

CSSE132

Introduction to Computer Systems

8 : Boolean Algebra

March 7, 2013

- Postulates
- Theorems
- DeMorgan's Theorem
- Some Definitions
- Canonical Forms
- Canonical Forms of the Half Adder
- Complements and Conversions
- In-class Examples
- Basic Gates

Boolean Algebra: Postulates

P1: $A = 0$ if $A \neq 1$ $A=1$ if $A \neq 0$

P2: if $A = 0$ then $A'=1$ if $A = 1$ then $A'=0$

P3: $0*0=0$ $1+1=1$

P4: $1*1=1$ $0+0=0$

P5: $0*1=1*0=0$ $1+0=0+1=1$

Boolean Algebra: Theorems

T1: $A+0=A$

$A*1=A$

T7: $A+(B+C)=(A+B)+C$

$A*(B*C)=(A*B)*C$

T2: $A+1=1$

$A*0=0$

T8: $A*B+A*C=A*(B+C)$

T3: $A+A=A$

$A*A=A$

$(A+B)*(A+C)=A+B*C$

T4: $(A')'=A$

T9: $A+A*B=A$ $A*(A+B)=A$

T5: $A+A'=1$

$A*A'=0$

T10: $A*B+A*B'=A$

$(A+B)*(A+B')=A$

T6: $A+B=B+A$

$A*B=B*A$

T11: Demorgan's theorem

$(A+B)'=A'*B'$

$(A*B)'=A'+B'$

Proof of T10 by Perfect Induction

- Enumerate all combinations and show two expressions are identical.

$$\text{T10: } A * B + A * B' = A$$

$$(A + B) * (A + B') = A$$

Proof of T10 by Perfect Induction

- Enumerate all combinations and show two expressions are identical.

$$\text{T10: } A*B+A*B'=A$$

$$(A+B)*(A+B')=A$$

A	B	A+B	A+B'	(A+B)(A+B')	AB	AB'	AB+AB'
0	0	0	1	0	0	0	0
0	1	1	0	0	0	0	0
1	0	1	1	1	0	1	1
1	1	1	1	1	1	0	1

DeMorgan's Theorem

- Prove by perfect induction
- Enumerate all combinations and show two expressions are identical.

$$(A+B)' = A'B'$$
$$(AB)' = A' + B'$$

$$\overline{(A+B)} = \bar{A} \cdot \bar{B}$$
$$\overline{A \cdot B} = \bar{A} + \bar{B}$$

DeMorgan's Theorem

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$$\overline{(A+B)} = \bar{A} \cdot \bar{B}$$

$$\overline{A \cdot B} = \bar{A} + \bar{B}$$

A	B	A'	B'	A+B	(A+B)'	A'B'	(AB)'	A'+B'
0	0	1	1	0	1	1	1	1
0	1	1	0	1	0	0	1	1
1	0	0	1	1	0	0	1	1
1	1	0	0	1	0	0	0	0

- Result : invert components, change operation!

Some Definitions

Product term:	all variables are ANDed, $A*B*C'*D$
Sum term:	all variables are ORed, $A+B+C'+D$
Sum of products:	$A*B+A*C*D$
Product of sums:	$(A+B)*(A+C)*(B+C+D)$
Normal term:	a product or sum in which each variable appears no more than once
Minterm:	a normal product term containing all variables
Maxterm:	a normal sum term containing all variables

Canonical Forms

- Form one: sum of minterms
- Form two: product of maxterms
- Minterm number and Maxterm number

A	B	C	Minterm	Minterm number	Maxterm	Maxterm number
0	0	0	$A'B'C'$	$\Sigma(0)$	$A+B+C$	$\Pi(0)$
0	0	1	$A'B'C$	$\Sigma(1)$	$A+B+C'$	$\Pi(1)$
0	1	0	$A'BC'$	$\Sigma(2)$	$A+B'+C$	$\Pi(2)$
0	1	1	$A'BC$	$\Sigma(3)$	$A+B'+C'$	$\Pi(3)$
1	0	0	$AB'C'$	$\Sigma(4)$	$A'+B+C$	$\Pi(4)$
1	0	1	$AB'C$	$\Sigma(5)$	$A'+B+C'$	$\Pi(5)$
1	1	0	ABC'	$\Sigma(6)$	$A'+B'+C$	$\Pi(6)$
1	1	1	$A*B*C$	$\Sigma(7)$	$A'+B'+C'$	$\Pi(7)$

Example: Canonical Forms of the Half Adder

A	B	Sum	Carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

$$\text{Sum} = A'B + AB' = \Sigma(1,2)$$

$$\text{Sum} = (A'B' + AB)' = (A+B)(A'+B') = \Pi(3,0)$$

$$\text{Carry} = AB = \Sigma(3)$$

$$\text{Carry} = (A'B' + A'B + AB)' = \Pi(0,1,2)$$

Complements and Conversions

- To complement a function, replace terms with those that are not present.
 - Complement of a function consists of all terms that causes “0”.
- To convert a function from the product form to the sum form, change the product symbol to the sum symbol and use the terms that are not present, vice versa.
 - Again, all those new terms form “0” entries in the truth table.

