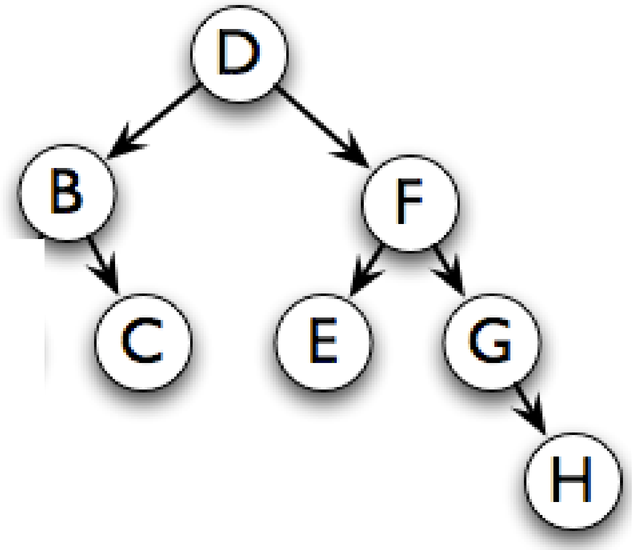
Q1-2



CSSE 230 Day 12

Height-Balanced Trees

After today, you should be able to…
…give the minimum number of nodes in a height-balanced tree
…explain why the height of a height-balanced trees is O(log n)
…help write an induction proof

Today's Agenda

- Announcements
  - EditorTrees team preferences due tonight
  - Exam 2 (programming only) in class on Thursday (day 14)

- Another induction example
- Recap: The need for balanced trees
- Analysis of worst case for height-balanced (AVL) trees

A useful result… by way of induction

- Recall our definition of the Fibonacci numbers:
  - $F_0 = 0$, $F_1 = 1$, $F_n = F_{n-1} + F_{n-2}$

- Prove the closed form:

7.8 Prove by induction the formula

$$F_N = \frac{1}{\sqrt{5}}\left(\left(\frac{(1+\sqrt{5})}{2}\right)^N - \left(\frac{1-\sqrt{5}}{2}\right)^N\right)$$

Recall: How to show that property P(n) is true for all $n \geq n_0$:

(1) Show the base case(s) directly

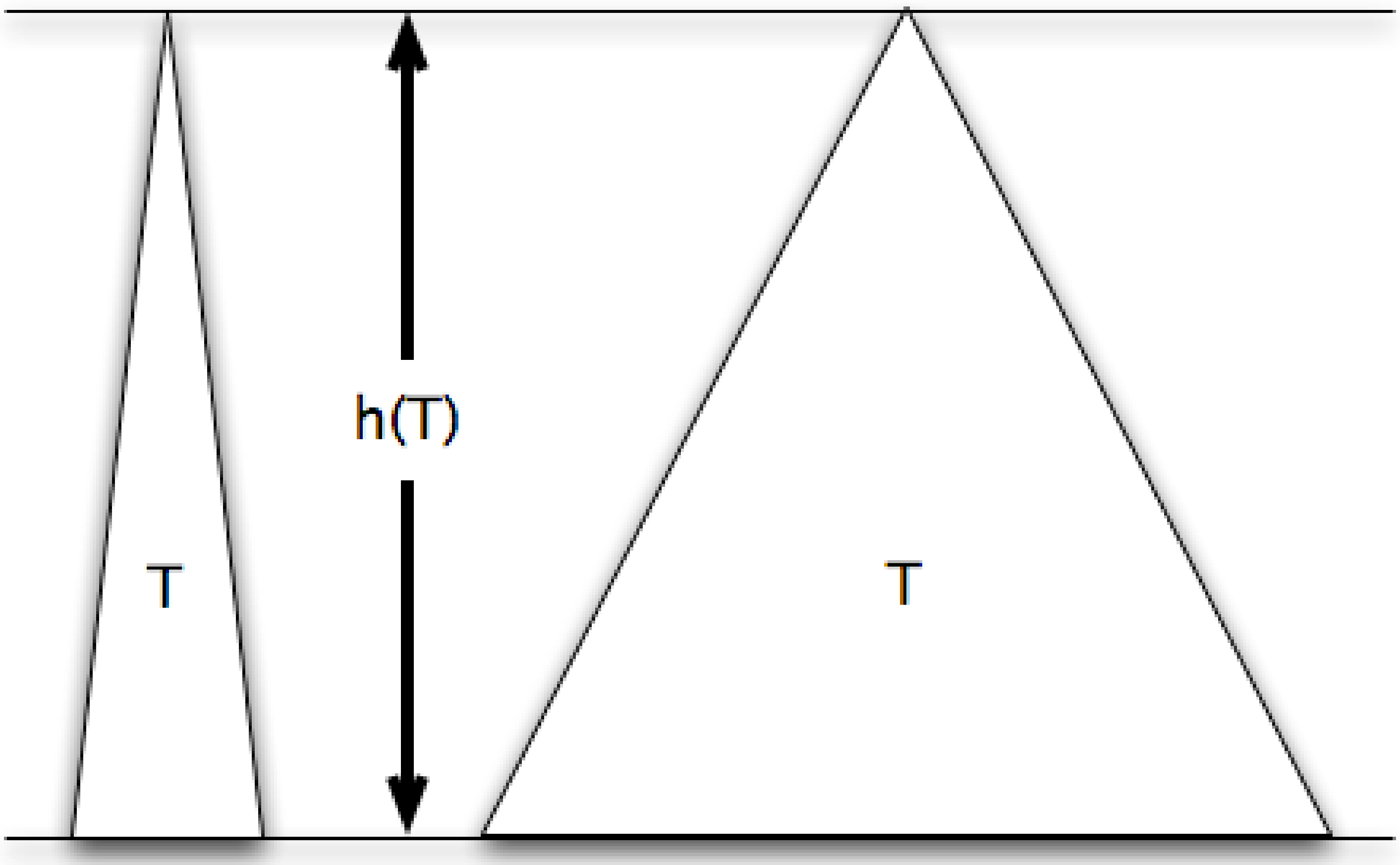
(2) Show that if P(j) is true for all j with $n_0 \leq j < k$, then P(k) is true also

