Pick up an in-class quiz from the table near the door

#### CSSE 230 Data Structures and Algorithm Analysis Day 1



- two visual representations



### **Student Introductions**

#### Roll call

- Introduce yourself to the person next to you
- I'll soon post an assignment to Moodle that asks you to share more with classmates on a Piazza discussion forum, e.g., what's your favorite food, what are your hobbies, types of work you've done, etc.

And intro to daily quizzes, worth 5% of grade: Q1

#### Introductions

#### Joe Hollingsworth, aka Dr. Holly

- At R-H since 2018. CSSE 220 in FallQ
- B.S. Indiana University, CS
- M.S. Purdue University, CS
- Ph.D. Ohio State University, Software Engineering
- Special interests in formal methods, software design, how to best teach computing
- Courses taught at Rose:
  - CSSE220 (FallQ 2018), CSSE230 (WinterQ 2019)
- Hobbies: cycling, running, learning Spanish, travel

# Goal: independently design, develop, and debug software that uses correct, clear, and efficient algorithms and data structures



#### Why *efficient* algorithms? Here's \$1,000,000,000:



- Find serial number KB46279860I
- If unsorted, you could look at all 10 million bills.
- If sorted by serial number, binary search finds it by only looking at \_\_\_\_\_ bills.

https://commons.wikimedia.org/wiki/File:Oenbilliondollar.jpg

### How to succeed in CSSE230

- Work hard
  - Re-do CSSE220 stuff as needed to make sure your foundations (recursion and linked lists) are strong
- Take initiative in learning
  - Read the text, search Javadocs, come for help
- Focus while in this class
  - <u>https://www.rose-</u> <u>hulman.edu/class/cs/csse230/201820/MiscDocuments/LaptopsA</u> <u>reGreatButNotDuringaLectureoraMeeting.pdf</u> (11/26/2017 NYT)
- Start early and plan for no all-nighters
  - Two assignments each week: 1 homework set and 1 major program
- Never give or use someone else's answers

### Tools

- Moodle Site: <u>https://moodle.rose-hulman.edu/course/view.php?id=49906</u>
  - schedule, reading/HW/program assignments, room #s!
  - Read the Syllabus: Tomorrow's quiz will start with questions about it.
  - gradebook, homework pdf turn-in, peer evaluations, solutions
- www.piazza.com, not email: homework questions and announcements
  - If you email me, I'll reply, "Great question! Please post it to Piazza"
  - It should auto-email you whenever there is a post.
- moodle.rose-hulman.edu: gradebook, homework pdf turn-in, peer evaluations, solutions

## After today's class, you will be able to...

- analyze runtimes of code snippets by counting instructions.
- explain why arrays need to grow as data is added.
- derive the average and worst case time to insert an item into an array [GrowableArray exercise]

#### Analysis/Math Review

#### Notation

Floor

 java.lang.Math, provides the static methods floor() and ceil()

#### Summations

- Summations
  - general definition:

$$\sum_{i=s}^{t} f(i) = f(s) + f(s+1) + f(s+2) + \dots + f(t)$$

 where f is a function, s is the start index, and t is the end index

#### Geometric progressions: each term is a constant multiple of the previous term

- Geometric progression:  $f(i) = a^{i}$ 
  - given an integer  $n \ge 0$  and a real number  $0 \le a \ne 1$

$$\sum_{i=0}^{n} a^{i} = 1 + a + a^{2} + \dots + a^{n} = \frac{1 - a^{n+1}}{1 - a}$$
 Memorize  
this  
formula!

geometric progressions exhibit exponential growth

Exercise: What is 
$$\sum_{i=2}^{6} 3^i$$
 ?

This will be useful for today's Growable Arrays exercise!

The sum can also be written:

$$a^{n+1} - 1$$

$$a - 1$$

09 - 10

orize

Arithmetic progressions: constant difference Q11-12 Most important to us: a difference of 1

• Arithmetic progressions:



#### Visual proofs of the summation formula

$$\sum_{i=1}^{n} i = 1 + 2 + 3 + \dots + n = \frac{n^2 + n}{2}$$

- two visual representations



**Application: Runtime of Selection Sort** 

1 for (int i = n-1; i > 0; i--) { 2 int maxPos = 0;3 **for** (int j = 0; j <= i; j++) { **if** (a[j] > a[maxPos]) { 4 5 maxPos = j;6 } 7 8 swap a[maxPos] with a[i]; 9

#### **Selection Sort**

- Basic idea:
  - Think of the array as having a sorted part (at the beginning) and an unsorted part (the rest)

0	1	2	3	4	5	6	7	8	9
38	44	87	2033	99	1500	100	90	239	748

- Find the smallest value in the unsorted part
- Move it to the end of the sorted part (making the sorted part bigger and the unsorted part smaller)

Repeat until unsorted part is empty

## Application: Find exact and big-Oh runtime of Selection Sort

```
1 for (int i = n-1; i > 0; i--) {
2     int maxPos = 0;
3     for (int j = 0; j <= i; j++) {
4         if (a[j] > a[maxPos]) {
5             maxPos = j;
6         }
7     }
8     swap a[maxPos] with a[i];
9 }
```

- On what line is comparison performed?
- How many comparisons of array elements are executed? Exact? Big-Oh?
- How many times are array elements copied?

## Growable Array Analysis

An exercise in doubling, done by pairs of students

#### Arrays are ubiquitous

- Basis for ArrayLists, sorting, and hash tables
- Why? O(1) access to any position, regardless of the size of the array.
- Limitation of ArrayLists:
  - Fixed capacity!
  - If it fills, you need to re-allocate memory and copy items
    - How efficient is this?
    - Consider two schemes: "add 1" and "double"

#### Work on Growable Array Exercise

- Work with a partner
- Hand in the document before you leave today if possible. Otherwise due start of day 2's class.
- Get help as needed from me and the assistants.

### Handy for Growable Arrays HW

Properties of logarithms

Properties of exponents

$$log_b(xy) = log_b(x) + log_b(y) \qquad a^{(b+c)} = a^b a^c$$

$$log_b(x/y) = log_b(x) - log_b(y) \qquad a^{bc} = (a^b)^c$$

$$log_b(x^{\alpha}) = \alpha log_b(x) \qquad a^b/_{a^c} = a^{(b-c)}$$

$$log_b(x) = \frac{log_a(x)}{log_a(b)} \qquad b = a^{log_a(b)}$$

$$a^{log_b(n)} = n^{log_b(a)} \qquad b^c = a^{c*log_a(b)}$$