

# CSSE 230 Day 6

## Intro to Trees

After today, you should be able to...

- ...use tree terminology
- ...write recursive tree functions

Checkout `BinarySearchTree` from SVN

# Announcements

- ▶ Review Day 5's quizzes on Java Collections and Data Structures
- ▶ Preview of HW3: includes an essay

# Observation about Stacks and Queues Infix $\rightarrow$ Postfix problem

- ▶ It must be  $O(n)$ , so you can't grow your strings character-by-character:
  - Strings are immutable, so characters must be copied. `s += "*"`  is as slow growing an array using the `+1` scheme
- ▶ Solution? Use a `StringBuilder`!
  - They have internal capacity, which doubles when full!
- ▶ See the example at the end of Warmup and Stretching's `ShapeText.java` for an example.

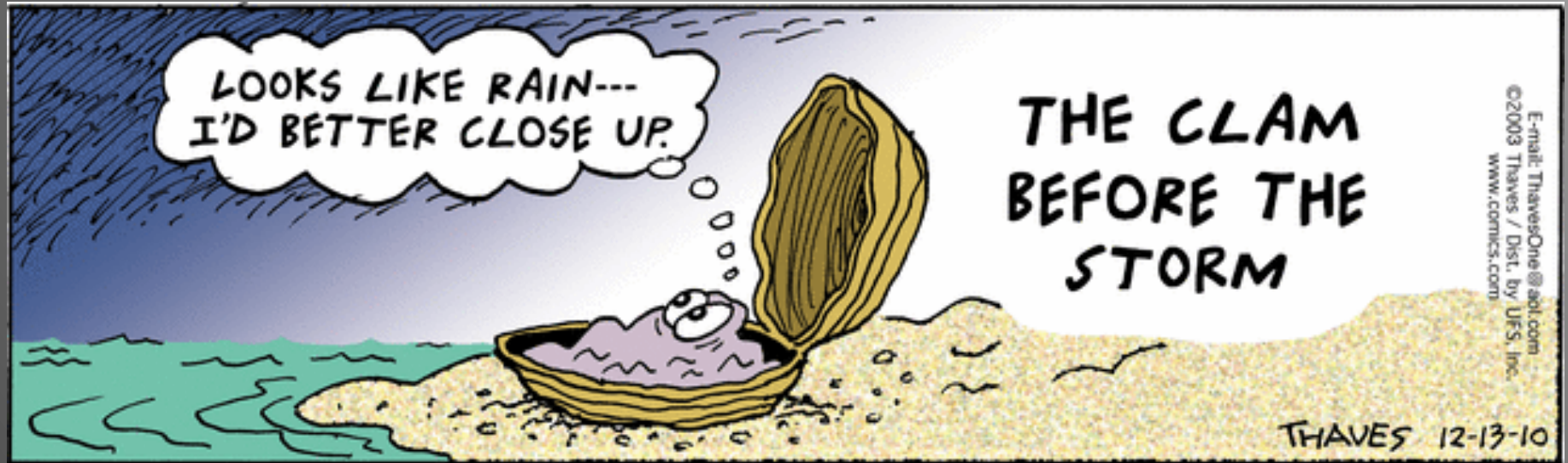
# Exam 1

- ▶ Exam 1 – Day 8: 7–9 pm
  - Coverage:
    - Everything from reading and lectures, Sessions 1–7
    - Programs: Warmup, Stacks and Queues
    - Homeworks 1–2
  - Allowed resources:
    - Written part:  $\frac{1}{2}$  of one side of 8.5 x 11 paper
      - Goal: to let you use formulas but force you to summarize.
    - Programming part:
      - Textbook
      - Eclipse (including programs you wrote in your repos)
      - Course web pages and materials on Moodle
      - Java API documentation
    - Two previous 230 Exam 1's are available in Moodle

# Exam 1 Possible Topics

- Written (50–70%):
  - Growable Arrays
  - MCSS
  - big  $O/\theta/\Omega$ : true/false, using definitions, code analysis
  - Binary search
  - ADT/Collections
  - Choosing an ADT to solve a given problem
- Programming (30–50%):
  - Implementing an ADT using an array, nodes, or another ADT
  - Writing an efficient algorithm to solve a simple array-based problem

# Questions?



# Next:

- ▶ an implementation that offers interesting benefits, but is more complex to code than arrays or lists...
- ▶ ... Trees!