Details of step 2:

a. Fix "arbitrary but specific" k ≥ _______.
b. Write the induction hypothesis: assume P(j) is true $\forall j : n_0 \leq j < k$
c. Prove P(k), using the induction hypothesis.

	A	B	C	D	E
1		1/SQRT(5) =	0.447213595		
2		(1 + SQRT(5))/2 =	1.618033989		
3		(1 - SQRT(5))/2 =	-0.618033989		
4					
5			**Fibonacci**		**Fibonacci**
6			**Open Form**		**Closed Form**
7	**n**	$\mathbf{f_n =}$	$\mathbf{f_{n-1} + f_{n-2}}$		**C$1*(POWER(C$2,n) - POWER(C$3,n))**
8					
9	0	f0 =	0		0
10	1	f1 =	1		1
11	2	f2 =	1		1
12	3	f3 =	2		2
13	4	f4 =	3		3
14	5	f5 =	5		5
15	6	f6 =	8		8
16	7	f7 =	13		13
17	8	f8 =	21		21
18	9	f9 =	34		34
19	10	f10 =	55		55
20	11	f11 =	89		89
21	12	f12 =	144		144
22	13	f13 =	233		233
23	14	f14 =	377		377
24	15	f15 =	610		610
25	16	f16 =	987		987
26	17	f17 =	1597		1597
27	18	f18 =	2584		2584
28	19	f19 =	4181		4181
29	20	f20 =	6765		6765

