CSSE 220 Day 22

Generics and Comparable Analysis of Algorithms intro Function Objects intro

Exam contents

- Exam will NOT include Chapter 14.
 - Except for the intro to analysis and big-oh which we will cover today.
 - I want to give you more time for the ideas to sink in.
 - Also I want to do a couple of other things before we get to the heart of chapter 14.
- The Computer part of the exam will not ask you to do any GUI programming.
 - There most likely will be GUI programming on the Final exam.
 - Likely things for you to do for Exam 2 Computer part:
 - Algorithms, recursion, classes, interfaces, inheritance, abstract classes, ArrayLists and Arrays.

VectorGraphics and Exam 2

- On the Written part, I may ask something about how your team did some particular aspect of the project
 - As a way of checking to make sure that everyone understands everything you did for the project
- Do you have questions about the exam?

Questions

Vector GraphicsExamRecursionAnything Else

Generic types and Collections

>>> Also Comparable interface

Generic Types and Collections

Before Java 1.5 (still supported, but gives warnings):

```
ArrayList a = new ArrayList(); Explicit class cast
Integer b = new Integer(7); required.
a.add(b);
Integer c = (Integer)(a.get(0));
```

New version (using Java generic type):

automatic unboxing: Integer → int.

Efficiency: Compile time vs run-time checking Q1

required.

Comparable review:

- interface java.lang.Comparable<T>
- ▶ Type Parameters: T the type of objects that this object may be compared to
- int compareTo(T other)
 - Compares this object with the specified object for ordering purposes.
 - Returns a negative integer, zero, or a positive integer as this object is less than, equal to, or greater than the specified object.

compareTo: the fine print

int compareTo (<u>T</u> o) from the JDK API documentation

Compares this object with the specified object for order. Returns a negative integer, zero, or a positive integer as this object is less than, equal to, or greater than the specified object.

The implementor must ensure sgn(x.compareTo(y)) == -sgn(y.compareTo(x)) for all x and y. (This implies that x.compareTo(y) must throw an exception iff y.compareTo(x) throws an exception.)

The implementor must also ensure that the relation is transitive: (x.compareTo(y)>0 && y.compareTo(z) >0) implies x.compareTo(z)>0.

Finally, the implementor must ensure that x.compareTo(y) == 0 implies that sgn(x.compareTo(z)) == sgn(y.compareTo(z)), for all z.

It is strongly recommended, but not strictly required that (x.compareTo(y) == 0) == (x.equals(y)). Generally speaking, any class that implements the Comparable interface and violates this condition should clearly indicate this fact. The recommended language is "Note: this class has a natural ordering that is inconsistent with equals."

In the foregoing description, the notation sgn (expression) designates the mathematical signum function, which is defined to return one of -1, 0, or 1 according to whether the value of expression is negative, zero or positive.

Interface Comparable<T>

Type Parameters:

- T the type of objects that this object may be compared to
- Any class that implements Comparable contracts to provide a compareTo() method

Method Detail

compareTo

int compareTo(T o)

String is a Comparable class.

If it did not already have a compareTo()
method, how would you write it?

Compares this object with the specified object for order. Returns a negative integer, zero, or a positive integer as this object is less than, equal to, or greater than the specified object.

Therefore, we can write generic methods on Comparable objects. For example, in the java.util.Arrays class:

```
Sort (Object [] a, int fromIndex, int toIndex)

Sorts the specified range of the specified array of objects into ascending order, according to the natural ordering of its elements.
```

Example of using Arrays.sort

```
import java.util.Arrays;
public class StringSort {
   public static void main(String[] args) {
      String [] toons = {"Mickey", "Minnie", "Donald",
                          "Pluto", "Goofy"};
      Arrays.sort(toons);
      for (String s:toons)
                                     Output:
         System.out.println(s);
                                     Donald
        Collections.sort can
                                     Goofy
        similarly be used to sort
                                     Mickey
        ArrayLists and other
                                     Minnie
        Collection objects.
                                     Pluto
```

Measuring program efficiency

General hints on efficiency ExamplesBig-oh and its cousins

Measuring program effciency

- What kinds of things should we measure?
 - CPU time
 - memory used
 - disk transfers
 - network bandwidth
- Mostly in this course, we focus on the first two, and especially on CPU time
- To measure running time, we can call System.currentTimeMillis()