# Trees

Introduction and terminology  
for three types





# Reminder: Specifying an ADT in Java

- ▶ Done with an interface, e.g., **java.util.Collection**

java.util

## Interface Collection<E>

boolean **add**(E o)

Ensures that this collection contains the specified element (optional operation).

boolean **contains**(Object o)

Returns true if this collection contains the specified element.

boolean **isEmpty**()

Returns true if this collection contains no elements.

boolean **remove**(Object o)

Removes a single instance of the specified element from this collection, if it is present (optional operation).

int **size**()

Returns the number of elements in this collection.

Iterator<E> **iterator**()

Returns an iterator over the elements in this collection.

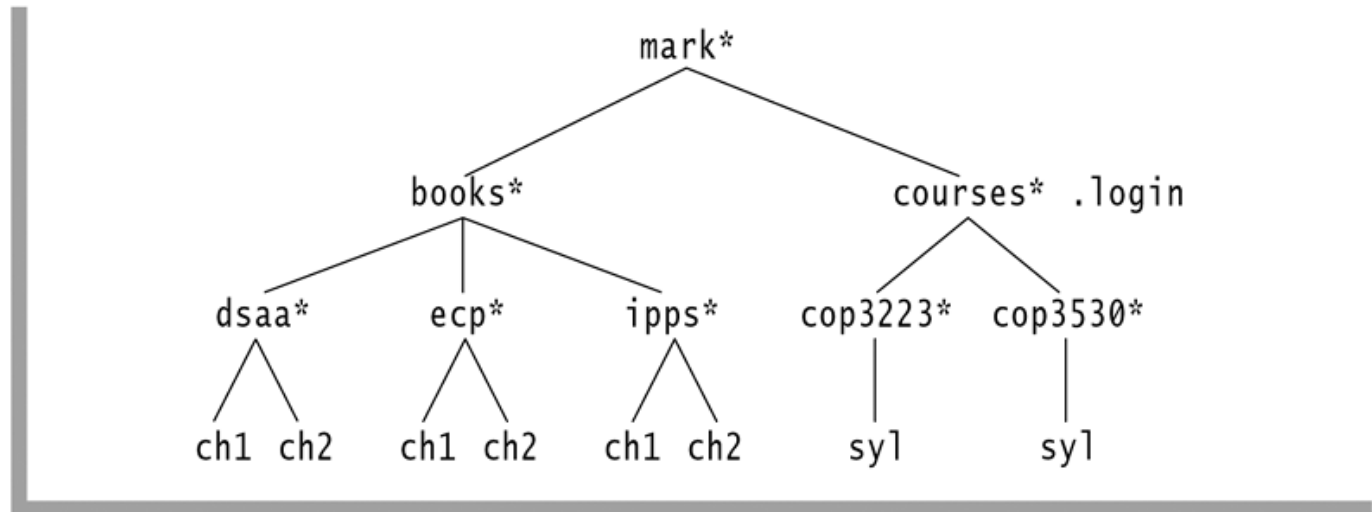
A “factory method”

# Trees in everyday life

- ▶ Class hierarchy tree (single inheritance only)
- ▶ Directory tree in a file system

**figure 18.4**

A Unix directory



# Traverse a Directory Tree

```
import java.io.File;

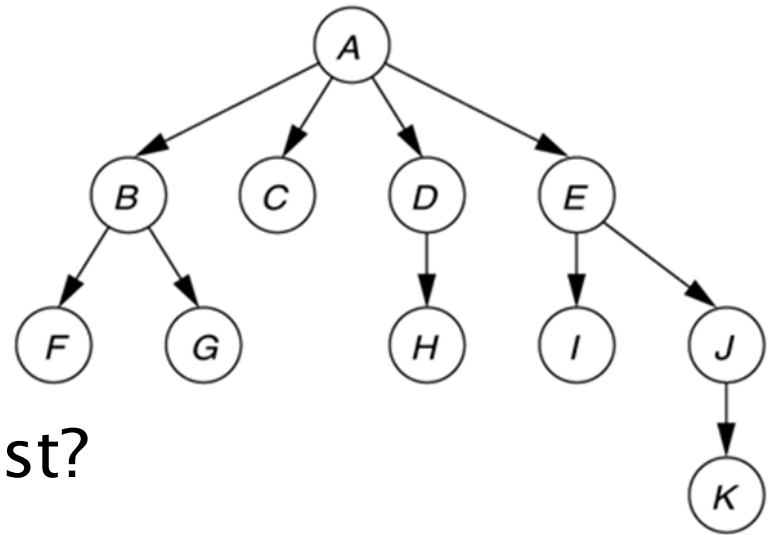
public class TraverseFiles {

    public static void main(String... args) {
        File[] files =
            new File("C:/EclipseWorkspaces/csse230-2014/BST2").listFiles();
        showFiles(files, 0);
    }

    public static void showFiles(File[] files, int indent) {
        for (File file : files) {
            if (file.isDirectory()) {
                System.out.println("        ".substring(0,indent) +
                    "Directory: " + file.getName());
                showFiles(file.listFiles(), indent+1); // Calls method again.
            } else {
                System.out.println("        ".substring(0,indent) +
                    "File: " + file.getName());
            }
        }
    }
}
```

