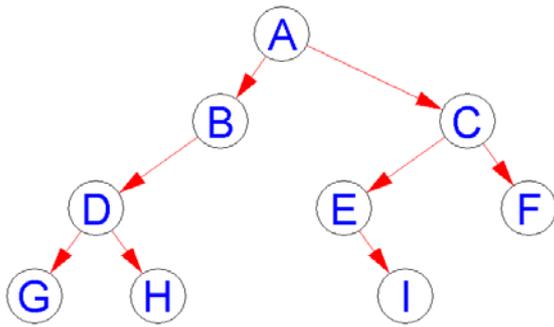


CSSE 230 Day 10

Binary Tree Iterators and Properties Displayable Binary Trees



Displayable Binary Trees
(next assignment)

Picnic time tomorrow afternoon!

DEPARTMENT OF COMPUTER SCIENCE &
SOFTWARE ENGINEERING

PICNIC

SATURDAY SEPTEMBER 28, 2013

HAWTHORN PARK
BURKEYBYLE SHELTER

3:00 – 6:00 P.M.

Outdoor games at 3:00 p.m.

Food served 4:30 p.m.

Food catered by Qdoba Mexican Grill

FOOD SPONSORED BY X-by-2



Agenda

- ▶ Displayable Binary Trees
- ▶ Binary Tree iterators
- ▶ Another induction example
- ▶ WA4 hints, questions

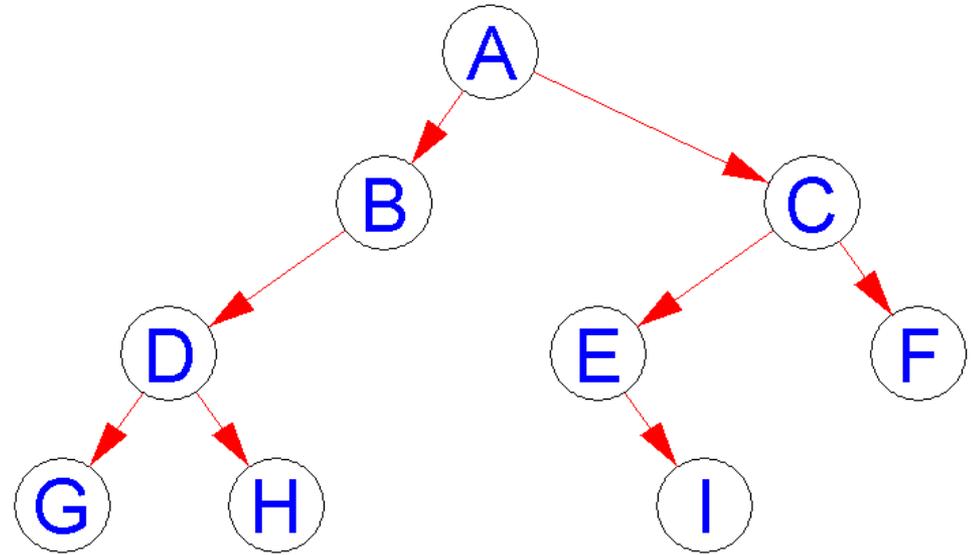
- ▶ # of nodes in Binary tree with height h

Displayable Binary Trees

Individual assignment

Gist of the assignment

- ▶ Levels all spaced evenly.
- ▶ Level-to-right spaced evenly, ordered by in-order traversal.
- ▶ Node and font sizes depend on window size



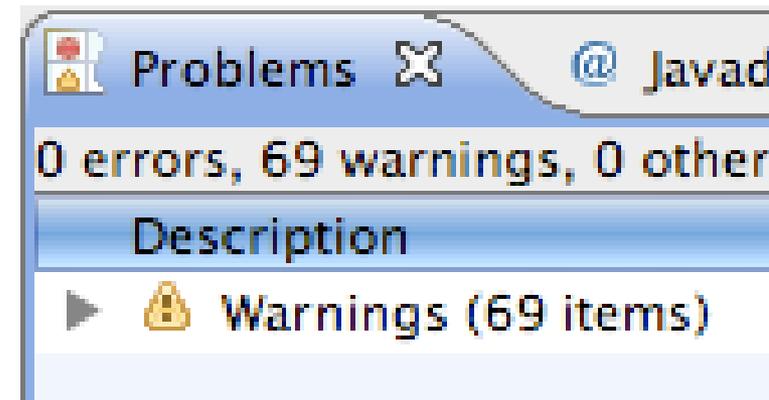
See demo of solution,
trees #9 and #7

Graphics Hints

- ▶ Suggested order for your graphics work:
 - Figure out how to calculate node locations
 - Get code to display correctly sized windows
 - Add code to draw nodes
 - Add code to draw lines
 - Only work on arrow heads if all the rest works!

