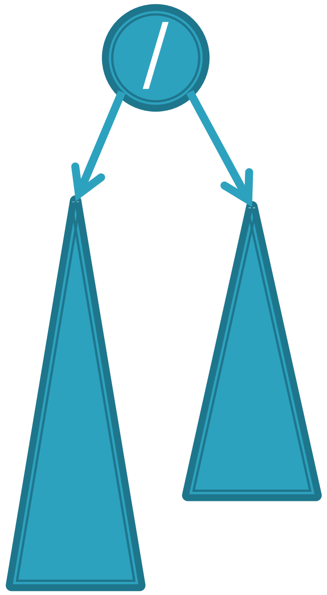


CSSE 230 Day 13

AVL trees and rotations

This week, you should be able to…
…perform rotations on height-balanced trees,
on paper and in code
… write a rotate() method
… search for the kth item in-order using rank

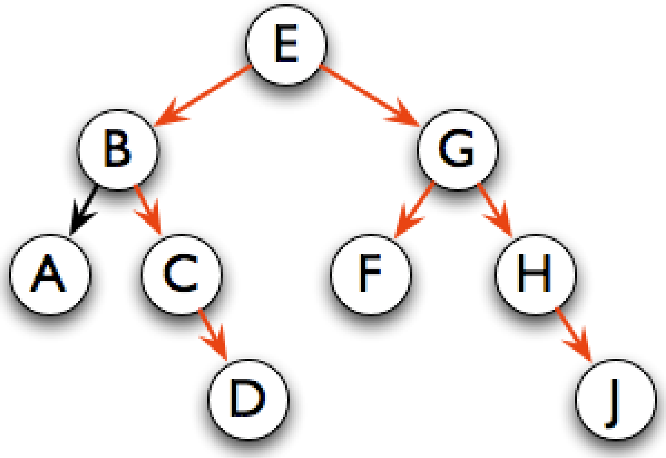
Announcements

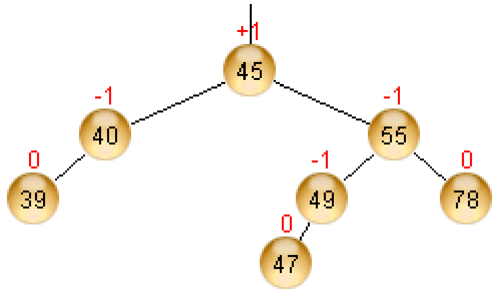
- Tests back later
- Mini-test next class
- See schedule page

Summary: for fast tree operations, we must keep tree somewhat balanced in O(log n) time

Q1

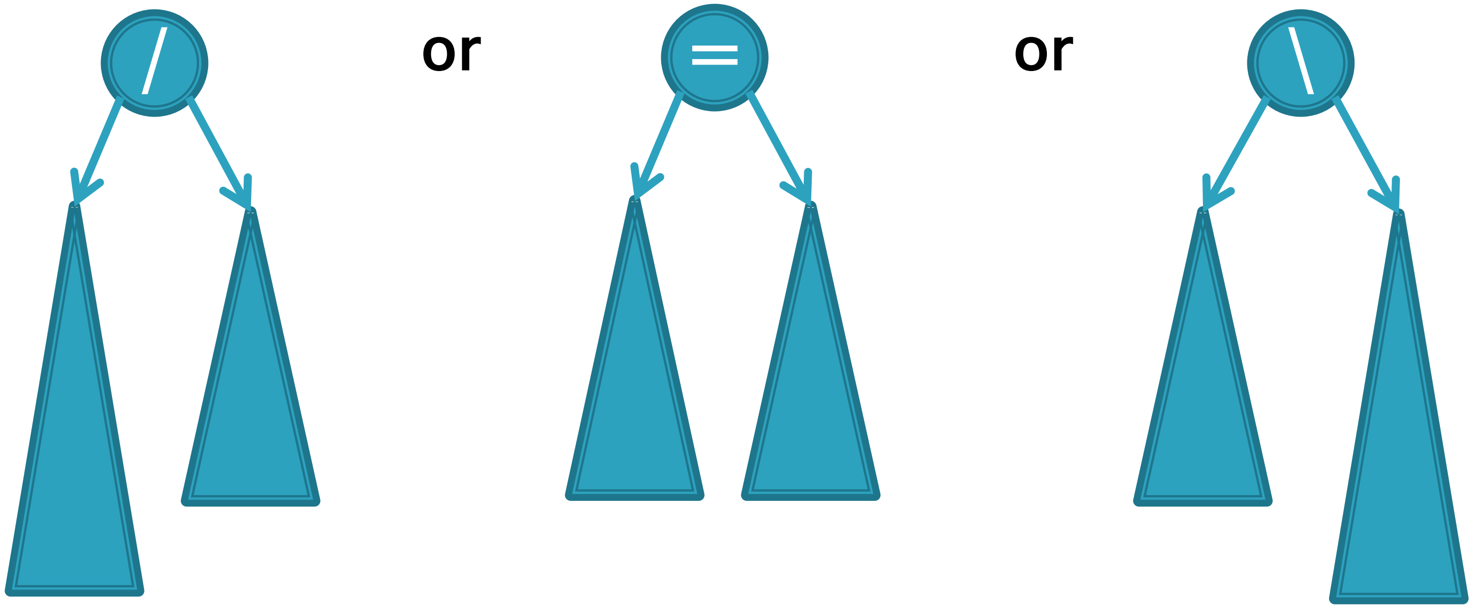
- Total time to do insert/delete =
  - Time to find the correct place to insert = O(height)
  - + time to detect an imbalance
  - + time to correct the imbalance
- If don't bother with balance:
- If try to keep perfect balance:
  - Height is O(log n) BUT …
  - But maintaining perfect balance is O(n)
- Height-balanced trees are still O(log n)
  - For T with height h, N(T) ≤ Fib(h+3) – 1
  - So H < 1.44 log (N+2) – 1.328 *
- AVL (Adelson-Velskii and Landis) trees maintain height-balance using rotations
- Are rotations O(log n)? We'll see…