The Half Adder: Complement and Conversion

A	B	Sum	Carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

Complement of Sum:

$$Sum = \bar{A}B + A\bar{B} = \sum (1,2)$$

$$\overline{Sum} = \overline{\bar{A}B + A\bar{B}} = \sum (0,3)$$

Convert Sum from SOP to POS:

$$\begin{aligned} Sum &= \bar{A}B + A\bar{B} = \sum (1,2) \\ &= \prod (0,3) = (A + B)(\bar{A} + \bar{B}) \end{aligned}$$

Boolean Algebra Properties

➤ Commutative

- $a + b = b + a$
- $a * b = b * a$

➤ Distributive

- $a * (b + c) = a * b + a * c$
- $a + (b * c) = (a + b) * (a + c)$

➤ Associative

- $(a + b) + c = a + (b + c)$
- $(a * b) * c = a * (b * c)$

➤ Identity

- $0 + a = a + 0 = a$
- $1 * a = a * 1 = a$

➤ Complement

- $a + a' = 1$
- $a * a' = 0$

➤ To prove, just evaluate all possibilities.

Example Applications of Boolean Algebra Properties

➤ Show abc' equivalent to $c'ba$.

- Use commutative property:
 $a*b*c' = a*c'*b = c'*a*b = c'*b*a$

➤ Show $abc + abc' = ab$.

- Use first distributive property
 $abc + abc' = ab(c+c')$.
- Complement property
Replace $c+c'$ by 1: $ab(c+c') = ab(1)$.
- Identity property
 $ab(1) = ab*1 = ab$.

➤ Show $x + x'z$ equivalent to $x + z$.

- Second distributive property
Replace $x+x'z$ by $(x+x')(x+z)$.
- Complement property
Replace $(x+x')$ by 1,
- Identity property
replace $1*(x+z)$ by $x+z$.

- Simplify the expression with Boolean algebra and indicate which theorems are used.***

$$Z(A, B, C) = A \cdot B \cdot \bar{C} + \bar{A} \cdot B + \bar{A} \cdot \bar{B} \cdot C$$

- Complement the following expression, using DeMorgan's theorem to produce a product of sums expression.***

$$Z(A, B, C) = A \cdot B \cdot \bar{C} + A \cdot C$$

- Obtain the simplified logic expression for the majority voting function.***

A	B	C	F
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

Solution for Problem 1 of In-class Exercise

- **Simplify with Boolean algebra and indicate which theorems are used.**

$$Z(A,B,C) = A \cdot B \cdot \bar{C} + \bar{A} \cdot B + \bar{A} \cdot \bar{B} \cdot C$$

$$= A \cdot B \cdot \bar{C} + \bar{A} \cdot B \cdot (C + \bar{C}) + \bar{A} \cdot \bar{B} \cdot C \quad \text{T5, T1}$$

$$= A \cdot B \cdot \bar{C} + \bar{A} \cdot B \cdot C + \bar{A} \cdot B \cdot \bar{C} + \bar{A} \cdot \bar{B} \cdot C \quad \text{T8}$$

$$= (A + \bar{A})B \cdot \bar{C} + \bar{A} \cdot C(B + \bar{B}) \quad \text{T8}$$

$$= B \cdot \bar{C} + \bar{A} \cdot C \quad \text{T10}$$

Solution for Problem 2 of In-class Exercise

- **Complement the following, using DeMorgan's theorem to produce a product of sums expression.**

$$Z(A,B,C) = A \cdot B \cdot \bar{C} + A \cdot C$$

$$\bar{Z}(A,B,C) = \overline{A \cdot B \cdot \bar{C} + A \cdot C} = \overline{A \cdot B \cdot \bar{C}} \cdot \overline{A \cdot C}$$

$$= (\bar{A} + \bar{B} + C)(\bar{A} + \bar{C})$$

$$Z(A,B,C) = \overline{(\bar{A} + \bar{B} + C)(\bar{A} + \bar{C})}$$

Solution for Problem 3 of In-class Exercise

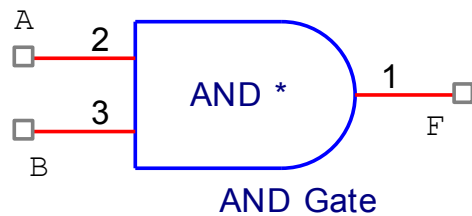
- **Obtain simplified logic expression for the majority voting function.**

$$\begin{aligned} & \bar{A} \cdot B \cdot C + A \cdot \bar{B} \cdot C + A \cdot B \cdot \bar{C} + A \cdot B \cdot C \\ &= \bar{A} \cdot B \cdot C + A \cdot B \cdot C + A \cdot \bar{B} \cdot C + A \cdot B \cdot \bar{C} + A \cdot B \cdot C + A \cdot B \cdot C \quad \text{T3} \\ &= BC + AC + AB \quad \text{T8, T5, T1} \end{aligned}$$

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

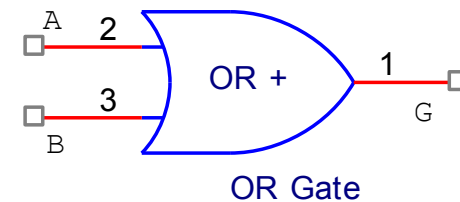
Basic Gates: AND Gate, OR Gate and NOT Gate (Inverter)

A	B	F
0	0	0
0	1	0
1	0	0
1	1	1



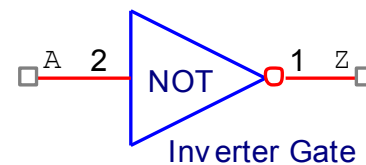
$$F = A * B$$

A	B	G
0	0	0
0	1	1
1	0	1
1	1	1



$$F = A + B$$

A	Z
0	1
1	0



$$Z = A' = \bar{A}$$

➤ Combine AND, OR, NOT to form new gates

- NAND : NOT(AND)
- NOR : NOT(OR)
- XOR : $A \oplus B = A * B' + A' * B$
- XNOR : NOT(XOR)

➤ NAND and NOR can be built without AND or OR gates

- All gates can be implemented with NAND and NOR
- Called 'Universal logic gate'