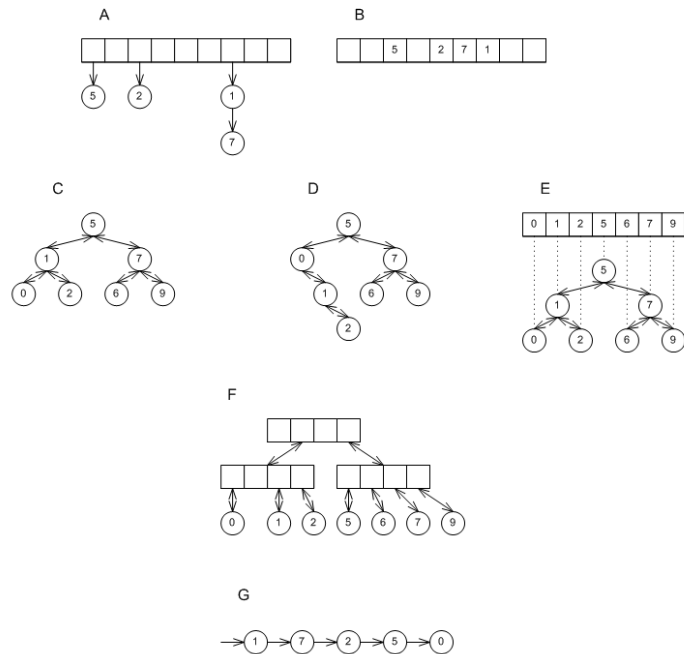


# CSSE 230 Day 5

Abstract Data Types  
Data Structure “Grand Tour”  
Java Collections



# Announcements

- ▶ Stacks and Queues

- Ideally, you have met with your partner to start
- Try your best to work well together, even if you have different amounts of programming experience.

Suggestion: Let the weaker programmer do most of the driving

- ▶ Finish day 4 + quiz with instructor if needed.

- ▶ Exam 1: next Tuesday, 7–9pm.

# How is Homework 2 coming?

- ▶ From question 3:

Suppose  $T_1(N)$  is  $O(f(N))$  and  $T_2(N)$  is  $O(f(N))$ . Prove that  $T_1(N) + T_2(N)$  is  $O(f(N))$  or give a counter-example.

- Hint: Supposing  $T_1(N)$  and  $T_2(N)$  are  $O(f(N))$ , that means there exist constants  $c_1, c_2, n_1, n_2$ , such that.....
  - How can you use these constants?
- ▶ What about the similar question for  $T_1(N) - T_2(N)$ ?
    - Remember,  $O$  isn't a tight bound.

# After today, you should be able to...

- ▶ explain what an Abstract Data Type (ADT) is
- ▶ List examples of ADTs in the Collections framework (from HW2 #1)
- ▶ List examples of data structures that implement the ADTs in the Collections framework
- ▶ Choose an ADT and data structure to solve a problem

# ADTs and Data Structures

# A *data type* is an interpretation of data (bits)



- “What is this data, and how does it work?”
- Primitive types (**int**, **double**): hardware-based
- Objects (such as **java.math.BigInteger**): require software interpretation
- Composite types (**int[]**): software + hardware

# What is an Abstract Data Type (ADT)?

- ▶ A mathematical model of a data type
- ▶ Specifies:
  - The type of data stored (but not *how* it's stored)
  - The operations supported
  - Argument types and return types of these operations (but not *how* they are implemented)

# An Example ADT: Stack

- ▶ Three basic operations:
  - **isEmpty**
  - **push**
  - **pop**
- ▶ Derived operations include **peek** (a.k.a. **top**)
  - How could we write it in terms of the basic operations?
  - We could have **peek** be a basic operation instead.
  - Advantages of each approach?
- ▶ Possible implementations:
  - Use a linked list.
  - Use a growable array.
  - Last time, we talked about implementation details for each.



# ADTs for collections of items

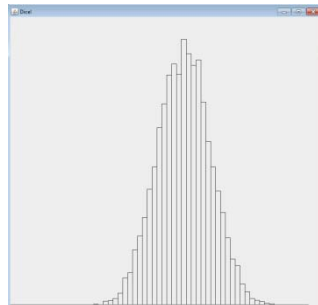
Application:  
“how can you use it?”

Specification  
“what can it do?”