AVL nodes are just like BinaryNodes, but also have an extra "balance code"

Q2



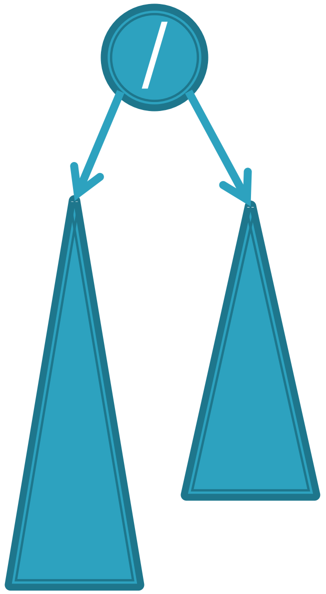
Different representations for / = \ :

- Just two bits in a low-level language
- Enum in a higher-level language

Using balance codes makes AVL Tree rebalancing efficient: O(log n)

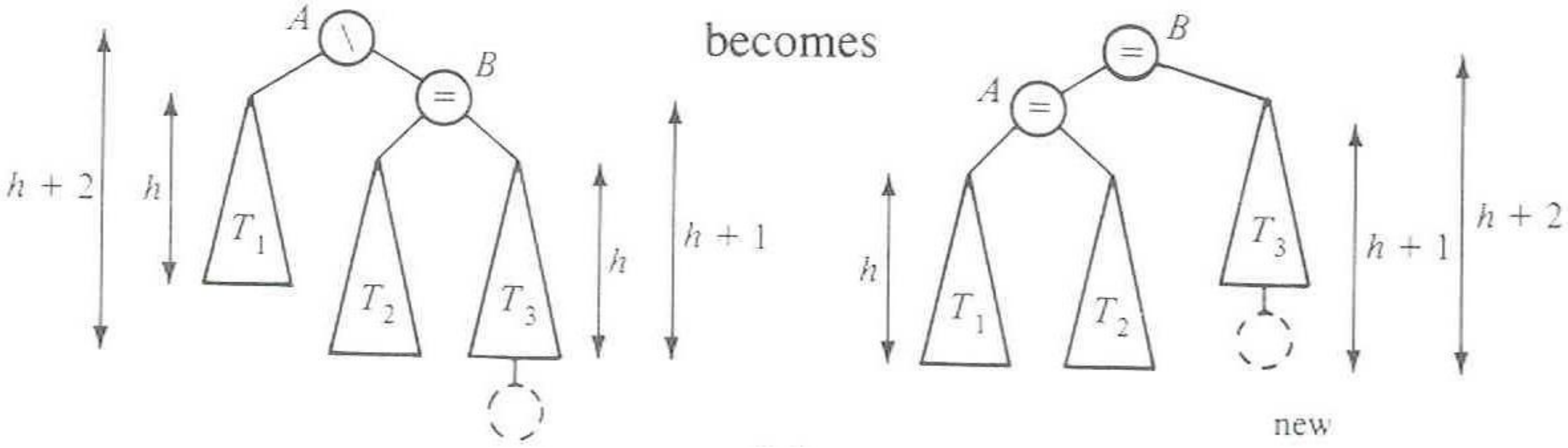
Q3

- Assume tree is height-balanced before insertion
- Insert as usual for a BST
- Move up from the newly inserted node to the lowest "unbalanced" node (if any)
