

<b>Exam 1</b>
---------------

**Name:** \_\_\_\_\_

**Instructions:**

- Write all answers on these pages. Use the back as necessary.
- Clearly indicate your final answer.
- For full credit, show your work, and document your code.
- Read the entire examination before starting, and then budget your time.

**Authorized resources:**

- Green reference card from the text.

**Unauthorized resources:**

You are NOT permitted to use any resources other than those identified above. In particular, you may NOT use books, notes, electronic files, calculators, PDAs, or computers.

Good luck!

<b>Problem Number</b>	<b>Maximum Points</b>	<b>Points Earned</b>
<b>1a</b>	<b>12</b>	
<b>1b</b>	<b>12</b>	
<b>1c</b>	<b>8</b>	
<b>1d</b>	<b>8</b>	
<b>2a</b>	<b>12</b>	
<b>2b</b>	<b>17</b>	
<b>3</b>	<b>16</b>	
<b>4</b>	<b>14</b>	
<b>Total</b>	<b>99 + 1 bonus</b>	

### Problems

**Problem 1a.** [12 points] List the machine language fields and obtain the hexadecimal representation for the following MIPS instructions:

```
lw  $t0, 4($t1)
```

```
addi $s2, $s1, -5
```

```
sub  $t3, $t5, $a0
```

**Problem 1b.** [12 points] Give the MIPS assembly language statements represented by each of the following:

```
0x27a50004
```

```
0x001a2082
```

```
0x116a0004
```

**Problem 1c.** [8 points] Assume that the MIPS instruction

```
beq $t0, $s0, Label
```

is located at address `0x0100 0040`, and that the 16-bit immediate field has the value

```
0000 0000 1000 0100
```

What is the 32-bit effective address value of `Label`? Express your answer in HEXADECIMAL. *Hint:* Remember that the offset is the signed number of instructions relative to the instruction FOLLOWING the branch.

**Problem 1d** - [8 points] Assume that the MIPS instruction `j Label` is located at address `0x8000 0000`, and `Label` is located at `0x8A00 0540`. What is the value of the 26-bit address field? Express your answer in BINARY.

**Problem 2.** You are designing an architecture with variable length instructions. Some instructions are 12 bits long, while others are 24 bits long. There are 8 general purpose registers and each of these registers is 12 bits wide. Addresses and immediates are also 12-bits wide. The instructions and their types are listed in the table below.

	Instruction	Type	Action	Length
1	Add rd, rs	A	$\text{Reg}[\text{rd}] = \text{Reg}[\text{rs}] + \text{Reg}[\text{rd}]$	12 bits
2	Sub rd, rs	A	$\text{Reg}[\text{rd}] = \text{Reg}[\text{rs}] - \text{Reg}[\text{rd}]$	12 bits
3	Jr rs	A	$\text{PC} = \text{Reg}[\text{rs}]$	12 bits
4	And rd, rs	A	$\text{Reg}[\text{rd}] = \text{Reg}[\text{rd}] \text{ and } \text{Reg}[\text{rs}]$	12 bits
5	Or rd, rs	A	$\text{Reg}[\text{rd}] = \text{Reg}[\text{rs}] \text{ or } \text{Reg}[\text{rd}]$	12 bits
6	Getinput rd, rs	A	$\text{Reg}[\text{rd}] = \text{SpecialRegister}[\text{rs}]$	12 bits
7.	Putoutput rd, rs	A	$\text{SpecialRegister}[\text{rd}] = \text{Reg}[\text{rs}]$	12 bits
8	Addi, rd, rs, 12-bit immediate	B	$\text{Reg}[\text{rd}] = \text{Reg}[\text{rs}] + 12\text{-bit immediate}$	24 bits
9	Beq rd, rs, 12-bit address	B	If ( $\text{Reg}[\text{rs}] = \text{reg}[\text{rd}]$ ) $\text{PC} = 12\text{-bit address}$	24 bits
10	Bne rd, rs, 12-bit address	B	If ( $\text{Reg}[\text{rs}] \neq \text{Reg}[\text{rd}]$ ) $\text{PC} = 12\text{-bit address}$	24 bits
11	Lw rd, 12-bit address (rs)	B	$\text{Reg}[\text{rd}] = \text{Mem}[\text{Reg}[\text{rs}] + 12\text{-bit address}]$	24 bits
12	Sw rd, 12-bit address (rs)	B	$\text{Mem}[\text{Reg}[\text{rs}] + 12\text{-bit address}] = \text{Reg}[\text{rd}]$	24 bits
13	Sl rd, rs, 5-bit immediate	B	If (5-bit immediate > 0) $\text{Reg}[\text{rd}] = \text{Reg}[\text{rs}] \ll 5\text{-bit immediate}$ else $\text{Reg}[\text{rd}] = \text{Reg}[\text{rs}] \gg 5\text{-bit immediate}$	24 bits
14	J 9-bit address	C	$\text{PC} = \text{PC}[11:9] \parallel 9\text{-bit address}$	12 bits
15	Jal 9-bit address	C	$\text{PC} = \text{PC}[11:9] \parallel 9\text{-bit address}$ $\text{Reg}[7] = \text{PC}$	12 bits
16	Rfe	C	$\text{PC} = \text{EPC}$	12 bits
17	Slt rd, rs, rt	D	If ( $\text{Reg}[\text{rs}] < \text{Reg}[\text{rt}]$ ) $\text{Reg}[\text{rd}] = 1$ ; Else $\text{Reg}[\text{rd}] = 0$	12-bits