Program Efficiency, part 2

- Some simple efficiency tips
 - If a statement in a loop calculates the same value each time through, move it outside (usually before) the loop
 - Store and retain data on a "need to know" basis
 - Don't store values that you won't reuse
 - Do store values that you need to reuse
 - Don't put everything into an array when you only need one or two consecutive items at a time
 - Don't declare a variable as a field if it can be a local variable of a method

Familiar example:

Linear search of a sorted array of Comparable items

```
for (int i=0; i < a.length; i++)
  if ( a[i].compareTo(soughtItem) > 0 )
    return NOT_FOUND; // perhaps NOT_FOUND == -1
  else if ( a[i].compareTo(soughtItem) == 0 )
    return i;
return NOT_FOUND;
```

- •What should we count?
- •Best case, worst case, average case?

Another algorithm analysis example

Does the following method actually create and return a copy of the string s?

What can we say about the running time of the method? (where N is the length of the string s)

What should we count?

```
public static String stringCopy(String s) {
   String result = "";
   for (int i=0; i<s.length(); i++)
      result += s.charAt(i);
   return result;</pre>
```

Don't be too quick to make assumptions when analyzing an algorithm!

How can we do the copy more efficiently?

Break



Interlude

Always code as if the guy who ends up maintaining your code will be a violent psychopath who knows where you live.

--Martin Golding

Figure 5.1
Running times for small inputs

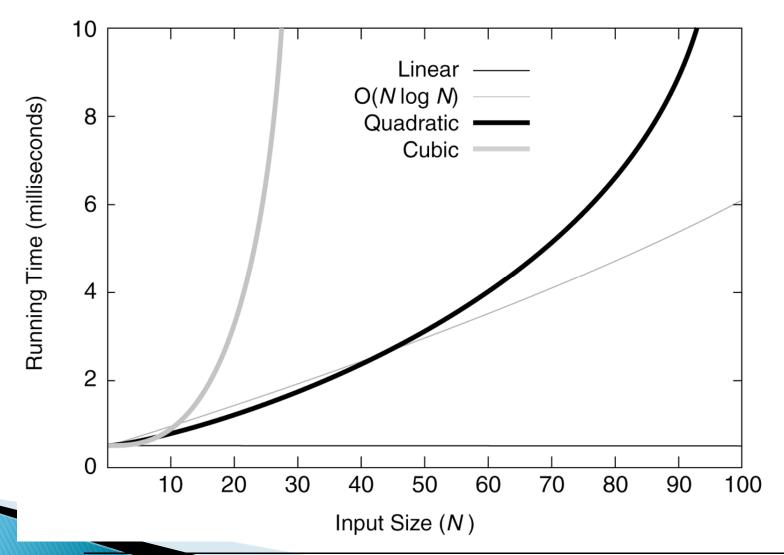


Figure 5.2
Running times for moderate inputs

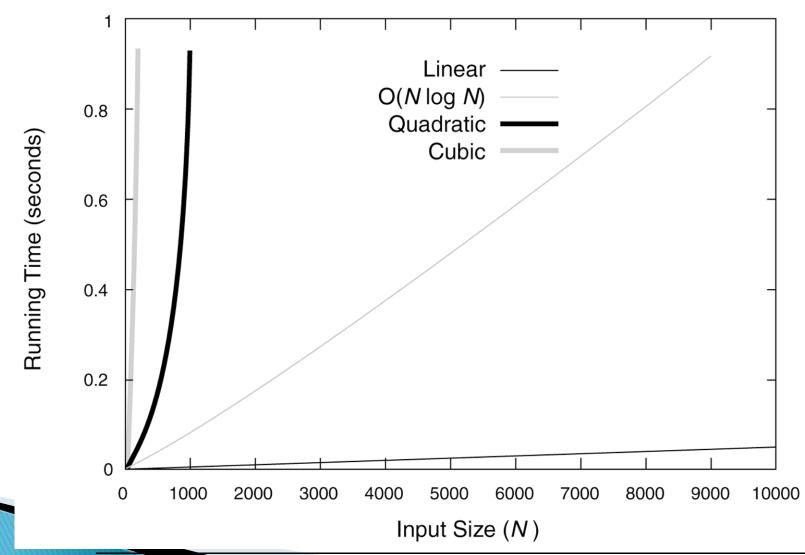


Figure 5.3 Functions in order of increasing growth rate

Function	Name
С	Constant
$\log N$	Logarithmic
$\log^2 N$	Log-squared
N	Linear
$N \log N$	N log N ——— a.k.a "log linear"
N^{2}	Quadratic
N^3	Cubic
2^N	Exponential