  - Use the **balance code** to detect unbalance – how?
  - Why is this O(log n)?
    - We move up the tree to the root in worst case, NOT recursing into subtrees to calculate heights
- Do an appropriate rotation (see next slides) to balance the sub-tree rooted at this unbalanced node



Four types of rotations are required to remove different cases of tree imbalances

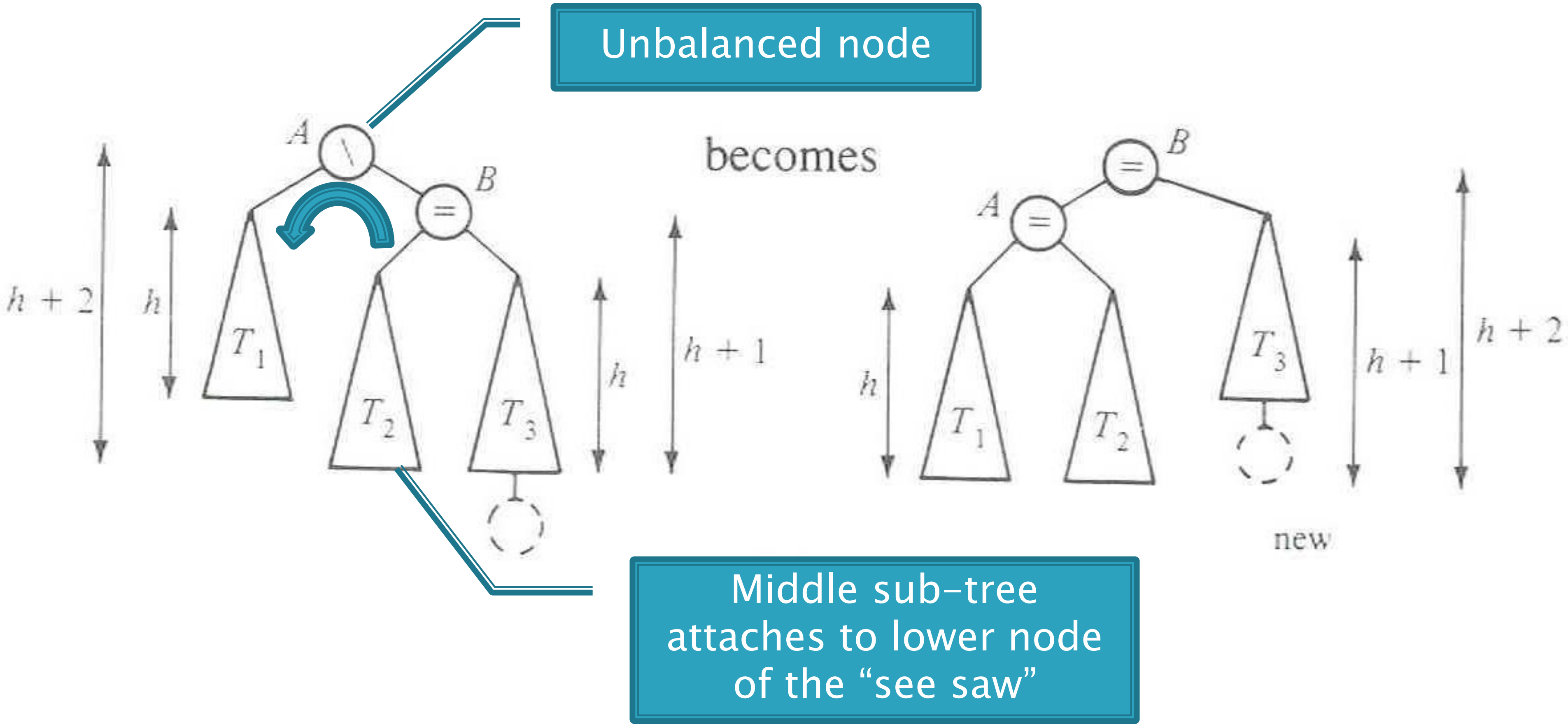
- For example, a *single left rotation*:



