## CSSE 304 Assignment 6 Updated for Fall, 2019-20

## **Objectives:** You should learn

- To begin to use first-class procedures effectively.
- To think a bit about efficiency of programs.

Before doing A6, be sure to read EoPL-1 from page 22 to the end of chapter 1.

#### Details for these instructions are in the previous assignment

Individual assignment. Comments at beginning, before each problem, when you do anything non-obvious. Submit to server (test offline first). Mutation not allowed.

Assignments 6a (problems 1-7) and 6b (problems 8-12) are due on different days.

**Reading Assignment: See the schedule page.** Have you been keeping up with the reading?

The first four problems are different than problems on previous assignments, in that I want you to read the EoPL-1 material and do related problems without first discussing the material in class. Of course, you are allowed to get help from other students, from the TAs, and from me.

## Problems to turn in: For many of these, you will want to write one or more helper procedures.

Some of the problems deal with *currying*. http://en.wikipedia.org/wiki/Currying describes this as:

In mathematics and computer science, currying (schönfinkeling) is the technique of transforming a function that takes multiple arguments (or a tuple of arguments) in such a way that it can be called as a chain of functions, each with a single argument (partial application). It was originated by Moses Schönfinkel and later worked out by Haskell Curry.

**Optional, not required knowledge for this course:** An interesting discussion of the advantages of currying (the language of discourse is Haskell, but I think you can still follow much of the discussion): <a href="http://www.reddit.com/r/programming/comments/181y2a/what is the advantage of currying/">http://www.reddit.com/r/programming/comments/181y2a/what is the advantage of currying/</a>

Some simple examples of currying appear on pages 26 (last sentence) through 28 of EoPL-1. The first two turnin-problems are from that section, and I recommend that you also think about problem 1.3.6.

**Reminder:** EoPL-1 is the first edition of EoPL. An excerpt was handed out on the first day of class. It is also available on Moodle.

#1 (10 points) curry2. This is EoPL-1 Exercise 1.3.4, page 28. Examples are on that page.

```
\textbf{curry2:} \ ((SchemeObject \times SchemeObject) \boldsymbol{\rightarrow} SchemeObject) \boldsymbol{\rightarrow} (SchemeObject \boldsymbol{\rightarrow} (SchemeObject) \boldsymbol{\rightarrow} SchemeObject))
```

```
#2 (10 points) EoPL-1 Exercise 1.3.5, page 28. Call your procedure curried-compose. For example, (((curried-compose car) cdr) '(a b c)) > b
```

**curried-compose:** (SchemeObject  $\rightarrow$  SchemeObject)  $\rightarrow$  ((SchemeObject)  $\rightarrow$  (SchemeObject)  $\rightarrow$  (SchemeObject)

```
#3 (10 points) compose. EoPL-1 Exercise 1.3.7, page 29. This one will most likely begin (define compose (lambda list-of-functions); notice the lack of parentheses around the argument name.
```

**compose**: listOf (SchemeObject  $\rightarrow$  SchemeObject)  $\rightarrow$  (SchemeObject  $\rightarrow$  SchemeObject)

#### **Examples:**

```
((compose list list) 'abc) → ((abc)) 
((compose car cdr cdr) '(a b c d)) → c
```

#4 (10 points) Write the procedure make-list-c that is a curried version of make-list. (Note that the original make-list is described in TSPL Exercise 2.8.3

 $make-list-c: Integer \rightarrow (SchemeObject \rightarrow Listof(SchemeObject))$ 

#### **Examples**:

#5 (10 points) Write (reverse-it lst) that takes a single-level list lst and returns a new list that has the elements of lst in reverse order. The original list should not be changed. Can you do this in O(n) time? [this is not a requirement]. You probably cannot do this if you use append. Obviously, you should not use Scheme's reverse procedure in your implementation.

#6 (10 points) Write (map-by-position fn-list arg-list) where

- fn-list and arg-list have the same length,
- each element of the list fn-list is a procedure that can take one argument,
- and each element of the list arg-list has a type that is appropriate to be an argument of the corresponding element of fn-list.

map-by-position returns a list (in their original order) of the results of applying each function from fn-list to the corresponding value from arg-list. You must use map to solve this problem; no explicit recursion is allowed.

```
(map-by-position (list cadr - length (lambda(x)(- x 3)))
 '((1 \ 2) \ -2 \ (3 \ 4) \ 5)) \Rightarrow (2 2 2 2)
```

#7 (45 points) Examples are in the test cases. A Binary Search Tree (BST) datatype is defined on page 10 of EoPL. Note that, as the book describes, the grammar alone is not sufficient to completely describe what constitutes a BST. We also need the "smaller values in left subtree, larger in right subtree" property, which cannot be described using a context-free grammar. Write the following procedures. Remember, no mutation of lists is allowed! If you cheat by representing a BST as a single-level sorted list or some other simpler structure, you will get 0 points, even if your code passes a lot of the test cases.

