

CSSE 220 Day 25

Sorting Algorithms
Algorithm Analysis and Big-O
Function Objects and the Comparator Interface

Checkout *SortingAndSearching* project from SVN

Questions

»» Exam results


Remember Selection Sort?

»» Let's see...

Why study sorting?

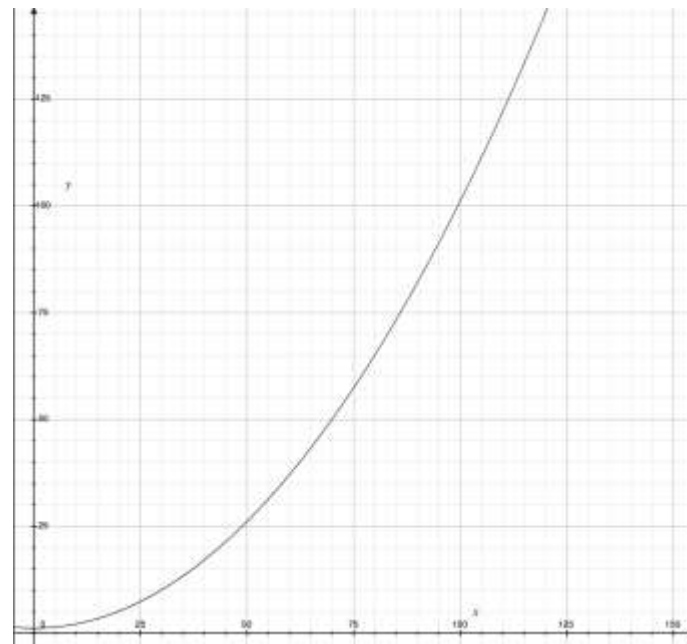
- » Remember
Shlemiel the Painter

Course Goals for Sorting: You should...


- ▶ Be able to **describe** basic sorting algorithms:
 - Selection sort
 - Insertion sort
 - Merge sort
 - Quicksort
 - ▶ Know the **run-time efficiency** of each
 - ▶ Know the **best and worst case** inputs for each
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Profiling Selection Sort

- ▶ **Profiling**: collecting data on the run-time behavior of an algorithm
- ▶ How long does selection sort take on:
 - 10,000 elements?
 - 20,000 elements?
 - ...
 - 80,000 elements?
- ▶ $O(n^2)$

