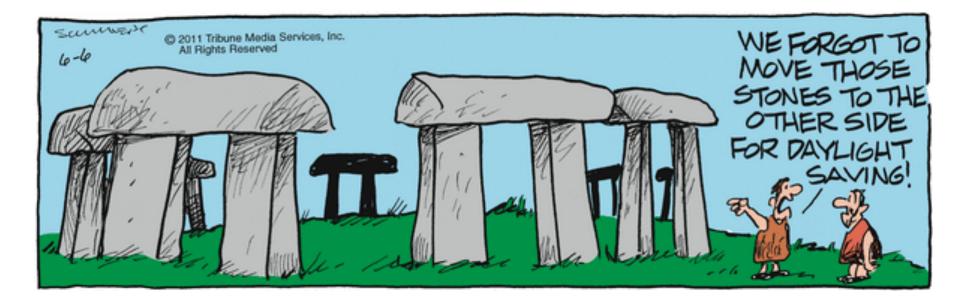


CSSE 230 Day 23 2D Trees

After today, you should be able to explain insert and nearest-neighbor in 2D trees ... implement these algorithms

Reminders/Announcements



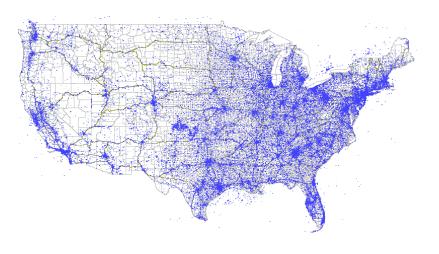
Reminders/Announcements

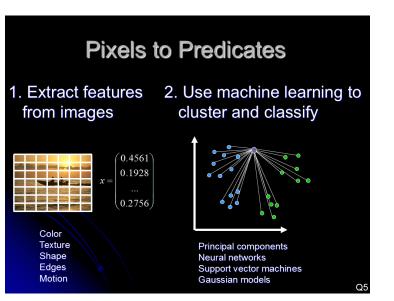
EditorTrees evals due tonight

- Would do ONLY if we want them to extend it on the exam:
- Before the final exam, copy your team's
 EditorTreesMilestone2 project to your individual
 CSSE 230 repository
 - Team > Update
 - Team > Disconnect
 - Before you press the Yes button, choose "Also Delete SVN metadata"
 - Team > Share Project > SVN > Next, choose your repo
 - Team>Commit
 - Just to be sure everything is there.

2D Data

- A large set of (x,y) points
- Which cell phone tower is closest to me?
- Which image is most like this one?
- In general:
 - Find the nearest neighbor of a query point (today).
 - Find or return all points in a certain range.





https://personalpages.manchester.ac.uk/staff/m.dodge/cybergeography/atlas/tower_maps_large.gif https://www.rose-hulman.edu/class/csse/csse463/201520/Slides/01%20Introduction.pdf

- List of points. Simple but slow
 - [p1, p2, ..., pN]
 - To find closest to q, find smallest of dist(q, p1), dist(q, p2), ...



2a

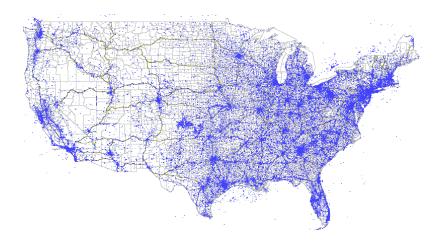
Representation	Average nearest-nbr efficiency
List of points	Ν

- List of points. Simple but slow
- Use a regular grid.
 - 2D array of lists
 - Faster, but which resolution?
 - Example, M=8

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Representation	Average nearest-nbr efficiency
List of points	Ν
Regular grid	1 + N/M ² but space = N + M ² , clustering degrades

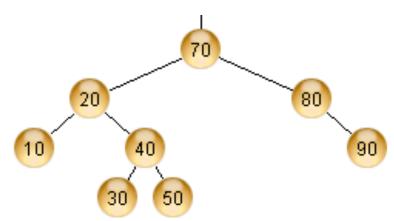
- List of points. Simple but slow
- Use a regular grid.
- ???