- 1. (empty-BST) takes no arguments and creates an empty tree, which is represented by an empty list.
- 2. (empty-BST? obj) takes a Scheme object obj. It returns #t if obj is an empty BST and #f otherwise.
- 3. (BST-insert num bst) returns a BST result. If num is already in bst, result is structurally equivalent to bst. If num is not already in bst, result adds num in its proper place relative to the other nodes in a tree whose shape is the same as the original. Like any BST insertion, this should have a worst-case running time that is O(height(bst)).
- 4. (BST-inorder bst) (can you do it in O(N) time?) produces an ordered list of the values in bst.
- 5. (BST? obj) returns #t if Scheme object obj is a BST and #f otherwise.
- 6. BST-element, BST-left, BST-right. Accessor procedures for the parts of a node.
- 7. (BST-insert-nodes bst nums) starts with tree bst and inserts each integer from the list nums, in the given order, returning the tree that includes all of the inserted nodes. The original tree is not changed, and this procedure does no mutation.
- 8. (BST-contains? bst num) determines, in time that is O(height(bst)), whether num is in bst.
- 9. (BST-height bst) returns the height of the BST bst. Examples:

```
(BST-height '()) \rightarrow -1
(BST-height '(3 () ())) \rightarrow 0
(BST-height '(2 ()(6 (4 () ()) ())) \rightarrow 2
```

#8 (10 points) Write let->application which takes a let expression (represented as a list) and returns the equivalent expression, also represented as a list, that represents an application of a procedure created by a lambda expression. Your solution should not change the *body* of the let expression. This procedure's output list replaces only the top-level let by an equivalent application of a lambda expression. You do not have to find and replace any non-top-level lets. You may assume that the let expression has the proper form; your procedure does not have to check for this. Furthermore, you may assume that the let expression is *not* a named let.

08/28/19

**let->application**:  $SchemeCode \rightarrow SchemeCode$ 

#### Example:

```
(let->application '(let ((x 4) (y 3))
```

```
(let ((z 5))

(+ x (+ y z)))) \Rightarrow

((lambda (x y)

(let ((z 5))

(+ x (+ y z))))

4 3)
```

#9 (10 points) Write let\*->let which takes a let\* expression (represented as a list) and returns the equivalent nested let expression. This procedure replaces only the **top-level** let\* by an equivalent nested let expression. You may assume that the let\* expression has the proper form.

**let\*->let:** SchemeCode → SchemeCode

**Example:** 

**#10** (20 points) (qsort pred ls) is a Scheme procedure that you will write whose arguments are a predicate (total ordering) which takes two arguments x and y, and returns #t if x is "less than" y, #f otherwise. a list whose items can be compared using this predicate.

qsort should produce the sorted list using a QuickSort algorithm (write your own; do not use Scheme's sort procedure).

## For example:

```
(qsort <= '(4 2 4 3 2 4 1 8 2 1 3 4)) → (1 1 2 2 2 3 3 4 4 4 4 8)

(qsort (lambda (x y) (<= (abs (- x 10)) (abs (- y 10))))

'(5 1 10 8 16 17 23 -1))

→ (10 8 5 16 17 1 -1 23)
```

If you do not remember how QuickSort works, see <a href="http://en.wikipedia.org/wiki/Quicksort">http://en.wikipedia.org/wiki/Quicksort</a> or Chapter 7 of the Weiss book used for CSSE230. There are quicksort algorithms that do fancy things when choosing the pivot in order to attempt to avoid the worst case. You do not need to do any of those things here; you can simply use the car of the list as the pivot. Since mutation is not allowed, your algorithm cannot do the sort in-place. Furthermore, you are not allowed to copy the list elements to a vector, then sort the vector and copy back to a list. All of your work should be done with lists.

#11 (15 points) Write a Scheme procedure (sort-list-of-symbols los) which takes a list of symbols and returns a list of the same symbols sorted as if they were strings. You will probably find the following procedures to be useful: symbol->string, map, string<?, sort (look it up in the <a href="Chez Scheme Users">Chez Scheme Users</a> 'Guide). Note that we have not covered specifics related to this problem, It is time for you to read some documentation and figure out how to use things.

**sort-list-of-symbols:**  $ListOf(Symbol) \rightarrow ListOf(Symbol)$ 

```
Example (sort-list-of-symbols '(b c d g ab f b r m)) → (ab b b c d f g m r)
```

#12 (10 points) (ribassoc s los v fail-value) returns the value in vector v that is associated with symbol s (or returns fail-value if there is no such associated value. What do we mean by "associated value"? If the first occurrence of s in los is in position n, return the value that is in position n in v. If s is not a member of los, then there is no associated value, so we return fail-value. You may assume that los and v have the same length; you do not have to test for this.

**ribassoc:** Symbol  $\times$  List-of-symbols  $\times$  Vector-of-values  $\times$  Alternate-answer  $\rightarrow$  value

#### **Examples:**

```
(ribassoc 'c '(a b c d) '\#(1 2 3 4) 'oops) \Rightarrow 3 (ribassoc 'e '(a b c d) '\#(1 2 3 4) 'oops) \Rightarrow oops
```

# Previous questions and answers on Piazza

# Map-by-position

I'm not sure I get what the procedure is supposed to do.

Is it supposed to apply each fn-list to each object in the arg-list? And does it do it in that order?