Hints and getting started

- ▶ Check out **Displayable** from your individual repo.
- ▶ If you get errors on the Weiss imports like `import weiss.nonstandard.Stack;` then install the Weiss packages now (see link from Schedule page)
- ▶ Should be no errors.
- ▶ If errors, see next slide.



Troubleshooting the Weiss install

- ▶ Close all Eclipse projects except *Displayable*
- ▶ Did you put jars in the right folder?
- ▶ Are they jars and not zips?
- ▶ Is Eclipse using that JRE?
 - See *Windows* → *Preferences*, then *Java* → *Installed JREs* → *Edit*.
 - They should be in that list.

Get help now if you're stuck.
Help others if you aren't.

Displayable Binary Trees Steps

- ▶ Solve the sub-problems in this order:
 - *BuildTree.preOrderBuild()*
 - *BinaryTree.inOrder()*
 - Graphics
- ▶ Run *CheckDisplayableBinaryTree* to test
 - Doesn't use JUnit
 - Tests *preOrderBuild* and *inOrder* first
 - Prompts for test case for which to display graphics
 - Each tree should be displayed in a separate window.

Better Exception Reporting in CheckDisplayableBinaryTrees

- ▶ Add a stack trace in *main()*

```
70         tp.dbTree.display();
71     } catch (InternalError e) {
72         System.out
73             .println("You must i
74     }
75
76     }
77 } catch (Exception e) {
78     System.out.println(e.toString());
79     e.printStackTrace();
80 }
81
82 }
83
84 private static String inOrder(int index) {
85     switch (index) {
```

preOrderBuild Hints

- ▶ Like WA4, problem 3
- ▶ Consider:
 - *chars* = 'ROSEHULMAN'
 - *children* = '22002RORLO'

inOrder Hints

- ▶ The iterators in `TestTreeIterators.java` are there for a reason!
- ▶ Recall how we can use Weiss iterators in a for loop:
 - ```
for(iter.first();iter.isValid();iter.advance()) {
 Object elem = iter.retrieve();
 // ... do something with elem ...
}
```

# Merge Method (from Weiss chapter 18)

- ▶ */\*\* Replaces the root element of this tree with the given item and the subtrees with the given ones. ... \*/*

```
public void merge(T rootItem,
 BinaryTreeNode<T> left,
 BinaryTreeNode<T> right)
```

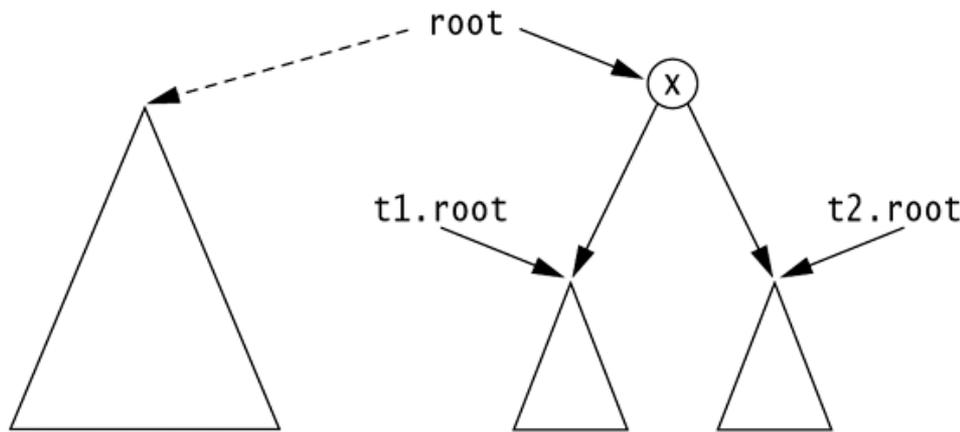
- ▶ Simple approach:

- *this.root = new BinaryTreeNode<T>(rootItem,  
 left.root,  
 right.root);*

What could go wrong?

# Problems With Naïve Merge

- ▶ A node should be part of one and only one tree.