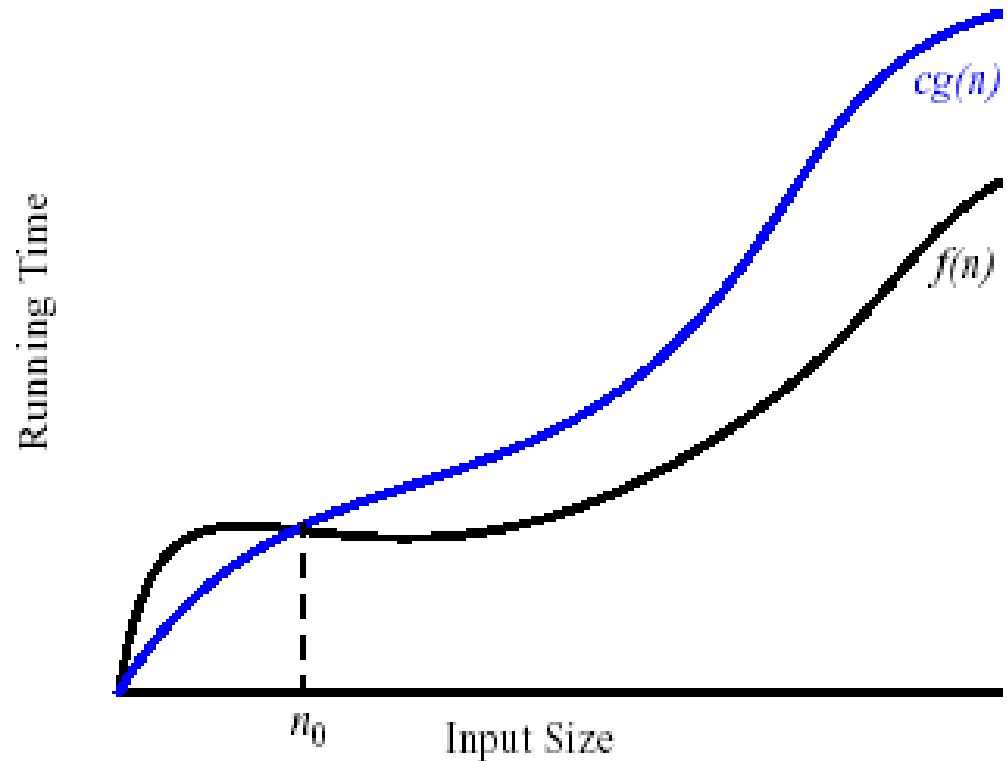


Big-Oh Notation

- ▶ In analysis of algorithms we care about differences between algorithms on very large inputs
 - ▶ We say, “selection sort takes on the order of n^2 steps”
 - ▶ Big-Oh gives a formal definition for “on the order of”
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Formally

- ▶ We write $f(n) = O(g(n))$, and say “ f is big-Oh of g ”
- ▶ if there exists positive constants c and n_0 such that
- ▶ $0 \leq f(n) \leq c g(n)$
for all $n > n_0$
- ▶ g is a **ceiling** on f



Rule of Thumb

- ▶ Suppose the number of operations is given by a polynomial:

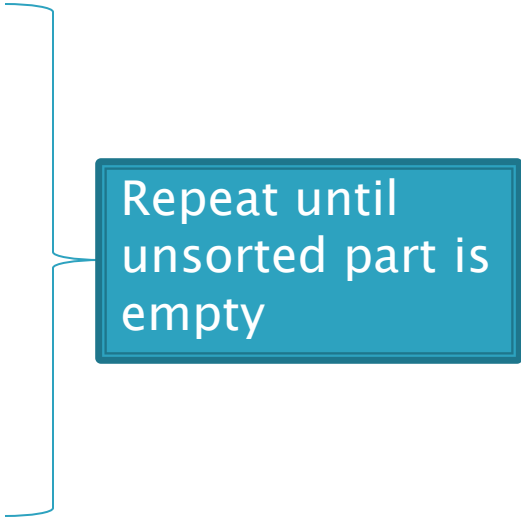
$$a_k * n^k + a_{k-1} * n^{k-1} + \dots + a_2 * n^2 + a_1 * n + a_0$$

- ▶ Then the algorithm is $O(n^k)$.
- ▶ That is, take the highest order term and drop the coefficient

Insertion Sort

▶ Basic idea:

- Think of the list as having a sorted part (at the beginning) and an unsorted part (the rest)
- Get the first number in the unsorted part
- Insert it into the correct location in the sorted part, moving larger values up to make room



Repeat until
unsorted part is
empty

Insertion Sort Exercise, Q4-11b


- ▶ **Profile** insertion sort
- ▶ **Analyze** insertion sort assuming the inner while loop runs that maximum number of times (count the array accesses)
- ▶ What input causes the worst case behavior?
The best case?
- ▶ Does the input affect selection sort?

Ask for help if you're stuck!

