

# CSSE 230 Day 23 <br> 2D Trees 

After today, you should be able to
... explain insert and nearest-neighbor in 2D trees
... implement these algorithms

## Reminders/Announcements



## Reminders/Announcements

- EditorTrees evals due tonight
- Would do ONLY if we want them to extend it on the exam:
- Before the final exam, copy your team's EditorTreesMilestone2 project to your individual CSSE 230 repository
- Team > Update
- Team > Disconnect
- Before you press the Yes button, choose "Also Delete SVN metadata"
- Team > Share Project > SVN > Next, choose your repo
- Team>Commit
- Just to be sure everything is there.


## 2D Data

- A large set of ( $x, y$ ) points
- Which cell phone tower is closest to me?

- Which image is most like this one?
- In general:
- Find the nearest neighbor of a query point (today).
- Find or return all points in a certain range.



## How to represent 2D data?

- List of points. Simple but slow
- [p1, p2, ..., pN]
- To find closest to q, find smallest of $\operatorname{dist}(q, p 1), \operatorname{dist}(q, p 2), \ldots$


| Representation | Average nearest-nbr efficiency |
| :--- | :--- |
| List of points | N |

## How to represent 2D data?

- List of points. Simple but slow
- Use a regular grid.
- 2D array of lists
- Faster, but which resolution?
- Example, $\mathrm{M}=8$


| Representation | Average nearest-nbr efficiency |
| :--- | :--- |
| List of points | N |
| Regular grid | $1+\mathrm{N} / \mathrm{M}^{2}$ <br> but space $=\mathrm{N}+\mathrm{M}^{2}$, clustering <br> degrades |
|  |  |

## How to represent 2D data?

- List of points. Simple but slow
- Use a regular grid.
- ???


| Representation | Average nearest-nbr efficiency |
| :--- | :--- |
| List of points | N |
| Regular grid | $1+\mathrm{N} / \mathrm{M}^{2}$ <br> but space $=\mathrm{N}+\mathrm{M}^{2}$, clustering <br> degrades |
| $? ? ?$ | log N |

# Binary search trees partition the number line 

- Split at 70
- Split at 20
- etc

- Any value inserted to the left of 30 must be in what range?


# You can partition the coordinate plane with a variation of BSTs 

- Root splits plane using x-coordinate and each level splits the plane in one direction only.
- Use the insert algorithm to build a tree from points:
A $(0.5,0.7)$
B $(0.75,0.5)$
C $(0.7,0.15)$
D $(0.8,0.25)$
E $(0.45,0.4)$
F $(0.9,0.15)$


## You can partition the coordinate plane with a variation of BSTs

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## Nearest neighbor using a 2D Tree

- Initialize the closest point as the root.
- Recursively go to each side if it could be closer:
- To left/top and update closest if one found
- To right/bottom and update closest if one found
- When hit a null node, just return
- New idea: don't always recurse to left/top first. Instead, first recurse to the same side as the query point, and then only recurse to the other side if it could yield a closer point
- To do this, each node also stores the bounds of rectangle it is part of
- I give you a Rectangle class with a method to find closest distance from a point that that rectangle.


## Nearest neighbor using a 2D Tree

Initialize the closest point as the root.
Recursively go to each side if it could be closer:

- To left/top and update closest (E) if one found
- To right/bottom and update closest if one found
- When hit a null node, just return

New idea: don't always recurse to left/top first. Instead, recurse to the same side as the query point, and then only recurse to the other side if it could yield a closer
 point

- To do this, each node will also store the bounds of rectangle it is part of


## How to represent 2D data?

- List of points. Simple but slow
- Use a regular grid.
- Use a 2D tree
- You can find the nearest neighbor efficiently


| Representation | Average nearest-nbr efficiency |
| :--- | :--- |
| List of points | N |
| Regular grid | $1+\mathrm{N} / \mathrm{M}^{2}$ <br> but space $=\mathrm{N} / \mathrm{M}^{2}+1$, <br> clustering degrades |
| 2D tree | $\log \mathrm{N}$ |

## 2D Trees are useful

- Questions for thought:
- How would you build a 3D tree?
- ... a kD tree for arbitrary dimension k?
- Summarize now
- Assignment for this week:
- Implement insert(Point), contains(Point), and nearest(Point) using a 2D tree.
- There are unit tests for correctness and efficiency. You must earn the correctness points to be considered for the efficiency points - efficiency first!