Implementation:  
“How is it built?”

```
public static void main(String[] args) {
    Scanner scanner = new Scanner(System.in);
    ArrayList<SingleDie> dice = new ArrayList<SingleDie>();
    while (true) {
        System.out.printf("How many sides (Q to quit):");
        String response = scanner.next();
        if (Character.toUpperCase(response.charAt(0)) == 'Q') {
            break;
        }
        int nSides = Integer.parseInt(response);
        nSides = (nSides < 4) ? 4: nSides;
        dice.add(new SingleDie(nSides));
    }

    scanner.close();
    int minSum = dice.size();
    int maxSum = 0;
    for (SingleDie die : dice) {
        maxSum += die.getNSides();
    }
}
```



Modifier and Type	Method and Description
boolean	<b>add(E e)</b> Appends the specified element to the collection.
void	<b>add(int index, E element)</b> Inserts the specified element at the specified position in the list.
boolean	<b>addAll(Collection c)</b> Appends all of the elements in the specified collection to the end of this list, in the order that they are returned by the collection's iterator.
boolean	<b>addAll(int index, Collection c)</b> Inserts all of the elements in the specified collection into this list at the position specified by the index.
void	<b>clear()</b> Removes all of the elements from the list.
boolean	<b>contains(Object o)</b> Returns true if the list contains the specified element.
boolean	<b>containsAll(Collection c)</b> Returns true if this list contains all of the elements in the specified collection.
boolean	<b>equals(Object o)</b> Compares the specified object to this list for equality.
E	<b>get(int index)</b> Returns the element at the specified position in this list.

```
public class ArrayList<E> extends AbstractList<E>
    implements List<E>, RandomAccess, Cloneable {
    private static final long serialVersionUID = 8
        645561451L;

    private transient Object[] elementData;

    private int size;

    public ArrayList(int initialCapacity) {
        super();
        if (initialCapacity < 0)
            throw new IllegalArgumentException("Illegal
            capacity: " + initialCapacity);
        this.elementData = new Object[initialCapacity];
    }

    public ArrayList() {
        this(10);
    }
}
```

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# Common collection ADTs and implementations (data structures!)

- ▶ List
  - Array List
  - Linked List
- ▶ Stack
- ▶ Queue
- ▶ Set
  - Tree Set
  - Hash Set
  - Linked Hash Set

- ▶ Map
  - Tree Map
  - Hash Map
- ▶ Priority Queue

Underlying data structures for many

- Array
- Tree

Implementations for almost all of these\* are provided by the **Java Collections Framework** in the **java.util** package.

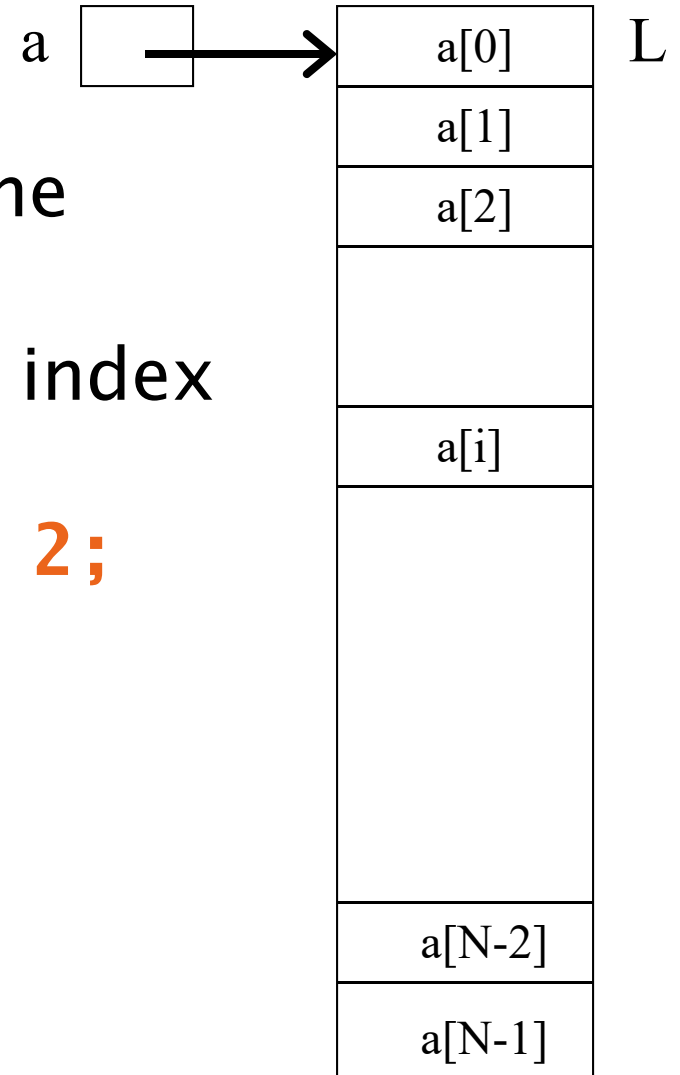
# Java Collections Framework

Reminder: Available, efficient, bug-free implementations of many key data structures

Most classes are in `java.util`

You started this in HW2 #1; Weiss Chapter 6 has more details

# Array



- ▶ Size must be declared when the array is constructed
- ▶ Can look up or store items by index

Example:

```
nums[i+1] = nums[i] + 2;
```

- ▶ How is this done?