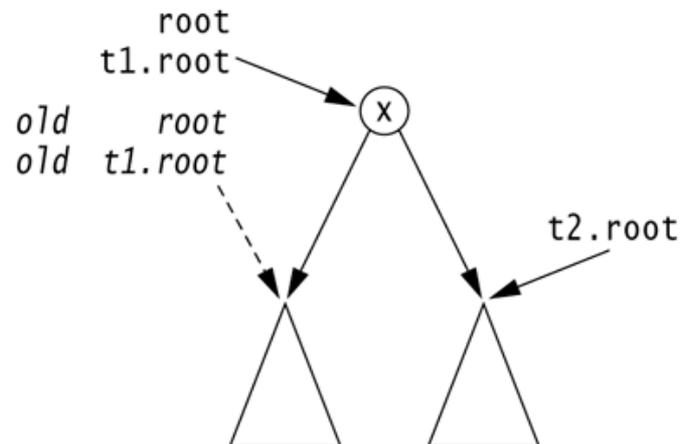


**figure 18.14**

Result of a naive merge operation:  
Subtrees are shared.

**figure 18.15**

Aliasing problems in the merge operation; t1 is also the current object.



# Correct Merge Method

```
1 /**
2 * Merge routine for BinaryTree class.
3 * Forms a new tree from rootItem, t1 and t2.
4 * Does not allow t1 and t2 to be the same.
5 * Correctly handles other aliasing conditions.
6 */
7 public void merge(AnyType rootItem,
8 BinaryTree<AnyType> t1, BinaryTree<AnyType> t2)
9 {
10 if(t1.root == t2.root && t1.root != null)
11 throw new IllegalArgumentException();
12
13 // Allocate new node
14 root = new BinaryNode<AnyType>(rootItem, t1.root, t2.root);
15
16 // Ensure that every node is in one tree
17 if(this != t1)
18 t1.root = null;
19 if(this != t2)
20 t2.root = null;
21 }
```

Weiss, figure 18.16

# Binary Tree Iterators

What if we want to iterate over the elements in the nodes of the tree one-at-a-time instead of just printing all of them?

# Big questions

- ▶ How do you “slow down” recursion to be one step at a time?
  - Hint: what data structure is used to hold recursive calls at runtime?
  
- ▶ How many times is each node visited in a traversal/iterator?
  - See the visualization linked to from day 9 schedule.

# Implementing Binary Tree Iterators

- ▶ What methods does an iterator typically provide?
- ▶ How do we get to the first item in:
  - a pre-order traversal?
  - an in-order traversal?
  - a post-order traversal?
- ▶ In what order should we advance?
- ▶ What instance variables do we need?
  
- ▶ The *Displayable* project has Weiss's *TestTreeIterators* implementation
  - Most of the code is on the next slides.

# Treeliterator abstract class

```
// TreeIterator class; maintains "current position"
//
// CONSTRUCTION: with tree to which iterator is bound
//
// *****PUBLIC OPERATIONS*****
// first and advance are abstract; others are final
// boolean isValid() --> True if at valid position in tree
// Object retrieve() --> Return item in current position
// void first() --> Set current position to first
// void advance() --> Advance (prefix)
// *****ERRORS*****
// Exceptions thrown for illegal access or advance
```

# Treeliterator fields and methods

```
protected BinaryTree t; // Tree
protected BinaryNode current; // Current position

public TreeIterator(BinaryTree theTree) {
 t = theTree;
 current = null;
}

abstract public void first();

final public boolean isValid() {
 return current != null;
}

final public Object retrieve() {
 if(current == null)
 throw new NoSuchElementException();
 return current.getElement();
}

abstract public void advance();
```

# Preorder: constructor and *first*

```
private Stack s; // Stack of TreeNode objects

public PreOrder(BinaryTree theTree) {
 super(theTree);
 s = new ArrayStack();
 s.push(theTree.getRoot());
}

public void first() {
 s.makeEmpty();
 if(t.getRoot() != null)
 s.push(t.getRoot());
 try
 { advance(); }
 catch(NoSuchElementException e) { } // Empty tree
}
```

# PreOrder: *advance*

```
public void advance() {
 if(s.isEmpty()) {
 if(current == null)
 throw new NoSuchElementException();
 current = null;
 return;
 }

 current = (BinaryNode) s.topAndPop();

 if(current.getRight() != null)
 s.push(current.getRight());
 if(current.getLeft() != null)
 s.push(current.getLeft());
}
```