Representation	Average nearest-nbr efficiency
List of points	Ν
Regular grid	1 + N/M ² but space = N + M ² , clustering degrades
???	log N

Binary search trees partition the number line

- Split at 70
- Split at 20
- etc



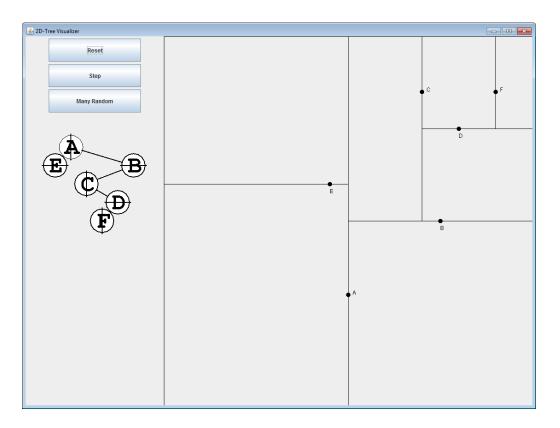
Any value inserted to the left of 30 must be in what range?

You can partition the coordinate plane with a variation of BSTs

- Root splits plane using x-coordinate and each level splits the plane in one direction only.
- Use the insert algorithm to build a tree from points: A (0.5, 0.7) B (0.75, 0.5) C (0.7, 0.15) D (0.8, 0.25) E (0.45, 0.4) F (0.9, 0.15)

You can partition the coordinate plane with a variation of BSTs

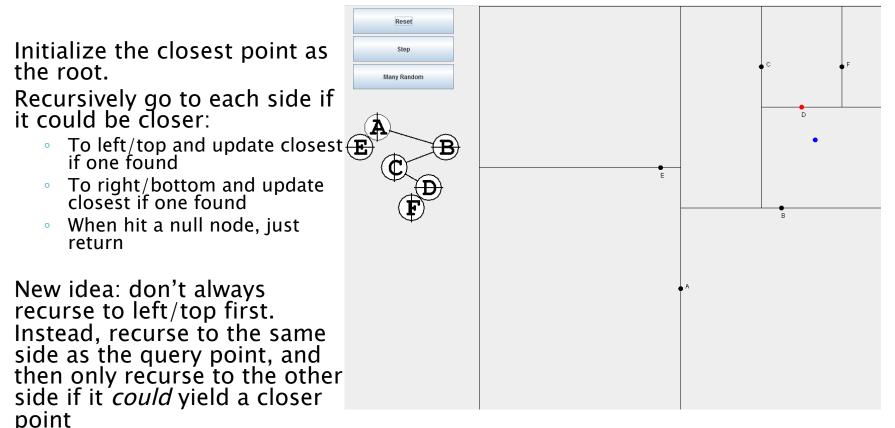
- Root splits plane using x-coordinate and each level splits the plane in one direction only.
- Use the insert algorithm to build a tree from points:
 A (0.5, 0.7)
 B (0.75, 0.5)
 C (0.7, 0.15)
 D (0.8, 0.25)



Nearest neighbor using a 2D Tree

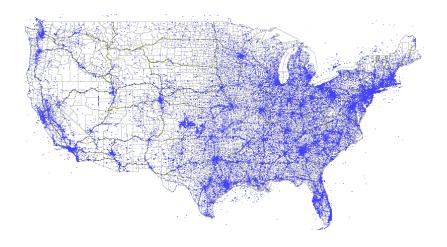
- Initialize the closest point as the root.
- Recursively go to each side if it could be closer:
 - To left/top and update closest if one found
 - To right/bottom and update closest if one found
 - When hit a null node, just return
- New idea: don't always recurse to left/top first. Instead, first recurse to the same side as the query point, and then only recurse to the other side if it could yield a closer point
 - To do this, each node also stores the bounds of rectangle it is part of
 - I give you a Rectangle class with a method to find closest distance from a point that that rectangle.

Nearest neighbor using a 2D Tree



 To do this, each node will also store the bounds of rectangle it is part of

- List of points. Simple but slow
- Use a regular grid.
- Use a 2D tree
 - You can find the nearest neighbor efficiently



Representation	Average nearest-nbr efficiency
List of points	Ν
Regular grid	1 + N/M ² but space = N/M ² +1, clustering degrades
2D tree	log N

2D Trees are useful

- Questions for thought:
 - How would you build a 3D tree?
 - ... a kD tree for arbitrary dimension k?
- Summarize now
- Assignment for this week:
 - Implement insert(Point), contains(Point), and nearest(Point) using a 2D tree.
 - There are unit tests for correctness and efficiency. You must earn the correctness points to be considered for the efficiency points – efficiency first!