We rotate by pulling the “too tall” sub-tree up and pushing the “too short” sub-tree down

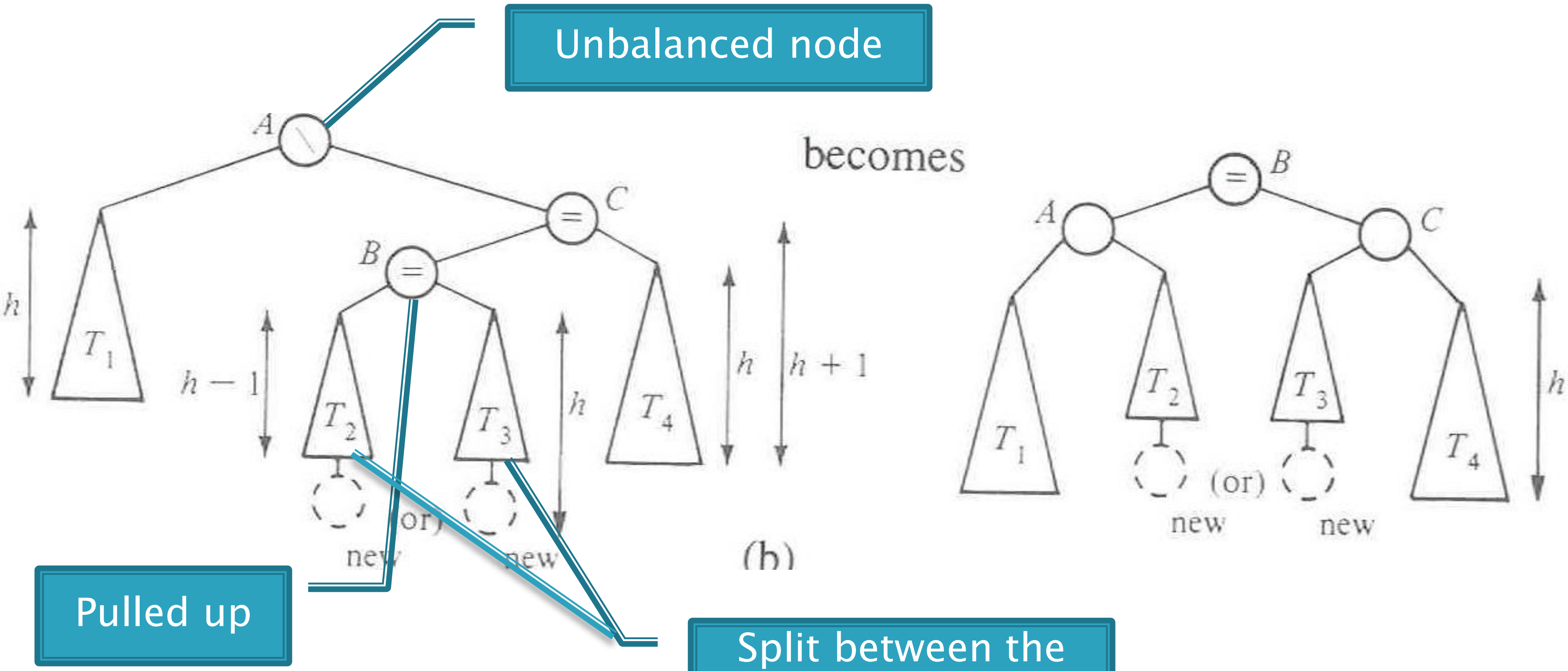
- Two basic cases
  - “See saw” case:
    - Too-tall sub-tree is on the outside
    - So tip the see saw so it’s level
  - “Suck in your gut” case:
    - Too-tall sub-tree is in the middle
    - Pull its root up a level