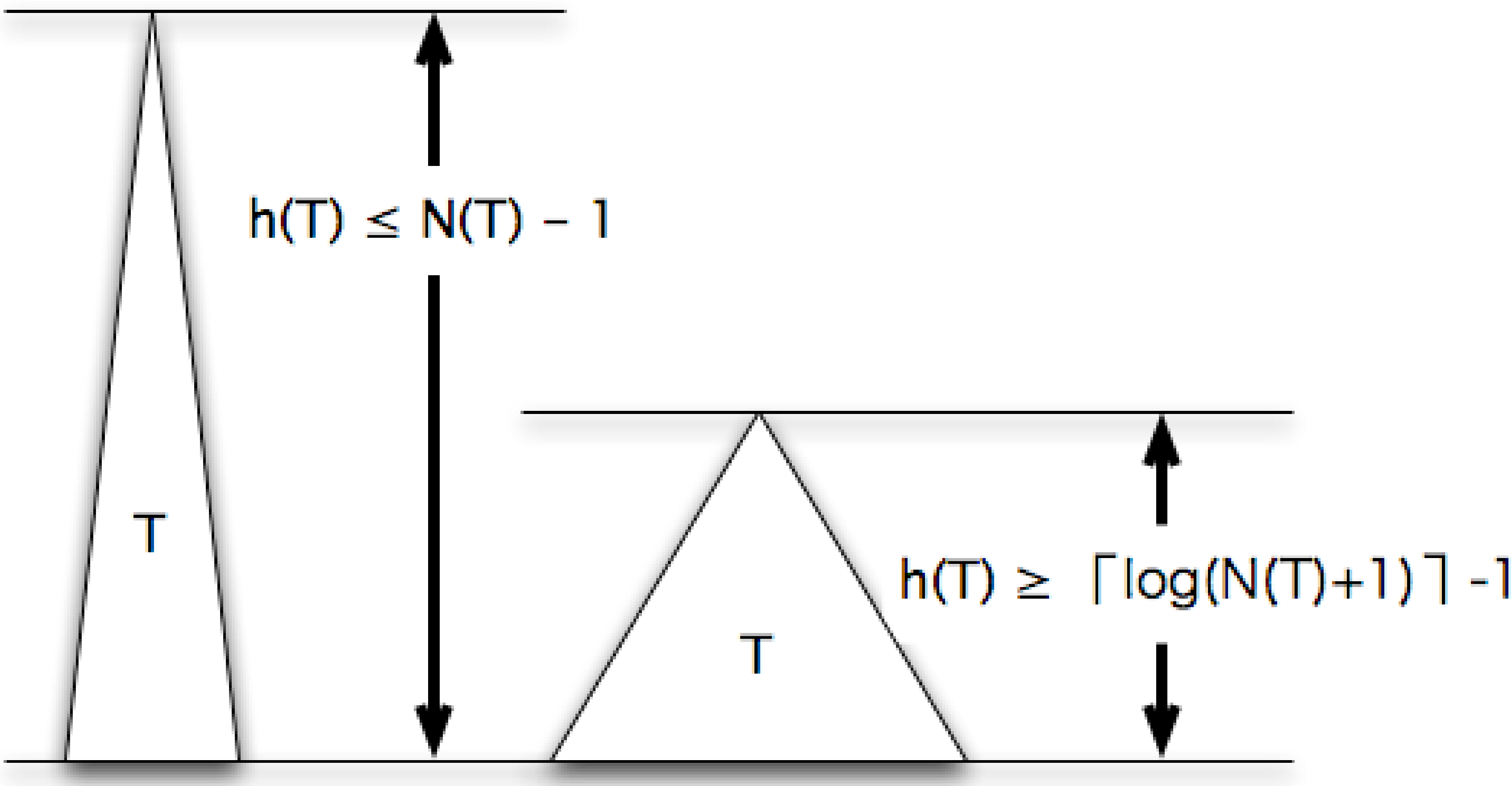
Review: The number of nodes in a tree with height h(T) is bounded