# LevelOrder: constructor and *first*

```
private Queue q; // Queue of TreeNode objects

public LevelOrder(BinaryTree theTree) {
 super(theTree);
 q = new ArrayQueue();
 q.enqueue(t.getRoot());
}

public void first() {
 q.makeEmpty();
 if(t.getRoot() != null)
 q.enqueue(t.getRoot());
 try
 { advance(); }
 catch(NoSuchElementException e) { } // Empty tree
}
```

# Preorder: constructor and *first*

```
private Stack s; // Stack of TreeNode objects
```

```
public PreOrder(BinaryTree theTree) {
 super(theTree);
 s = new ArrayStack();
 s.push(theTree.getRoot());
}
```

```
public void first() {
 s.makeEmpty();
 if(t.getRoot() != null)
 s.push(t.getRoot());
 try
 { advance(); }
 catch(NoSuchElementException e) { } // Empty tree
}
```

# LevelOrder: *advance*

```
public void advance() {
 if(q.isEmpty()) {
 if(current == null)
 throw new NoSuchElementException();
 current = null;
 return;
 }

 current = (BinaryNode) q.dequeue();

 if(current.getLeft() != null)
 q.enqueue(current.getLeft());
 if(current.getRight() != null)
 q.enqueue(current.getRight());
}
```

# PreOrder: *advance*

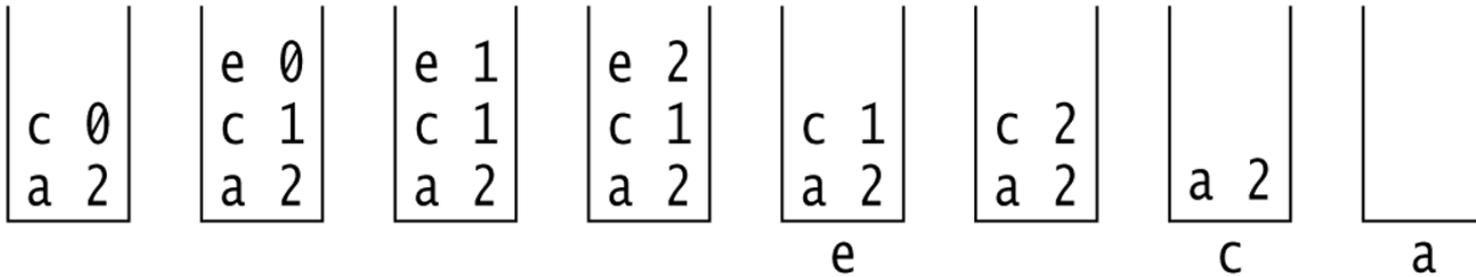
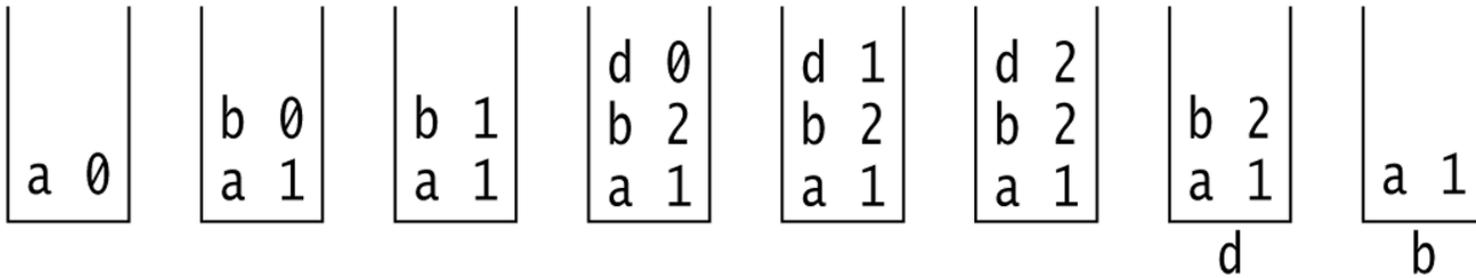
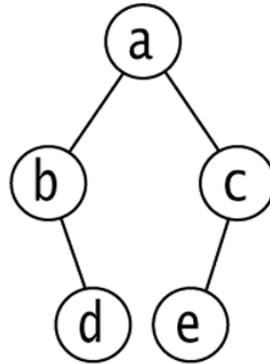
```
public void advance() {
 if(s.isEmpty()) {
 if(current == null)
 throw new NoSuchElementException();
 current = null;
 return;
 }
}
```

```
current = (BinaryNode) s.topAndPop();
```

```
if(current.getRight() != null)
 s.push(current.getRight());
if(current.getLeft() != null)
 s.push(current.getLeft());
}
```

