

CSSE 230
Red-black trees

After today, you should be able to...

- ...determine if a tree is a valid red/black tree
- ...perform top-down insertion in a red/black tree

EditorTrees Milestone 1 due tomorrow night

- If submitted tonight, all will earn a late day.
- If submitted late, everyone is charged a late day.
 - Does everyone on my team have a late day?
 - See link from Moodle
 - One team has a member with 0 late days left
- Tomorrow's class will be project work time
 - Don't let your team down! Be here and be on time.

Feedback to help as you move on...

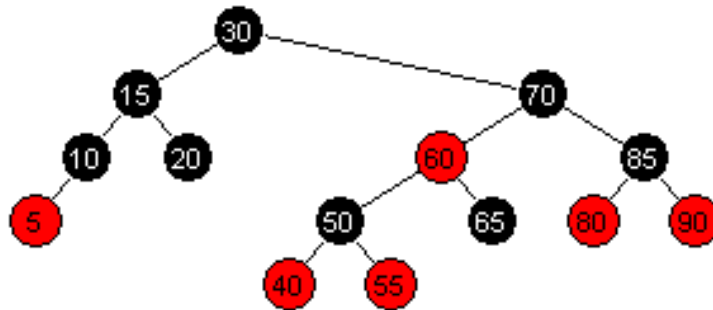
- Milestone 1 is graded on unit tests only.
- But...be sure to fix efficiency issues for the future
 - See final notes in specification
 - Cannot recalculate size or height to update balance codes or handle rotations.
 - You can recalculate **rank and balance codes**: these are $O(1)$ computations per node.
 - Suggestion: update rank (++) on the way down the tree.
 - Update balance codes and do rotations (which change rank and balance codes) on the way back up.
 - So each is $O(\log n)$ total
 - Know when you can stop! (day 14 slides have the algorithm for insertion, you'll have to think about deletion)

Red-Black Trees

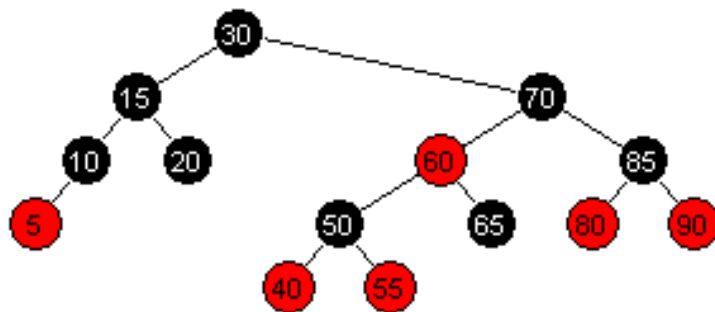
Another type of self-balancing search tree
with $O(\log N)$ performance

A red-black tree is a binary tree with 5 properties: ¹

1. It is a BST
2. Every node is either colored **red** or black.
3. The root is black.
4. No two successive nodes are **red**.
5. **Every path from the root to a null node** has the same number of black nodes (“perfect black balance”)



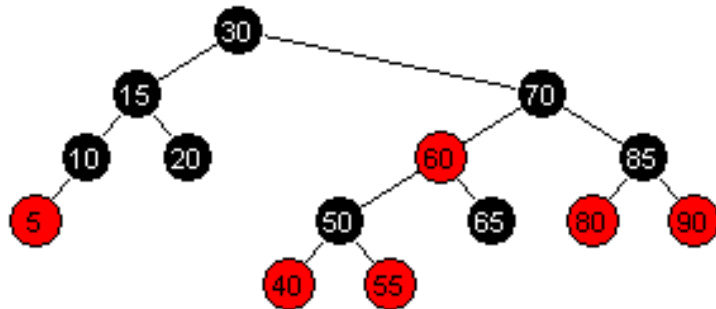
To search a red-black tree, just ignore the colors



Runtime is $O(\text{height})$

Since it's a BST, runtime of insert and delete should also be $O(\text{height})$

How tall is a red-black tree?



Best-case: if all nodes black, it is $\sim \log n$.

Worst case: every other node on the longest path is red. Height $\sim 2 \log n$.

Note: **Not height-balanced:**

Sometimes taller but often shorter on average.

Bottom-Up Insertion Strategy

2

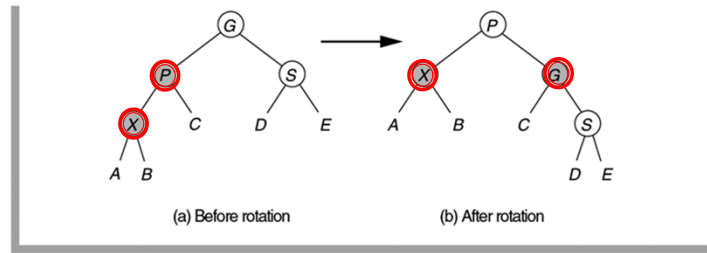
- Like BST:
 - Insert at leaf
 - Color it red (to keep perfect black balance)
- But could make two reds in a row?
 - On the recursive travel back up the tree (like AVL),
 - **rotate** (single- and double-, like AVL)
 - and **recolor** (new)
 - Show now that various “rotation+recoloring”s fix two reds in a row while maintaining black balance.
- At end of insert, always make root **of the entire tree** black (to fix property 3).