# A General Tree—Global View

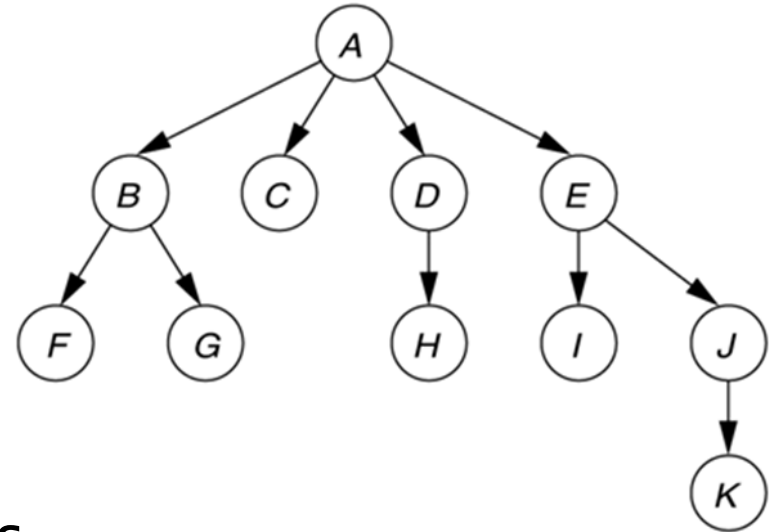
- ▶ A collection of **nodes**
- ▶ Nodes are connected by **directed** edges.
  - One special **root node** has no incoming edges
  - All other nodes have exactly one incoming edge
- ▶ One way that Computer Scientists are odd is that our trees usually have their root at the top!



- ▶ How are trees like a linked list?
- ▶ How are they different?

# Tree Terminology

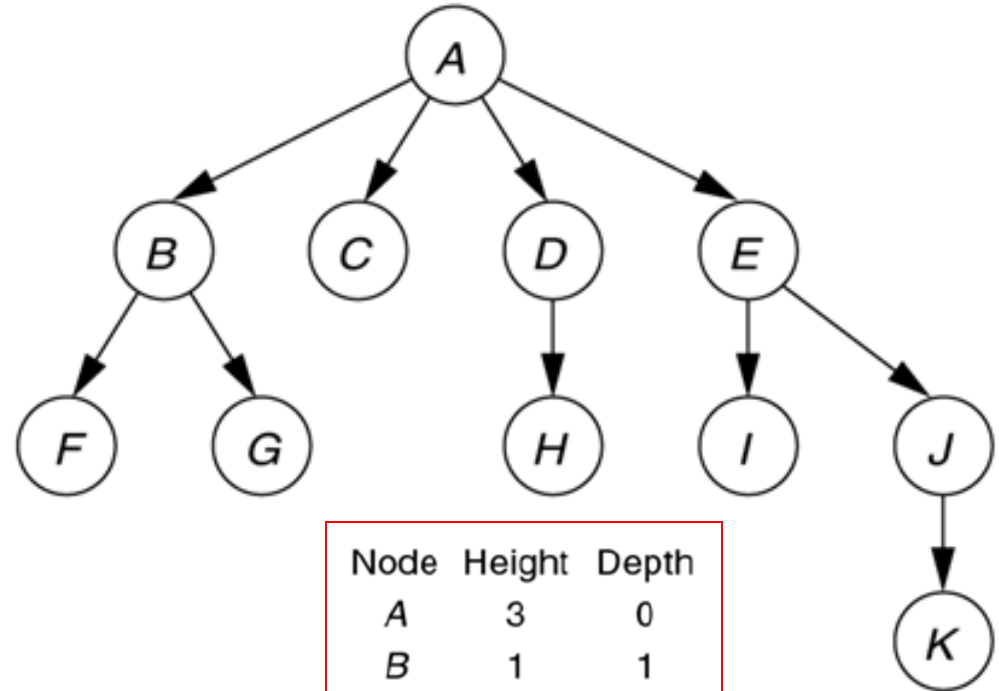
- ▶ Parent
- ▶ Child
- ▶ Grandparent
- ▶ Sibling
- ▶ Ancestors and descendants
- ▶ Proper ancestors, proper descendants
- ▶ Subtree
- ▶ Leaf, interior node
- ▶ Depth and height of a node
- ▶ Height of a tree