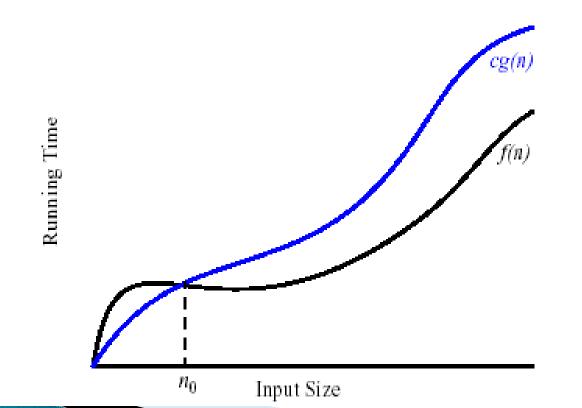
Asymptotic analysis

- We only really care what happens when N (the size of a problem) gets large
- Is the function basically linear, quadratic, etc.?
- ▶ For example, when n is large, the difference between n² and n² – 3 is negligible

The "Big-Oh" Notation

- given functions f(n) and g(n), we say that f(n) is O(g(n)) if and only if $f(n) \le c g(n)$ for $n \ge n_0$
- c and n_0 are constants, f(n) and g(n) are functions over non-negative integers

In this course, we won't be so formal. We'll just say that f(N) is O(g(N)) means that f(n) is eventually smaller than a constant times g(n).



- Simple Rule: Drop lower order terms and constant factors.
 - 7n 3 is O(n)
 - $8n^2\log n + 5n^2 + n$ is $O(n^2\log n)$

- Special classes of algorithms:
 - logarithmic: $O(\log n)$
 - linear O(n)
 - quadratic $O(n^2)$
 - polynomial $O(n^k)$, $k \ge 1$
 - exponential $O(a^n)$, n > 1

- · "Relatives" of the Big-Oh
 - $-\Omega(f(n))$: Big Omega
 - $-\Theta(f(n))$: Big Theta

Recap: O, Ω, Θ

- f(N) is O(g(N)) if there is a constant c such that for sufficiently large N, $f(N) \le cg(N)$
 - Informally, as N gets large the growth rate of f is bounded above by the growth rate of g
- f(N) is $\Omega(g(N))$ if there is a constant c such that for sufficiently large N, $f(N) \ge cg(N)$
 - Informally, as N gets large the growth rate of f is bounded below by the growth rate of g
- f(N) is $\Theta(g(N))$ if f(N) is O(g(n)) and f(N) is $\Omega(g(N))$
 - Informally, as N gets large the growth rate of f is the same as the growth rate of g

Limits and asymptotics

consider the limit

$$\lim_{n \to \infty} \frac{f(n)}{g(n)}$$

- What does it say about asymptotics if this limit is zero, nonzero, infinite?
- We could say that knowing the limit is a sufficient but not necessary condition for recognizing big-oh relationships.
- It will be all we need for all examples in this course.

Apply this limit property to the following pairs of functions

- 1. N and N^2
- 2. $N^2 + 3N + 2$ and N^2
- 3. $N + \sin(N)$ and N
- 4. log N and N
- 5. N log N and N^2
- 6. Na and Nn
- 7. a^N and b^N (a < b)
- 8. $\log_a N$ and $\log_b N$ (a < b)
- 9. N! and N^N

Big-Oh Style

Give tightest bound you can

• Saying that 3N+2 is $O(N^3)$ is true, but not as useful as saying it's O(N) [What about $O(N^3)$?]

Simplify:

- You could say:
- \circ 3n+2 is O(5n-3log(n) + 17)
- and it would be technically correct...
- It would also be poor taste ... and put me in a bad mood.

▶ But... if I ask "true or false: 3n+2 is O(n³)", what's the answer?

- True!
- There may be "trick" questions like this on assignments and exams.
- But they aren't really tricks, just following the big-Ohderinition!

Work Time

VectorGraphics cycle.
Finish before next class meeting.
Get help as needed.