# The Stack in a PostOrder iterator

Q1-4



# Other Approaches to Tree Iterators

Weiss's way isn't the only one

# Alternative:

- ▶ Each node can store pointer to the next node in a traversal
- ▶ Must update extra info in constant time as tree changes

An upcoming written assignment will include these “threaded binary trees”

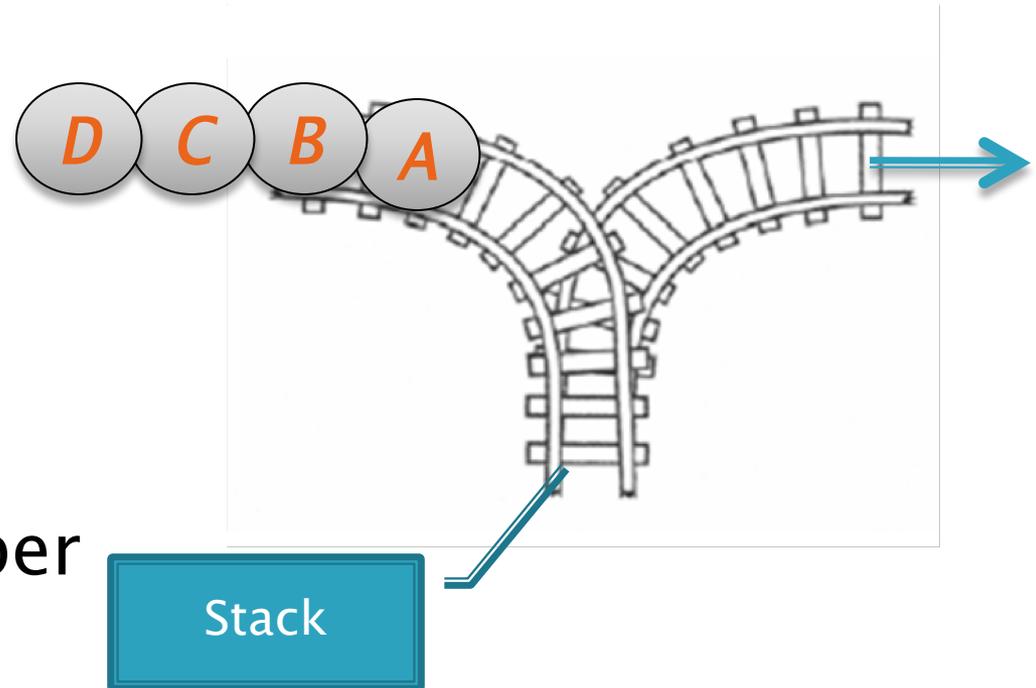
# Wouldn't it be nice?

- ▶ If we did not have to maintain the stack for these iterators?
- ▶ If we could somehow “tap into” the stack used in the recursive traversal?
  - I.e. Take a “snapshot of that call stack, and restore it later when we need it.
  - This is called a **continuation**.
    - A big subject in the PLC course, CSSE 304

# Tips on WA4

# WA4, Problem 2 Application

- ▶ Railroad switching
- ▶ Problem is equivalent to counting the number of possible orders the cars can leave the station