- a. [12 points] Show the instruction format for each type i.e. A, B, C, and D. Indicate clearly how many bits are allocated for each field. *Hint:* Start with the C and D types.

- b. [17 points] For each of the instructions below, assign a value for the opcode field, as well as values for any fields that augment the opcode (e.g. the funct field for MIPS R-type instructions)

Instruction	Opcode	Values of fields that augment opcode
Add		
Sub		
Jr		
And		
Or		
Getinput		
Putoutput		
Addi		
Beq		
Bne		
Lw		
Sw		
Sl		
J		
Jal		
Rfe		
Slt		

**Problem 3**[16 pts] For each pseudoinstruction listed below, give a minimal sequence of actual MIPS instructions to accomplish the same thing. You may need to use `$at` for some of the sequences. “big” refers to a 32-bit immediate value and “small” refers to a 16-bit immediate value.

a. `move $sp, $s0`

b. `li $a0, big`

c. `lw $v0, big($k0)`

d. `ble $t8, $t9, label`

**Problem 4** [14 pts] Complete the MIPS program on the following pages by filling in the provided spaces with MIPS assembly language statements such that:

- The procedure `dotProduct` accepts three input parameters: the addresses of two arrays and the size of the arrays.
- The procedure `dotProduct` returns the dot product of the vectors represented by the two arrays.
- The procedure `dotProduct` follows MIPS register usage conventions.
- The main program calls the procedure `dotProduct` in such a way that the dot product  $(0, 1, 2, 3, 4, 5, 6, 7, 8, 9) \cdot (1, 2, 3, 4, 5, 6, 7, 8, 9, 10) = 0 + 2 + 6 + 12 + \dots + 90 = 330$  is displayed.

Assume the existence of another MIPS procedure `product` that returns in `$v0` the product of its two integer parameters, which are passed in `$a0` and `$a1`. In other words, even though `product` is not shown, you may call it, and you do not need to write it.