Or does it apply the corresponding function to the corresponding object?

I thought that's what it was, but I don't see which elements are corresponding since the example one gives functions that don't apply to the argument it's matched up with.

the instructors' answer, where instructors collectively construct a single answer It is your second case, corresponding.
Look more closely. The procedures and arguments in my example DO correspond. Just as in the example from the beginning of yesterday's class, <b>list</b> is not one of the procedures in the list.
-good answer0
Updated 5 months ago by Claude Anderson
followup discussions
I think I see my mistake now, is "list" not one of the functions? That would be my problem
Reply to this followup discussion
Resolved Unresolved
2
Claude Anderson 5 months ago

## qsort: odd error 'incorrect argument count in call'

In the partition method I wrote for qsort, I've been getting an odd error that tells me there may be an incorrect argument count in call "(partition pred? (cdr ls) (list (car ls)))". It's the same line of code every time. But that method is defined right under qsort, as having 3 arguments "(define partition

(lambda (pred? Is pivot)" Oddly enough, if I cut the code out, load it, then put the code back in, it will load perfectly fine without errors and work correctly. However, I'm getting this error when uploading to the server now and have no idea how to fix it. The cases that fail on the server work fine locally, they just give that error on the server.

hw7		
	-good question0	

## the instructors' answer,

list just makes a list of the functions

where instructors collectively construct a single answer

There is a link on the Schedule page (Day 1, resources column) called "What if the grading program gives zero points for something that works on your computer?" Follow it and read the document.

I believe that this document describes your problem and tells you how to fix it.

The built-in procedure that you are redefining is "partition".

## Sort-list-of-symbols: Clarification

So to clarify, we aren't allowed to use any sort of recursion or looping for this, or are we allowed to use recursion with map?

#### the instructors' answer,

where instructors collectively construct a single answer

You may use map. The function that you are mapping over the list can (I think must) be one that you write, but it cannot be call itself or call another procedure that calls it.

## Let->let\*: getting let to not combine

Im working on the let\*->let problem and i made a helper function that take in the chunk of the passed list that holds the various let expressions (i.e '((a 1) (b 2) (c 3))) and is supposed to put a let in form of each expression (i.e. create '(let (a 1) (let (b 2) (let (c 3))))) but no matter what combination of cons/list/append I use scheme/emacs is combining my let expressions into one (even though each recursive call i use an individual (quote let)). is there a way to circumvent this?

this is my helper function code:

#### the instructors' answer.

where instructors collectively construct a single answer

I believe I see what's going wrong for you here. The initial issue is your function add-lists is using a lambda with what's termed "var-args". If you don't enclose list-of-lets in parentheses, then the function will 1) accept any number of arguments, and 2) collect all its arguments, themselves as a list, in list-of-lets.

To take a small example:

```
> (simple-example 'foo 'bar 'baz)
'((foo bar baz) (foo bar baz))
>
```

These so-called 'var-args' functions can be useful. For instance

```
> (+ 1 2)
3
> (+ 1 2 3)
6
```

We just saw another use above, with list. As a fun exercise, think about how to implement list as a var-args function.

The upshot, I think, is that if you add a set of parens around the function argument, Scheme should instead treat this as an ordinary, one-argument function and I believe will give the behavior you're intending.

## BST problem:

So does scheme have INT\_MAX and INT\_MIN?

#### the instructors' answer,

where instructors collectively construct a single answer

As far as I know, the size of integers is limited only by the amount of memory that is available to Scheme. For example, use the definition of factorial from the day05 folder in Live-in-class. I tried computing the factorial of 20,00, and it worked just fine. The answer ts more than 77,000 digits long. None of the assigned problems require knowing these max and min numbers. Just use the numbers in the tree (perhaps with a post-order traversal).

## Questions for assignment 6

I am having trouble finding the questions for assignment 6a. The homework sheet for assignment 6 says the question is exercise 1.3.4 on page 28 in EoPL, but all I can find on page 28 are exercises 1.xx and above, but none in the format of 1.x.x. Furthermore, none of the questions mention currying. Is there something I'm not understanding?

## the students' answer,

where students collectively construct a single answer

You'll find the problems in EoPL-1, which is the handout/packet we received in class that had four book pages copied onto each page.

## bt-inorder-list Can I use filter-in?

Can I just make a list of in-order traversal of the whole tree then filter in all the symbols? Or I need to check if the value is a symbol when I make the list?

In this problem, does it mean all interior nodes symbols and all leaves integers? If so it would be easier not to use filter-in.

## the instructors' answer,

where instructors collectively construct a single answer

Best: When you traverse the tree to make the list, don't include the leaf nodes. Then you don't have to remove

# Quicksort Algorithm: Does the runtime have to be strictly N(logN) best case?

The implementation I have in mind behaves like Java Quicksort, but has a longer run time than best-case N(log N).

hw6
edit good question0
Updated 2 months ago by
Fric Tu
LIIC TU
the instructors' answer, where instructors collectively construct a single answer
I can't think of any reasonable approach that would not be N log N in the best case. If it is not N log N, is it really

quicksort?