2 Reds in a row, with red outer grandchild and black sibling

2

figure 19.35

If *S* is black, a single rotation between parent and grandparent, with appropriate color changes, restores property 3 if *X* is an outside grandchild.



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2 Reds in a row, with red inner grandchild and black sibling

2

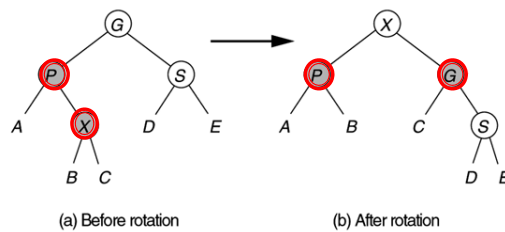


figure 19.36

If *S* is black, a double rotation involving *X*, the parent, and the grandparent, with appropriate color changes, restores property 3 if *X* is an inside grandchild.

2 Reds in a row, with red outer grandchild and red sibling

2

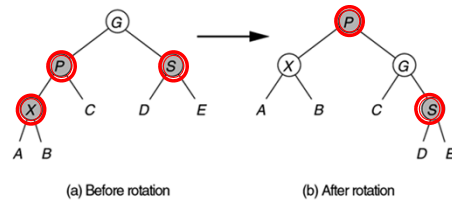


figure 19.37

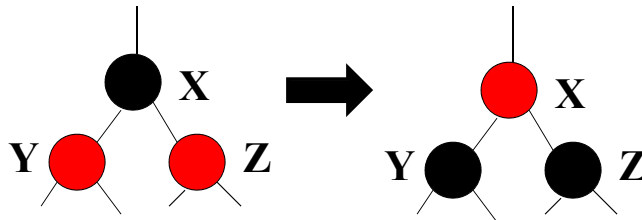
If *S* is red, a single rotation between parent and grandparent, with appropriate color changes, restores property 3 between *X* and *P*.

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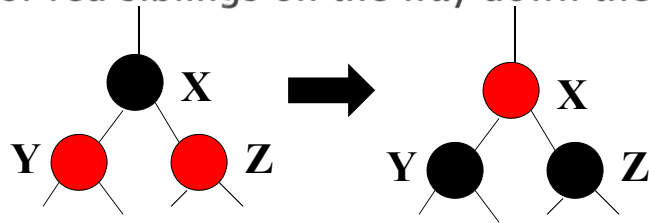
Case 3 (red sibling) can force us to do multiple rotations recursively

2

- Bottom-Up insertion strategy must be recursive.
- An alternative:
 - If we ever had a black node with two red children, swap the colors and black balance stays.
 - Details next...



Top-down insertion strategy:
 Recolor red siblings on the way down the tree



Situation: A black node with two red children.

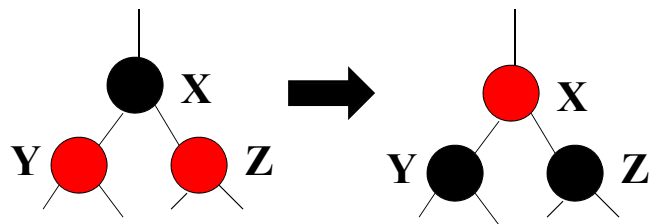
Action:

- Recolor the node **red** and the children **black** (if root, make black).
- If the parent is **red**, perform rotations, otherwise continue down the tree

Does this change black balance? No.

Top-Down Insertion Strategy

2



- On the way down the tree to the insertion point, if ever see a black node with two red children, swap the colors.
 If X's parent is **red**, perform rotations, otherwise continue down the tree
- The rotations are done **while traversing down the tree** to the insertion point.
 - **Avoid rotating into case (c) (2 red siblings) altogether.**
- Top-Down insertion can be done with loops without recursion or parent pointers, **so is slightly faster.**

Insertion summary

- Rotate when an insertion or color flip produces two successive red nodes.
- Rotations are just like those for AVL trees:
 - If the two red nodes are both left children or both right children, perform a *single rotation*.
 - Otherwise, perform a *double rotation*.
- Except we **recolor nodes** instead of adjusting their heights or balance codes.

Testing

3

1. Insert: 1, 2, 3, 4, 5, 6, 7, 8
2. Insert: 7, 6, 5, 4, 3, 2, 1, 1
 - Relationship with (1)?
 - Duplicates not inserted.
3. Insert: 10, 85, 15, 70, 20, 60, 30, 50, 65, 80, 90, 40, 5, 55
4. Use applet [linked to in Schedule for today] to check your work.

Summary

- Java uses:
- Slightly faster than AVL trees
- What's the catch?
 - Need to maintain pointers to lots of nodes (child, parent, grandparent, great-grandparent, great-great-grandparent)
- The deletion algorithm is *nasty*.

```

java.util
Class TreeMap<K,V>

java.lang.Object
  java.util.AbstractMap<K,V>
    java.util.TreeMap<K,V>

Type Parameters:
  K - the type of keys maintained by this map
  V - the type of mapped values

All Implemented Interfaces:
  Serializable, Cloneable, Map<K,V>, NavigableMap<K,V>

public class TreeMap<K,V>
  extends AbstractMap<K,V>
  implements NavigableMap<K,V>, Cloneable, Se
A Red-Black tree based NavigableMap implementation. T
This implementation provides guaranteed log(n) time cost fo

```