$N(T) \geq h(T) + 1$

$N(T) \leq 2^{h(T)+1} - 1$

Review: Therefore the height of a tree with N(T) nodes is also bounded



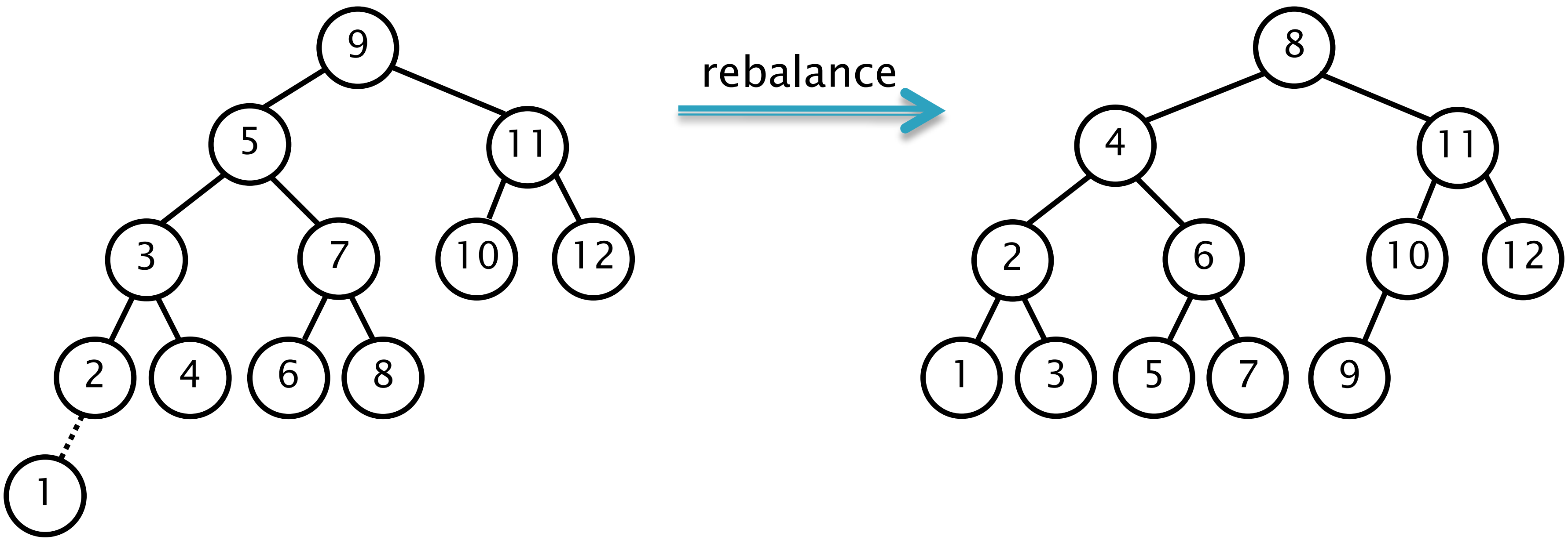
We want to keep trees balanced so that the run time of BST algorithms is minimized

- BST algorithms are O(h(T))
- Minimum value of h(T) is $\lceil \log(N(T) + 1) \rceil - 1$
- Can we rearrange the tree after an insertion to guarantee that h(T) is always **minimized**?

But keeping complete balance is too expensive!

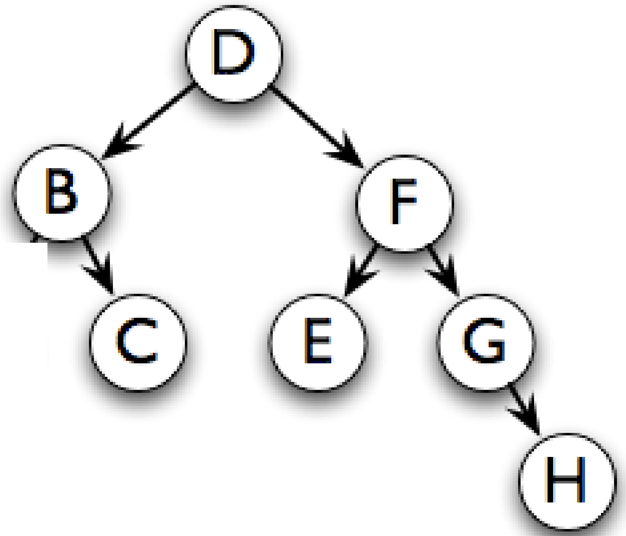
Q5

- Consider inserting 1 in the following tree.
- What does it take to get back to complete balance?
- Keeping completely balanced is too expensive:
  - O(N) to rebalance after insertion or deletion



Solution: Height Balanced Trees (less is more)

Height-Balanced Trees have subtrees whose heights differ by at most 1



Still height-balanced?

More precisely, a binary tree **T** is height balanced if

T is empty, or if

| **height**(T_L) – **height**(T_R) | $\leq$ 1, and

T_L and T_R are both height balanced.

What is the tallest height-balanced tree with N nodes?

Is it taller than a completely balanced tree?

- Consider the dual concept: find the minimum number of nodes for height h.

A binary search tree **T** is height balanced if

T is empty, or if

| **height(** T_L **)** – **height(** T_R **)** | ≤ **1**, and

T_L and T_R are both height balanced.

An AVL tree is a height-balanced BST that maintains balance using “rotations”

- Named for authors of original paper, **A**delson-**V**elskii and **L**andis (1962).
- Max. height of an AVL tree with **N** nodes is:
  $$\mathbf{H < 1.44 \log (N+2) - 1.328 = O(\log N)}$$

Our goal is to rebalance an AVL tree after insert/delete in O(log n) time

Q10

- Why?
- Worst cases for BST operations are **O(h(T))**
  - find, insert, and delete
- **h(T)** can vary from **O(log N)** to **O(N)**

- Height of a height-balanced tree is **O(log N)**

- So if we can rebalance after insert or delete in **O(log N)**, then **all** operations are **O(log N)**