Single Left Rotation



Diagrams are from *Data Structures* by E.M. Reingold and W.J. Hansen

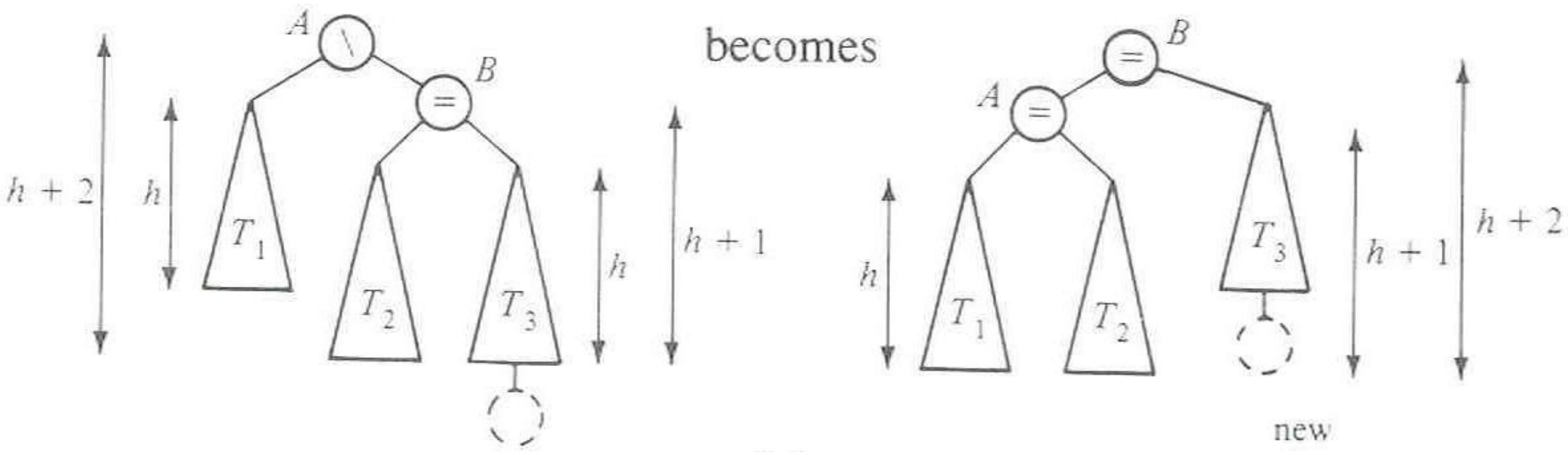
Double Left Rotation



Weiss calls this “right-left double rotation”

Your turn — work with a partner

Q8



- Write the method:
-
```
static BalancedBinaryNode singleRotateLeft (
    BalancedBinaryNode parent,    /* A */
    BalancedBinaryNode child      /* B */   ) {

}
```
- Returns a reference to the new root of this subtree.
- Don’t forget to set the balanceCode fields of the nodes.

More practice—(sometime after class)

- Write the method:

-
  ```
  BalancedBinaryNode doubleRotateRight (
     BalancedBinaryNode parent,      /* A */
     BalancedBinaryNode child,       /* C */
     BalancedBinaryNode grandChild   /* B */ ) {


  }
  ```
- Returns a reference to the new root of this subtree.
- Rotation is mirror image of double rotation from an earlier slide

O(log N)?

- If you have to rotate after insertion, you can stop moving up the tree:
  - Both kinds of rotation leave height the same as before the insertion!

- Is insertion plus rotation cost really O(log N)?

Insertion/deletion in AVL Tree:	O(log n)
Find the imbalance point (if any):	O(log n)
Single or double rotation:	O(1)
(looking ahead) *for deletion, may have to do O(log N) rotations*	
Total work:	O(log n)

Term Project: EditorTrees

Like BST, except:

1. Keep height-balanced
2. Insertion/deletion by **index**, not by comparing elements. So not sorted

Examples:

- EditorTree et = new EditorTree()
- et.add(‘a’) // append to end
- et.add(‘b’) // same
- et.add(‘c’) // same. Rebalance!
- et.add(‘d’, 2) // where does it go?
- et.add(‘e’)
- et.add(‘f’, 3)

