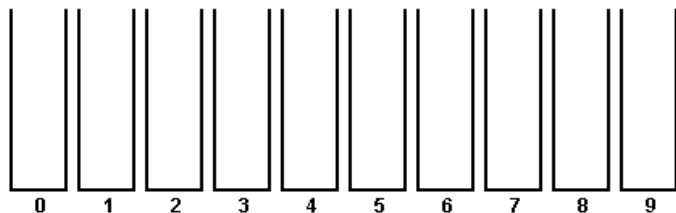


What is the min height of a tree with X external nodes?

CSSE 230

Sorting Lower Bound Radix Sort

Radix sort to the rescue ... sort of...



After today, you should be able to...

- ...explain why comparison-based sorts need at least $O(n \log n)$ time
- ... explain bucket sort
- ... explain radix sort
- ... explain the situations in which radix sort is faster than $O(n \log n)$

Announcements

- ▶ SortingRaces due Friday
- ▶ The sounds of sorting. Radix sort later.
 - <https://www.youtube.com/watch?v=kPRA0W1kECg>

A Lower-Bound on Sorting Time

We can't do much better than
what we already know how to
do.

What's the best best case?

- ▶ Lower bound for best case?
- ▶ A particular algorithm that achieves this?

What's the best worst case?

- ▶ Want a function $f(N)$ such that the worst case running time for **all sorting algorithms** is $\Omega(f(N))$
- ▶ How do we get a handle on “all sorting algorithms”?

Tricky!

What are “all sorting algorithms”?

- ▶ We can't list all sorting algorithms and analyze all of them
 - Why not?
- ▶ But we can find a **uniform representation** of any sorting algorithm that is based on **comparing** elements of the array to each other

First of all...

- ▶ The problem of sorting N elements is at least as hard as determining their ordering

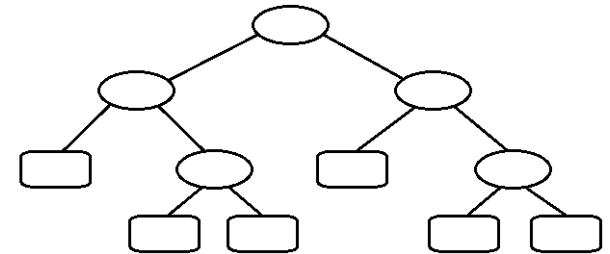
- e.g., determining that $a_3 < a_4 < a_1 < a_0 < a_2$

0	1	2	3	4
58	55	73	5	10

- sorting = determining order, then movement
- ▶ So any lower bound on all "order-determination" algorithms is also a lower bound on "all sorting algorithms"

Sort Decision Trees

- ▶ Let A be any **comparison-based algorithm** for sorting an array of distinct elements
- ▶ We can draw an EBT that corresponds to the comparisons that will be used by A to sort an array of N elements
 - This is called a **sort decision tree**
 - Internal nodes are comparisons
 - External nodes are orderings
- Different algorithms will have different trees



Insertion Sort

- ▶ Basic idea:

- Think of the array as having a **sorted part** (at the beginning) and an **unsorted part** (the rest)

0	1	2	3	4	5	6	7	8	9
38	44	87	2033	99	1500	100	90	239	748

- Get the **first** value in the unsorted part
- Insert it into the **correct** location in the sorted part, moving larger values up to make room

Repeat until
unsorted
part is
empty

So what?

- ▶ Minimum number of external nodes in a sort decision tree? (As a function of N)
- ▶ Is this number dependent on the algorithm?
- ▶ What's the height of the shortest EBT with that many external nodes?

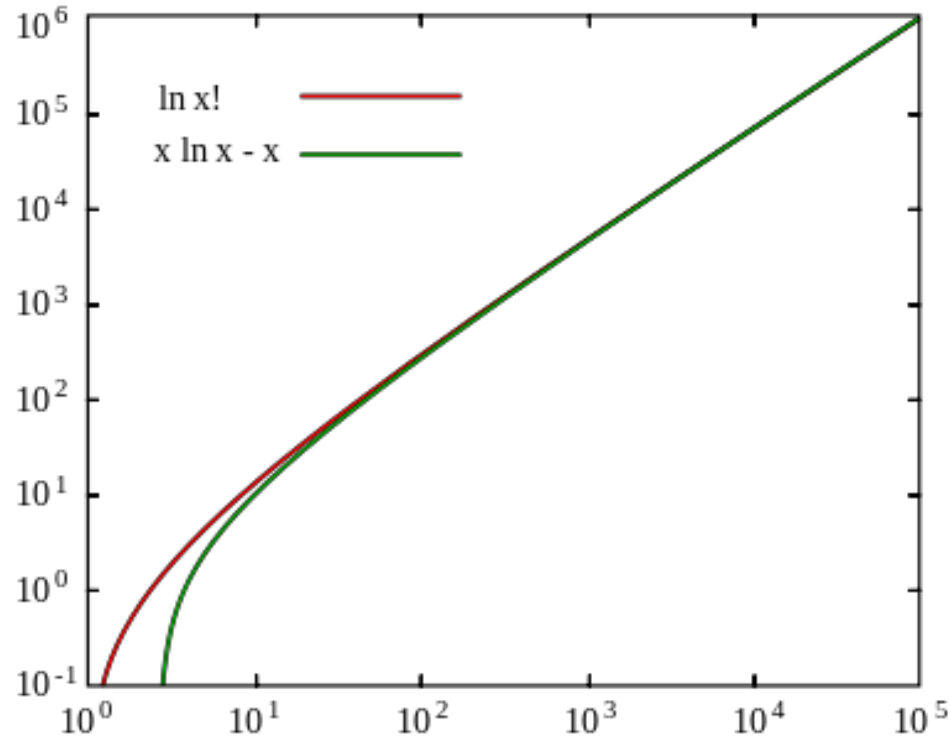
$$\lceil \log N! \rceil \approx N \log N - 1.44N = \Omega(N \log N)$$

No comparison-based sorting algorithm, known or not yet discovered, can ever do better than this!

An approximation for $\log(n!)$

- ▶ Use **Stirling's approximation**:

$$\ln n! = n \ln n - n + O(\ln(n))$$



Can we do better than $N \log N$?

- ▶ $\Omega(N \log N)$ is the best we can do if we compare items
- ▶ Can we sort without comparing items?

Yes, we can! We can avoid comparing items and still sort. This is fast if the range of data is small. Q5

- ▶ Observation:
 - For N items, if the range of data is less than N , then we have duplicates

- ▶ $O(N)$ sort: Bucket sort
 - Works if possible values come from limited range and have a uniform distribution over the range
 - Example: Exam grades histogram

- ▶ A variation: Radix sort

Radix sort

- ▶ A picture is worth 10^3 words, but an animation is worth 2^{10} pictures, so we will look at one.
- ▶ <http://www.cs.auckland.ac.nz/software/AlgAnim/radixsort.html> (good but blocked)
- ▶ https://www.youtube.com/watch?v=xuU-DS_5Z4g&src_vid=4S1L-pyQm7Y&feature=iv&annotation_id=annotation_133993417 (video, good basic idea, distracting zooms)
- ▶ <http://www.cs.usfca.edu/~galles/visualization/RadixSort.html> (good, uses single array)

RadixSort is almost $O(n)$

- ▶ It is $O(kn)$
 - Looking back at the radix sort algorithm, what is k ?
- ▶ Look at some extreme cases:
 - If all integers in range 0–99 (so, many duplicates if N is large), then $k = \text{-----}$
 - If all N integers are distinct, $k = \text{-----}$