Read all the provided parts of the program, before you begin.

```

# Procedure dotProduct calculates and returns x (dot) y,
# where the address of x, the address of y, and their
# (common) size are the three parameters of the procedure.
#
# A high-level language description of the procedure follows:
#
#
#     /**
#     * @param x First array
#     * @param y Second array
#     * @param size Number of elements in each array
#     * @return sum The dot product
#     */
#     private static int dotProduct(
#         int[] x, int[] y, int size ) {
#         int sum = 0;
#         for( int count = 0; count < size; count++) {
#             sum = sum + product( x[ count ] , y[ count ] );
#         }
#         return sum;
#     }
#
# "Public" register usage:
#
# $a0 - address of x
# $a1 - address of y
# $a2 - number of elements in each array
# $v0 - x (dot) y
#
# "Private" register usage:
#
# $s0 - address of x[ count ]
# $s1 - address of y[ count ]
# $s2 - size
# $s3 - count
# $s4 - sum
#
# $t0 - flag
#
# Procedure entrance and initialization
#
dotProduct:
    addi    $sp, $sp, -16    # Create space on the stack

    sw     $ra, 0($sp)      # Place values to be preserved
    sw     $s0, 4($sp)      # on the stack.
    sw     $s1, 8($sp)
    sw     $s2, 12($sp)

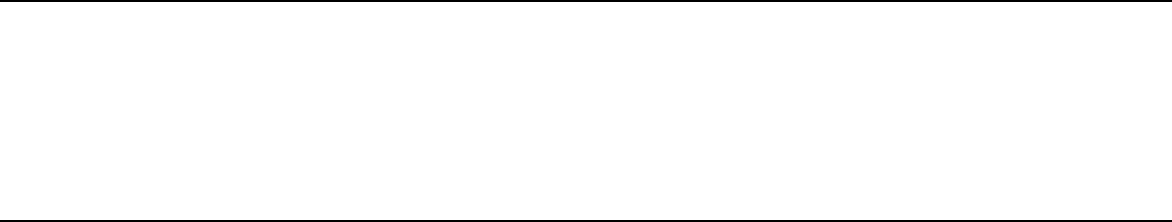
                                # Read input arguments from the
                                # argument registers
    move   $s0, $a0        # $s0 = address of x[ count ]
    move   $s1, $a1        # $s1 = address of y[ count ]
    move   $s2, $a2        # $s2 = size

    move   $s3, $0         # count = 0
    move   $s4, $0         # sum = 0

loop:     slt     $t0, $s3, $s2    # if( count < size)
    beq   $t0, $0, exit1    # continue

    lw    $t1, 0($s0)        # Read x[ count ] from memory
    lw    $t2, 0($s1)        # Read y[ count ] from memory
                                # Call "product" and pass
                                # parameters

```



```
        add      $s4, $s4, $t3      # sum = sum +
                                       #      product (x[count],y[count])
        addi     $s0, $s0, 4         # $s0 = $s0 + 4
        addi     $s1, $s1, 4         # $s1 = $s1 + 4
        addi     $s3, $s3, 1         # count++

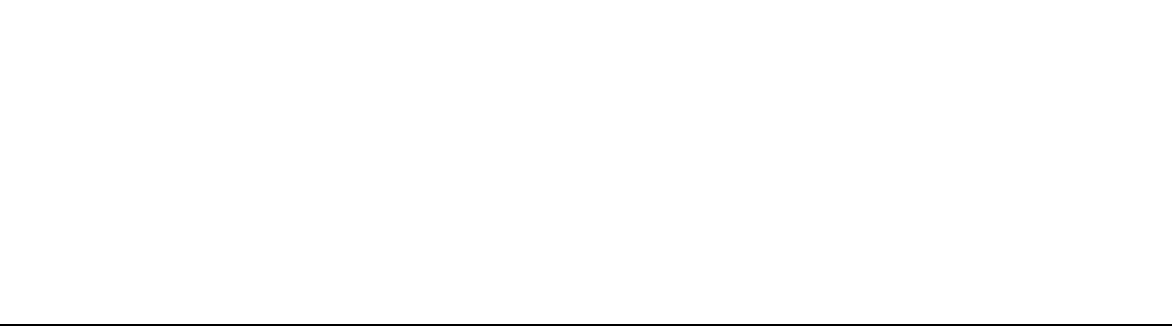
        j       loop

#
# Exit from dotProduct procedure
#
exit1:   move    $v0, $s4           # Move sum to the return value
                                       # register
        lw      $ra, 0($sp)        # Restore the values of the
        lw      $s0, 4($sp)        # preserved registers from the
        lw      $s1, 8($sp)        # stack.
        lw      $s2, 12($sp)
        addi    $sp, $sp, 16       # Restore the stack pointer

        jr     $ra                # Return to calling procedure

#
# Main program starts here
#
main:   la $s0, x                  # x
        la $s1, y                  # y
        la $s2, N                  # size
        lw $s2, 0($s2)

                                       # Call "dotproduct" and pass
                                       # parameters (The registers to use
                                       # are listed at the beginning of the
                                       # program on page 8.)
```



```
        la $a0, Msg                # Print message
        li $v0, 4
        syscall

        move $a0, $s0              # Print the value of the value
                                       # returned by the procedure "dotproduct"
        li $v0, 1
        syscall

        li $v0, 10                 # Quit program
        syscall
```

**# The program continues on the next page**

```
#  
# Sample data starts here  
#  
      .data  
x:    .word 0 1 2 3 4 5 6 7 8 9  
y:    .word 1 2 3 4 5 6 7 8 9 10  
N:    .word 10  
Msg:  .asciiiz "The dot product is"
```