- Notice the tree is height-balanced (so height = O(log n)), but not a BST

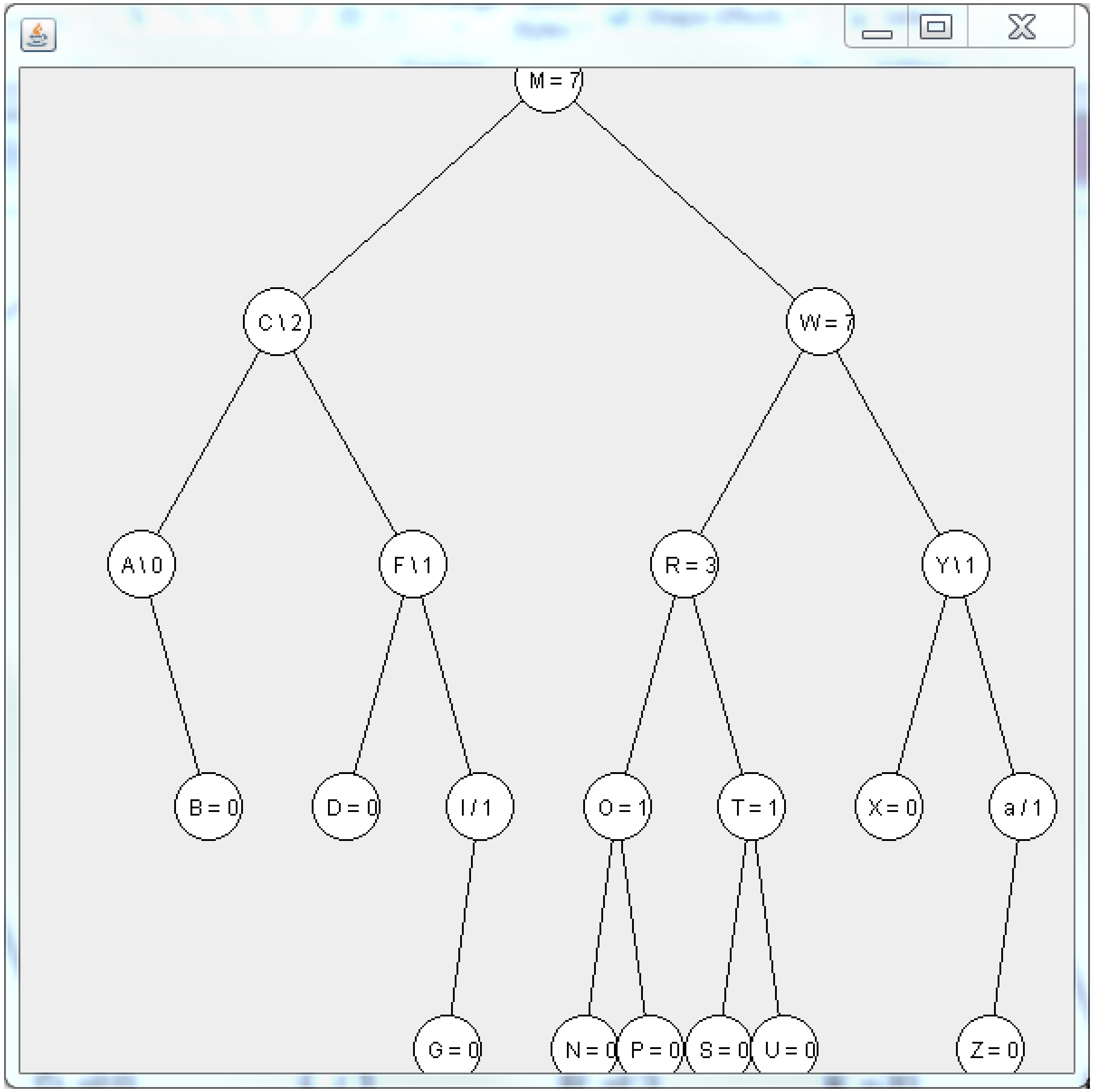
To find index quickly, add a **rank** field to BinaryNode

- Gives the in-order position of this node within its own subtree
  - i.e., the size of its left subtree

0-based indexing

- How would we do `findK`$_{th}$?

- `Insert` and `delete` start similarly


M = 7
C \ 2
W = 7
A \ 0
F \ 1
R = 3
Y \ 1
B = 0
D = 0
I / 1
O = 1
T = 1
X = 0
a / 1
G = 0
N = 0
P = 0
S = 0
U = 0
Z = 0

Get with your EditorTrees team

Read the specification and check out the starting code

Milestone 1 due soon.
Get started before next class!

Test 1 summary:

- Goals
  - Runtime of code with loops, including divide and conquer (cut in half = logs)
    - Big-Oh and cousins
  - Using common ADTs
    - Difference between sets and maps, hash and tree implementations
    - Decisions about which ADT is best to use for a given problem
      - For correctness and efficiency
    - Nice job with PurgeableStack
- Overall a good start!

Test 2a next class: Recursive tree methods all follow this format

- Consider an arbitrary method named foo()

foo()

If base case, return the appropriate value

  - 1. Compute a value for the node
  - 2. Call left.foo()
  - 3. Call right.foo()
  - Combine the results and return them

- This is O(n) if the computation on the node is constant-time
- When searching in a BST, you only need to recurse left **or** right, so it is O(height)

If you submitted HW4, you will receive a solution in your repo.

HW5 is very relevant – I encourage you to start before the test!

Let’s discuss now