# List

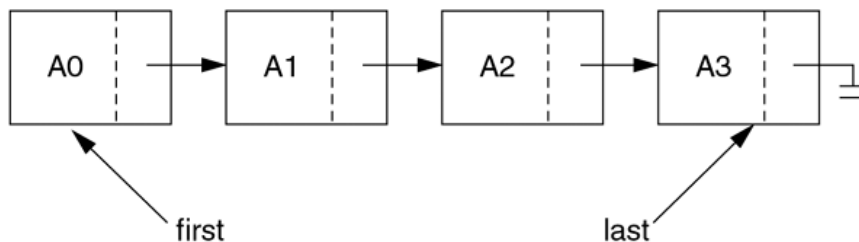
- ▶ A list is an indexed collection where elements may be added anywhere, and any elements may be deleted or replaced.
- ▶ Accessed by **index**
- ▶ **Implementations:**
  - ArrayList
  - LinkedList

# Array Lists and Linked Lists

Operations Provided	ArrayList Efficiency	LinkedList Efficiency
Random access	$O(1)$	$O(n)$
Add/remove at end	amortized $O(1)$ , worst $O(n)$	$O(1)$
Add/remove at iterator location	$O(n)$	$O(1)$



ArrayList



**figure 6.19**

A simple linked list

# Stack

- ▶ A last-in, first-out (LIFO) data structure

- ▶ Real-world stacks

- Plate dispensers in the cafeteria
- Pancakes!

- ▶ Some uses:

- Tracking paths through a maze
- Providing “unlimited undo” in an application

- ▶ `java.util.Stack` uses `LinkedList` implementation

```
public static void printInReverse(List<String> words) {  
    // TODO: implement  
    Stack<String> stack = new Stack<String>();  
    for (String w : words) {  
        stack.push(w);  
    }  
    while (!stack.isEmpty()) {  
        System.out.println(stack.pop());  
    }  
}
```

Operations Provided	Efficiency
Push item	$O(1)$
Pop item	$O(1)$

Implemented by **Stack**, **LinkedList**, and **ArrayDeque** in Java

# Queue

- ▶ first-in, first-out (FIFO)

data structure

- ▶ Real-world queues

- Waiting line at the BMW
- Character on Star Trek TNG

- ▶ Some uses:

- Scheduling access to shared resource (e.g., printer)

```
/**
 * Uses a queue to print pairs of words consisting of
 * a word in the input and the word that appeared five
 * words before it.
 *
 * @param words
 */
public static void printCurrentAndPreceding(List<String> words) {
    // TODO: implement
    ArrayDeque<String> queue = new ArrayDeque<String>();
    // Preloads the queue:
    for (int i = 0; i < 5; i++) {
        queue.add("NotAWord");
    }
    for (String w : words) {
        queue.add(w);
        String fiveAgo = queue.remove();
        System.out.println(w + ", " + fiveAgo);
    }
}
```

Operations Provided	Efficiency
Enqueue item	O(1)
Dequeue item	O(1)

Implemented by  
**LinkedList** and  
**ArrayDeque** in  
Java



# Set

- ▶ A collection of items without duplicates (in general, order does not matter)
  - If **a** and **b** are both in set, then **!a.equals(b)**
- ▶ Real-world sets:
  - Students
  - Collectibles
- ▶ One possible use:
  - Quickly checking if an item is in a collection

```
public static void printSortedWords(List<String> words) {  
    TreeSet<String> ts = new TreeSet<String>();  
    for (String w : words) {  
        ts.add(w);  
    }  
    for (String s : ts) {  
        System.out.println(s);  
    }  
}
```

Example from 220

Operations	HashSet	TreeSet
Add/remove item	amort. $O(1)$ , worst $O(n)$	$O(\log n)$
Contains?	$O(1)$	$O(\log n)$

Sorts items!

# Map

How is a TreeMap like a TreeSet?  
How is it different?

- ▶ Associate **keys** with **values**
- ▶ Real-world “maps”
  - Dictionary
  - Phone book
- ▶ Some uses:
  - Associating student ID with transcript
  - Associating name with high scores

Operations	HashMap	TreeMap
Insert key-value pair	amort. $O(1)$ , worst $O(n)$	$O(\log n)$
Look up the value associated with a given key	$O(1)$	$O(\log n)$

Sorts items by key!

