

CSSE 230 Day 23

2D Trees

After today, you should be able to ...

... explain insert and nearest-neighbor in 2D trees

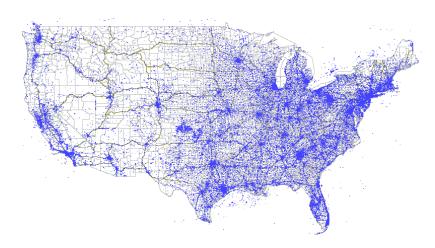
... implement these algorithms

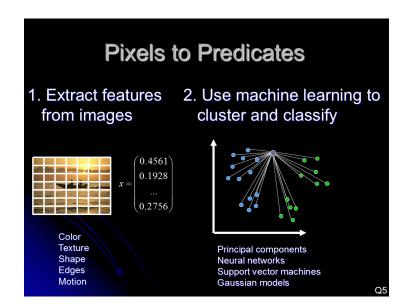
Reminders/Announcements

EditorTrees evals due tonight

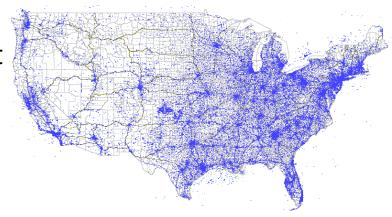
2D Data

- A large set of (x,y) points
- Which cell phone tower is closest to me?
- Which image is most like this one?
- In general:
 - Find the nearest neighbor of a query point (today).
 - Find or return all points in a certain range.



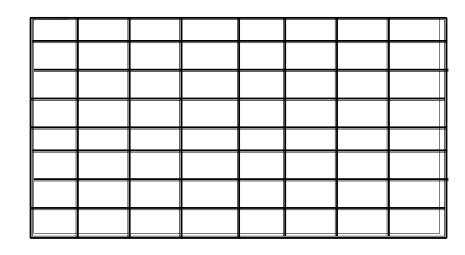


- List of points. Simple but slow
 - [p1, p2, ..., pN]
 - To find closest to q, find smallest of dist(q, p1), dist(q, p2), ...



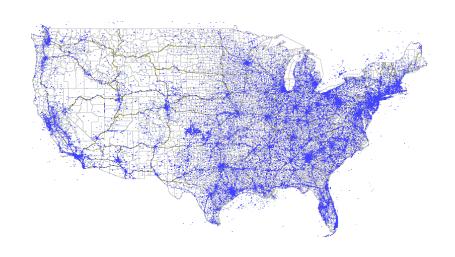
Representation	Average nearest-nbr efficiency
List of points	N

- List of points. Simple but slow
- Use a regular grid.
 - 2D array of lists
 - Faster, but which resolution?
 - Example, M=8



Representation	Average nearest-nbr efficiency
List of points	N
Regular grid	1 + N/M ² but space = N +M ² , clustering degrades

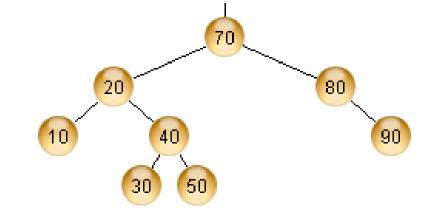
- List of points. Simple but slow
- Use a regular grid.
- > ???



Representation	Average nearest-nbr efficiency
List of points	N
Regular grid	$1 + N/M^2$ but space = $N + M^2$, clustering degrades
???	log N

Binary search trees partition the number line

- Split at 70
- Split at 20
- etc



Any value inserted to the left of 30 must be in what range?

You can partition the coordinate plane with a variation of BSTs

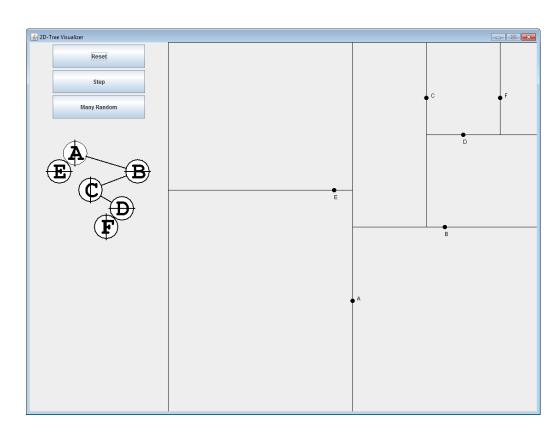
- Root splits plane using x-coordinate and each level splits the plane in one direction only.
- Use the insert algorithm to build a tree from points:

```
A (0.5, 0.7)
B (0.75, 0.5)
C (0.7, 0.15)
D (0.8, 0.25)
E (0.45, 0.4)
F (0.9, 0.15)
```

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Nearest neighbor using a 2D Tree

- Initialize the closest point as the root.
- Recursively go to each side if it could be closer:
 - To left/top and update closest if one found
 - To right/bottom and update closest if one found
 - When hit a null node, just return
- New idea: don't always recurse to left/top first. Instead, first recurse to the same side as the query point, and then only recurse to the other side if it could yield a closer point
 - To do this, each node also stores the bounds of rectangle it is part of
 - I give you a Rectangle class with a method to find closest distance from a point that that rectangle.

Nearest neighbor using a 2D Tree

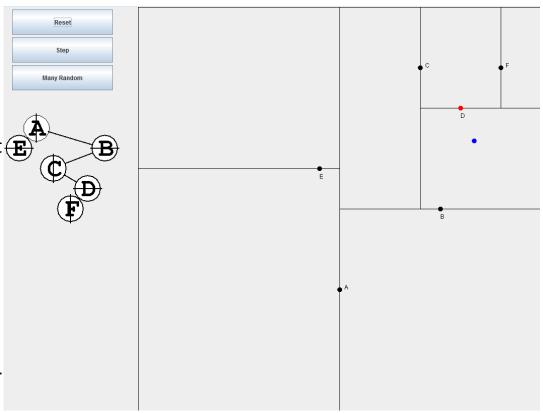
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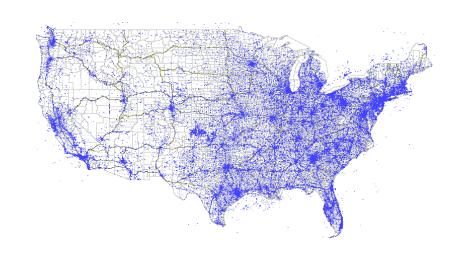
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New idea: don't always recurse to left/top first. Instead, recurse to the same side as the query point, and then only recurse to the other side if it *could* yield a closer point

 To do this, each node will also store the bounds of rectangle it is part of



- List of points. Simple but slow
- Use a regular grid.
- Use a 2D tree
 - You can find the nearest neighbor efficiently



Representation	Average nearest-nbr efficiency
List of points	N
Regular grid	1 + N/M ² but space = N/M ² +1, clustering degrades
2D tree	log N

2D Trees are useful

- Questions for thought:
 - How would you build a 3D tree?
 - ... a kD tree for arbitrary dimension D?
- Summarize now
- Assignment for this week:
 - Implement insert(Point), contains(Point), and nearest(Point) using a 2D tree