# Node height and depth examples

**figure 18.1**

A tree, with height and depth information



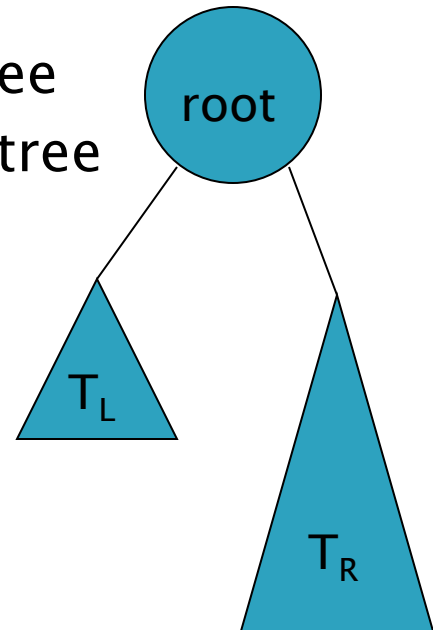
The height of a tree is the height of its root node.

Which is larger, the sum of the heights or the sum of the depths of all nodes in a tree?

Node	Height	Depth
A	3	0
B	1	1
C	0	1
D	1	1
E	2	1
F	0	2
G	0	2
H	0	2
I	0	2
J	1	2
K	0	3

# Binary Tree: Recursive definition

- ▶ A **Binary Tree** is either
  - **empty**, or
  - **consists** of:
    - a distinguished node called the root, which contains an element, and
    - A left subtree  $T_L$ , which is a binary tree
    - A right subtree  $T_R$ , which is a binary tree
- ▶ **Binary trees** contain at most 2 children



# Binary Search Trees (BST)

- ▶ Q: What property enables us to search BSTs efficiently?
- ▶ A: Every element in the left subtree is smaller than the root, and every element in the right subtree is larger than the root. And this is true at **every node**, not just the root.



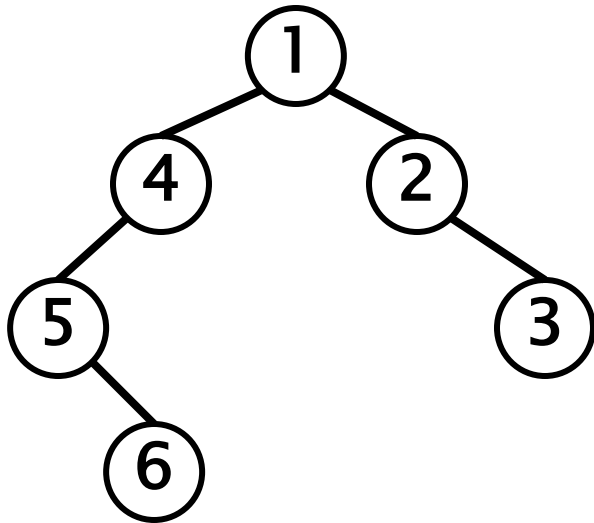
# Connections with Linked Lists

- ▶ Write `size()` for linked list
  - Non-`recursively`
  - `Recursively`
  
- ▶ Write `size()` for a tree
  - `Recursively`
  - Non-`recursively` (later)

# Growing Trees

- ▶ Let's start the BinarySearchTrees assignment: implement a **BinaryTree<T>** class

Test tree:



A single tiny recursive method for size will touch **every node in the tree**. Let's write, then watch in debugger.