# HashMap/HashSet Example (220)

```
public static void printWordCountsByLength(List<String> words) {
    HashMap<Integer, HashSet<String>> map =
        new HashMap<Integer, HashSet<String>>();

    for (String w : words) {
        int len = w.length();
        HashSet<String> set;
        if (map.containsKey(len)) {
            set = map.get(len);
        } else {
            set = new HashSet<String>();
            map.put(len, set);
        }
        set.add(w);
    }
    System.out.printf("%d unique words of length 3.%n", getCount(map, 3));
    System.out.printf("%d unique words of length 7.%n", getCount(map, 7));
    System.out.printf("%d unique words of length 9.%n", getCount(map, 9));
    System.out.printf("%d unique words of length 15.%n", getCount(map, 15));
}
```

```
public static int getCount(HashMap<Integer, HashSet<String>> map, int key) {
    if (map.containsKey(key)) {
        return map.get(key).size();
    } else {
        return 0;
    }
}
```

# Priority Queue

Not like regular queues!

- ▶ Each **item** stored **has an** associated **priority**
  - Only item with “minimum” priority is accessible
  - Operations: **insert**, **findMin**, **deleteMin**
- ▶ Real-world “priority queue”:
  - Airport ticketing counter
- ▶ Some uses
  - Simulations
  - Scheduling in an OS
  - Huffman coding

```
PriorityQueue<String> stringQueue =  
    new PriorityQueue<String>();  
  
stringQueue.add("ab");  
stringQueue.add("abcd");  
stringQueue.add("abc");  
stringQueue.add("a");  
  
while(stringQueue.size() > 0)  
    System.out.println(stringQueue.remove());
```

Assumes a binary heap implementation. The version in Warm Up and Stretching isn't this efficient.

Operations Provided	Efficiency
Insert/ Delete Min	amort. $O(\log n)$ , worst $O(n)$
Find Min	$O(1)$

# Trees, Not Just For Sorting

- ▶ Collection of nodes
  - One specialized node is the root.
  - A node has one parent (unless it is the root)
  - A node has zero or more children.
- ▶ Real-world “trees”:
  - Organizational hierarchies
  - Some family trees
- ▶ Some uses:
  - Directory structure on a hard drive
  - Sorted collections

Only if tree is  
“balanced”

Operations Provided	Efficiency
Find	$O(\log n)$
Add/remove	$O(\log n)$

# Graphs

- ▶ A collection of nodes and edges
  - Each edge joins two nodes
  - Edges can be directed or undirected
- ▶ Real-world “graph”:
  - Road map
- ▶ Some uses:
  - Tracking links between web pages
  - Facebook

Operations Provided	Efficiency
Find	$O(n)$
Add/remove	$O(1)$ or $O(n)$ or $O(n^2)$

Depends on implementation  
(time/space trade off)

# Networks

- ▶ Graph whose edges have numeric labels
- ▶ Examples (labels):
  - Road map (mileage)
  - Airline's flight map (flying time)
  - Plumbing system (gallons per minute)
  - Computer network (bits/second)
- ▶ Famous problems:
  - Shortest path
  - Maximum flow
  - Minimal spanning tree
  - Traveling salesman
  - Four-coloring problem for planar graphs

# Common ADTs

- ▶ Array
- ▶ List
  - Array List
  - Linked List
- ▶ Stack
- ▶ Queue
- ▶ Set
  - Tree Set
  - Hash Set
- ▶ Map
  - Tree Map
  - Hash Map
- ▶ Priority Queue
- ▶ Tree
- ▶ Graph

We'll implement and use nearly all of these, some multiple ways. And a few other data structures.



# Data Structure Summary

Structure	find	insert/remove	Comments
Array	$O(n)$	can't do it	Constant-time access by position
Stack	top only $O(1)$	top only $O(1)$	Easy to implement as an array.
Queue	front only $O(1)$	$O(1)$	insert rear, remove front.
ArrayList	$O(N)$ $O(\log N)$ if sorted	$O(N)$	Constant-time access by position Add at end: am. $O(1)$ , worst $O(N)$
Linked List	$O(N)$	$O(1)$	$O(N)$ to find insertion position.
HashSet/Map	$O(1)$	amort. $O(1)$ , worst $O(N)$	Not traversable in sorted order
TreeSet/Map	$O(\log N)$	$O(\log N)$	Traversable in sorted order
PriorityQueue	$O(1)$	$O(\log N)$	Can only find/remove smallest
Search Tree	$O(\log N)$	$O(\log N)$	If tree is balanced, $O(N)$ otherwise

\*Some of these are amortized, not worst-case.

# Often, one particular ADT and implementation is best for the problem

- ▶ Which ADT to use?
  - It depends. How do you access your data? By position? By key? Do you need to iterate through it? Do you need the min/max?
- ▶ Which implementation to use?
  - It also depends. How important is fast access vs fast add/remove? Does the data need to be ordered in any way? How much space do you have?
- ▶ But real life is often messier...

# How to figure this out?

- ▶ Use Java's Collections Framework.
  - Search for *Java 8 Collection*
  - With a partner, read the javadocs to answer the quiz questions. You only need to submit one quiz per pair. (Put both names at top)
- ▶ If you finish, you may work on your current CSSE230 assignments