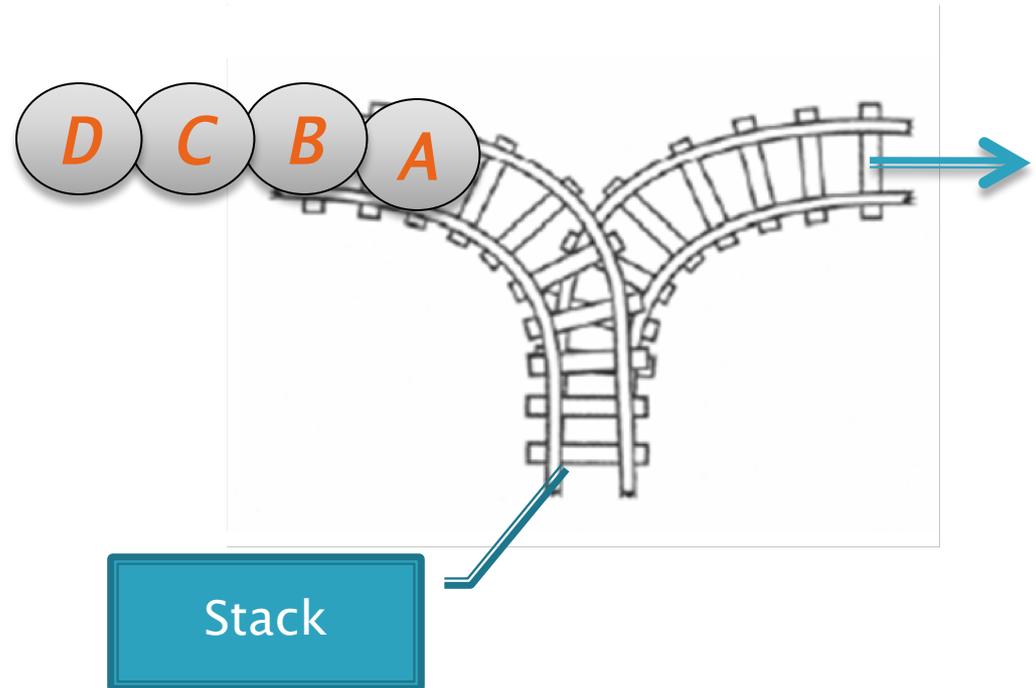


# General Approach to Puzzle Problems

- ▶ Make up tiny examples like the given problem
  - No really tiny, I'm serious
- ▶ Solve the tiny problem
- ▶ Solve a slightly larger problem
- ▶ Solve a slightly larger problem than that
- ▶ Once you see the pattern, then try to solve the given problem

# What's the smallest problem like this?

- ▶ In how many possible orders can the cars leave the station?



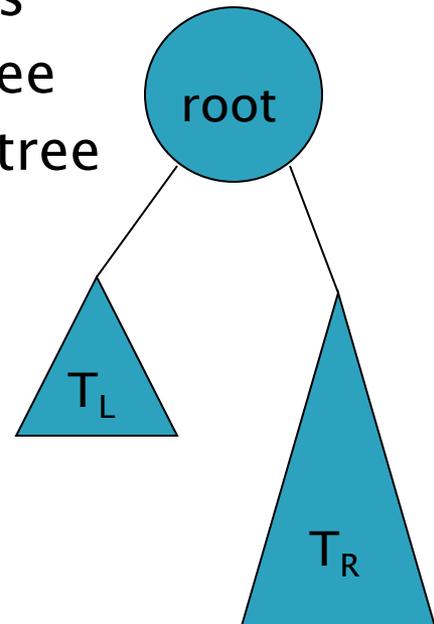
# More Binary Trees

If a tree falls in the forest and there are two people around to hear it...

# Size vs. Height in Binary Trees

# Binary Tree: Recursive definition

- ▶ A Binary Tree is either
  - **empty**, or
  - **consists of:**
    - a distinguished node called the root, which contains an element, and two disjoint subtrees
    - A left subtree  $T_L$ , which is a binary tree
    - A right subtree  $T_R$ , which is a binary tree



# Size and Height of Binary Trees

- ▶ Notation:
  - Let  $T$  be a tree
  - Write  $h(T)$  for the height of the tree, and
  - $N(T)$  for the size (i.e., number of nodes) of the tree
- ▶ Given  $h(T)$ , what are the **bounds** on  $N(T)$ ?
- ▶ Given  $N(T)$ , what are the bounds on  $h(T)$ ?

# Extreme Trees

- ▶ A tree with the maximum number of nodes for its height is a **full tree**.
  - Its height is  **$O(\log N)$**
- ▶ A tree with the minimum number of nodes for its height is essentially a \_\_\_\_\_
  - Its height is  **$O(N)$**
- ▶ Height matters!
  - We will see that the algorithms for search, insertion, and deletion in a Binary search tree are  **$O(h(T))$**

# Time out for math!

- ▶ Want to prove some properties about trees
- ▶ Weak induction isn't enough
- ▶ Need strong induction instead:



The former  
governor of  
California

# Strong Induction

- ▶ To prove that  $p(n)$  is true for all  $n \geq n_0$ :
  - Prove that  $p(n_0)$  is true, and
  - For all  $k > n_0$ , prove that if we assume  $p(j)$  is true for  $n_0 \leq j < k$ , then  $p(k)$  is also true
- ▶ Weak induction uses the previous domino to knock down the next
- ▶ Strong induction uses a whole box of dominoes to knock down the rest!