index blocking factors

## Issues in Cost-based Optimization

- □ Accuracy of statistics
- □ Cost of calculating costs
- Accuracy of estimates of disk speed, memory available
- □ Shear number of possible execution strategies

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- □ SQL Server uses cost-based optimization
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- □ Cost-based optimization is "taking over"
- □ SQL Server uses cost-based optimization
- Does NOT try to minimize total cost!
- □ Tries to minimize time to initial results