

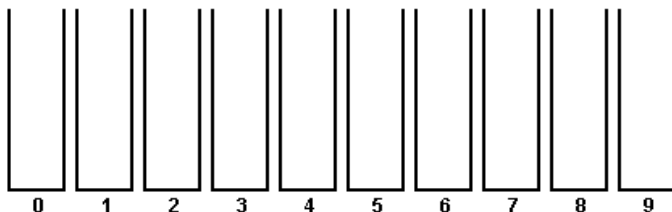
What is the min height of a tree with X external nodes?

CSSE 230

Sorting Lower Bound Radix Sort

Radix sort to the rescue ... sort of...

2805 1081 7556 3684 6428 6321 1456 8850 8022 7919



After today, you should be able to...

- ...explain why comparison-based sorts need at least $O(n \log n)$ time
- ... explain bucket sort
- ... explain radix sort
- ... explain the situations in which radix sort is faster than $O(n \log n)$

Announcements

- ▶ EditorTree evals due last night – late is better than never on these, though!
- ▶ Doublets demo
- ▶ Questions on HW8?
- ▶ Once in a lifetime opportunity
 - ...tomorrow 7th hour in the GM Room.

A Lower-Bound on Sorting Time

We can't do much better than
what we already know how to
do.

What's the best best case?

- ▶ Lower bound for best case?
- ▶ A particular algorithm that achieves this?

What's the best worst case?

- ▶ Want a function $f(N)$ such that the worst case running time for **all sorting algorithms** is $\Omega(f(N))$
- ▶ How do we get a handle on “all sorting algorithms”?

Tricky!

What are “all sorting algorithms”?

- ▶ We can't list all sorting algorithms and analyze all of them
 - Why not?
- ▶ But we can find a **uniform representation** of any sorting algorithm that is based on **comparing** elements of the array to each other

First of all...

- ▶ The problem of sorting N elements is at least as hard as determining their ordering
 - e.g., determining that $a_3 < a_4 < a_1 < a_5 < a_2$
 - sorting = determining order, then movement
- ▶ So any lower bound on all "order-determination" algorithms is also a lower bound on "all sorting algorithms"

Sort Decision Trees

- ▶ Let A be any **comparison-based algorithm** for sorting an array of distinct elements
- ▶ We can draw an EBT that corresponds to the comparisons that will be used by A to sort an array of N elements
 - This is called a **sort decision tree**
 - Internal nodes are comparisons
 - External nodes are orderings
- Different algorithms will have different trees

So what?

- ▶ Minimum number of external nodes in a sort decision tree? (As a function of N)
- ▶ Is this number dependent on the algorithm?
- ▶ What's the height of the shortest EBT with that many external nodes?

$$\lceil \log N! \rceil \approx N \log N - 1.44N = \Omega(N \log N)$$

No comparison-based sorting algorithm, known or not yet discovered, can ever do better than this!

Can we do better than $N \log N$?

- ▶ $\Omega(N \log N)$ is the best we can do if we compare items
- ▶ Can we sort without comparing items?

Yes, we can! We can avoid comparing items and still sort. This is fast if the range of data is small. Q5

- ▶ Observation:
 - For N items, if the range of data is less than N , then we have duplicates
- ▶ $O(N)$ sort: Bucket sort
 - Works if possible values come from limited range
 - Example: Exam grades histogram
- ▶ A variation: Radix sort

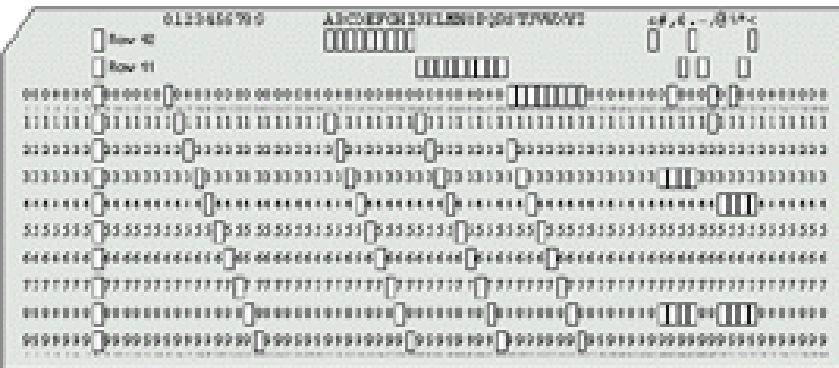
Radix sort

- ▶ A picture is worth 10^3 words, but an animation is worth 2^{10} pictures, so we will look at one.
- ▶ <http://www.cs.auckland.ac.nz/software/AlgAnim/radixsort.html> (good but blocked)
- ▶ https://www.youtube.com/watch?v=xuU-DS_5Z4g&src_vid=4S1L-pyQm7Y&feature=iv&annotation_id=annotation_133993417 (video, good basic idea, distracting zooms)
- ▶ <http://www.cs.usfca.edu/~galles/visualization/RadixSort.html> (good, uses single array)

RadixSort is almost $O(n)$

- ▶ It is $O(kn)$
 - Looking back at the radix sort algorithm, what is k ?
- ▶ Look at some extreme cases:
 - If all integers in range 0–100 (so many duplicates if N is large), then $k = \text{-----}$
 - If all N integers are distinct, $k = \text{----}$

Radix sort example: card sorter



Type 82 Electric Punched Card Sorting Machine

Used an appropriate
combo of
mechanical, digital,
and human effort to
get the job done.

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 - Hear Fred Brooks, Turing Award winner!