Searching

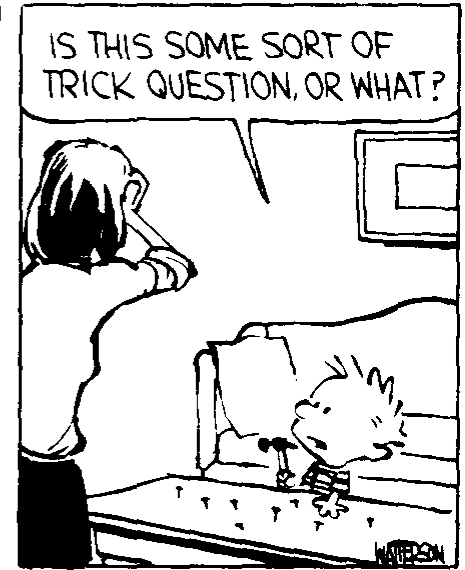
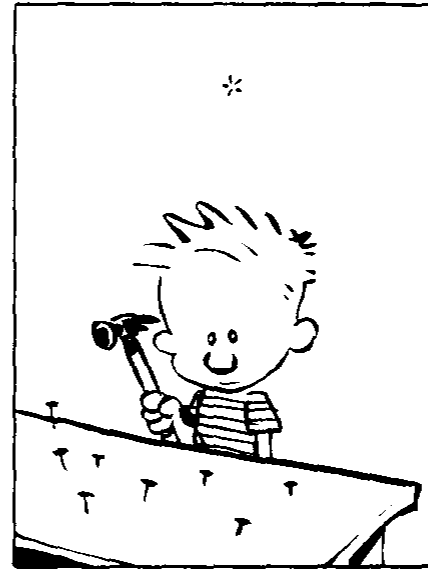
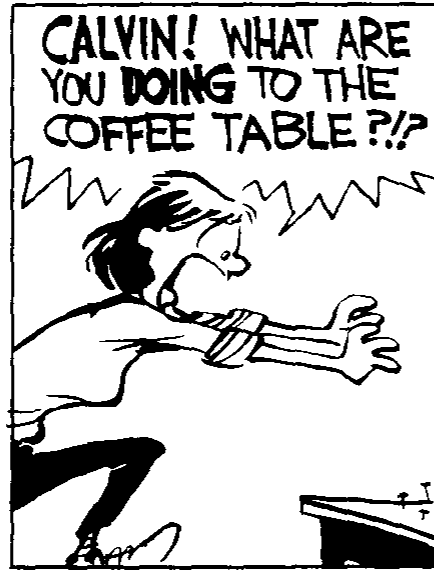
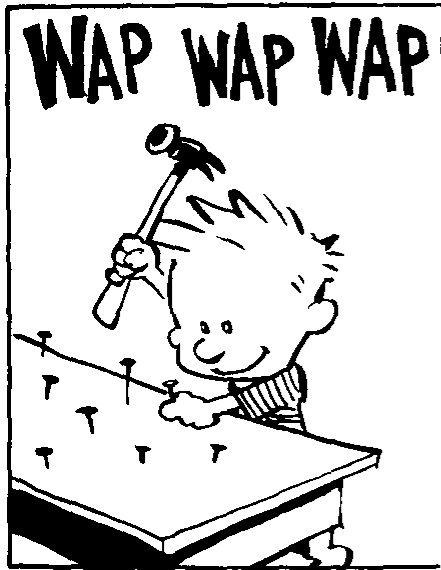
- ▶ For searching unsorted data, what's the worst case number of comparisons we would have to make?

Binary Search of Sorted Data

- ▶ A **divide and conquer** strategy
 - ▶ Basic idea:
 - Divide the list in half
 - Should result be in first or second half?
 - Recursively search that half
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Analyzing Binary Search

- ▶ What's the best case?
- ▶ What's the worst case?



Perhaps it's time for a break.

Merge Sort

- ▶ Basic recursive idea:
 - If list is length 0 or 1, then it's already sorted
 - Otherwise:
 - Divide list into two halves
 - Recursively sort the two halves
 - **Merge** the sorted halves back together
- ▶ Let's profile it...

Analyzing Merge Sort

- ▶ More trees

Quicksort

- ▶ Basic recursive idea:
 - If length is 0 or 1, then it's already sorted
 - Otherwise:
 - Pick a "pivot"
 - Shuffle the items around so all those less than the pivot are to its left and greater are to its right
 - Recursively sort the two "partitions"
- ▶ Let's profile it...




Analyzing Quicksort

- ▶ This one is trickier
- ▶ How should we choose the “pivot”


Function Objects

- » Another way of creating reusable code

A Sort of a Different Order

- ▶ Java libraries provide efficient sorting algorithms
 - `Arrays.sort(...)` and `Collections.sort(...)`
 - ▶ But suppose we want to sort by something other than the “natural order” given by `compareTo()`
 - ▶ Function Objects to the rescue!
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Function Objects

- ▶ Objects defined to just “wrap up” functions so we can pass them to other (library) code
 - ▶ We’ve been using these for awhile now
 - Can you think where?
 - ▶ For sorting we can create a function object that implements Comparator
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Data Structures

- »» Understanding the engineering trade-offs when storing data

Data Structures

- ▶ Efficient ways to store data based on how we'll use it
- ▶ So far we've seen `ArrayLists`
 - Fast addition to end of list
 - Fast access to any existing position
 - Slow inserts to and deletes from middle of list

Another List Data Structure

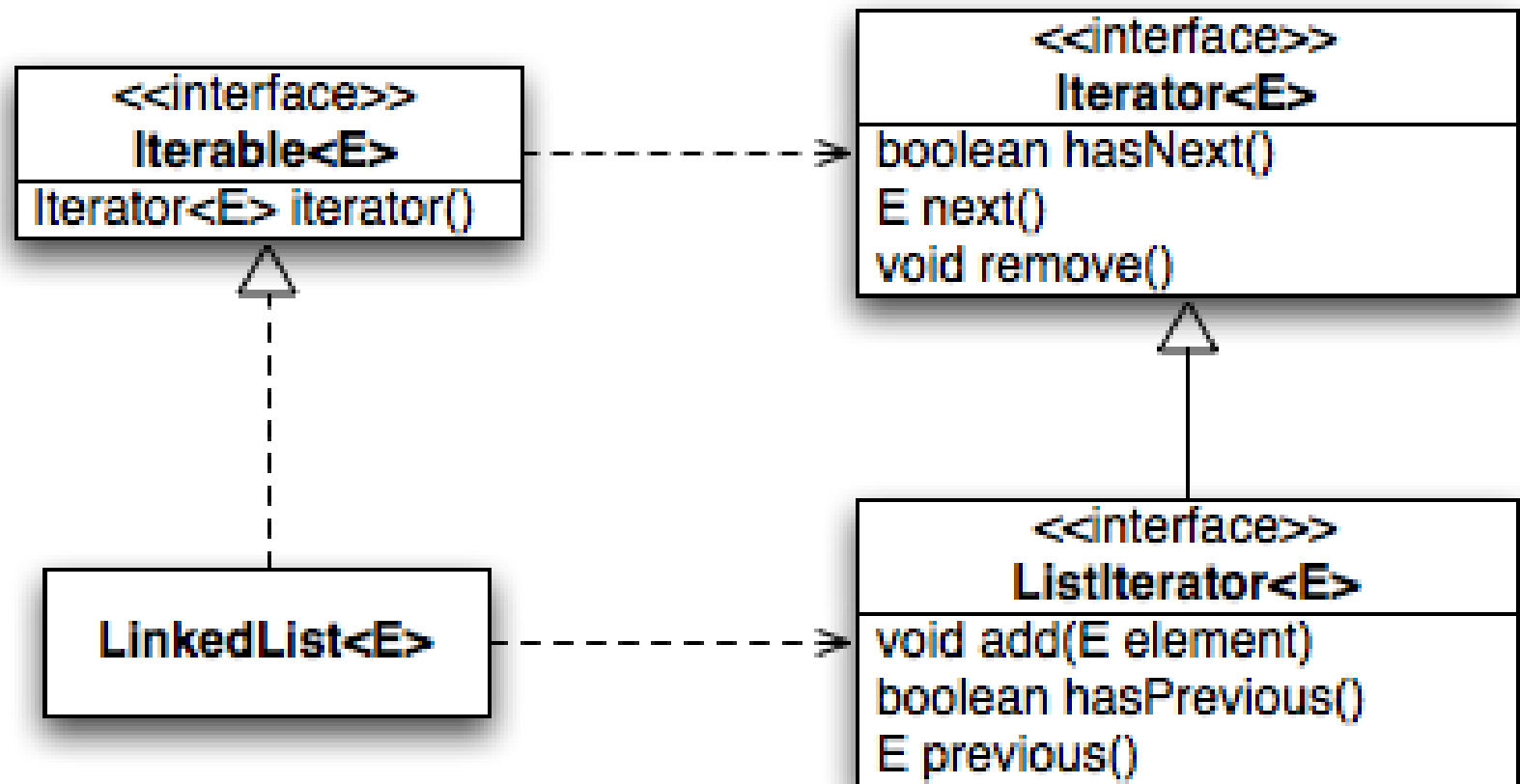
- ▶ What if we have to add/remove data from a list frequently?
- ▶ `LinkedLists` support this:
 - Fast insertion and removal of elements
 - Once we know where they go
 - Slow access to arbitrary elements

LinkedList<E> Methods

- ▶ **void addFirst(E element)**
- ▶ **void addLast(E element)**
- ▶ **E getFirst()**
- ▶ **E getLast()**
- ▶ **E removeFirst()**
- ▶ **E removeLast()**

- ▶ What about the middle of the list?
 - **LinkedList<E> implements Iterable<E>**

Accessing the Middle of a LinkedList



An Insider's View

```
for (String s : list) {  
    // do something  
}
```

```
Iterator<String> iter =  
    list.iterator();
```

```
while (iter.hasNext()) {  
    String s = iter.next();  
    // do something  
}
```

Enhanced For Loop

What Compiler Generates

Next Time

- ▶ Implementing ArrayList and LinkedList
 - ▶ A tour of some data structures
 - ▶ VectorGraphics work time
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