

CSSE 220 Day 14

Sorting Algorithms
Algorithm Analysis and Big-O
Searching

Checkout *SortingAndSearching* project from SVN

Exam recap

- ▶ Look at a few problems together
 - Computer part, and problems 3, 5, 6.
- ▶ Student questions?
- ▶ If you think one of your problems was incorrectly graded, talk with me during the work time at the end of class.

- ▶ Questions about anything in the course?

What is sorting?

»» Let's see...

Why study sorting?

»» Shlemiel the Painter

Course Goals for Sorting: You should...

- ▶ Be able to **describe** basic sorting algorithms:
 - Selection sort
 - Insertion sort
 - Merge sort
 - Quicksort
- ▶ Know the **run-time efficiency** of each
- ▶ Know the **best and worst case** inputs for each

Selection Sort

- ▶ Basic idea:
 - Think of the list as having a **sorted part** (at the beginning) and an **unsorted part** (the rest)
 - Find the **smallest** value in the unsorted part
 - Move it to the **end** of the sorted part (making the sorted part bigger and the unsorted part smaller)

Repeat until
unsorted part is
empty

Profiling Selection Sort

- ▶ **Profiling**: collecting data on the run-time behavior of an algorithm
- ▶ How long does selection sort take on:
 - 10,000 elements?
 - 20,000 elements?
 - ...
 - 80,000 elements?

Q1

Analyzing Selection Sort

- ▶ **Analyzing**: calculating the performance of an algorithm by studying how it works, typically mathematically
- ▶ Typically we want the **relative** performance as a function of input size
- ▶ Example: For an array of length n , how many times does `selecti onSort ()` call `compareTo ()`?

Handy Fact

$$1 + 2 + \dots + (n - 1) + n = \frac{n(n + 1)}{2}$$

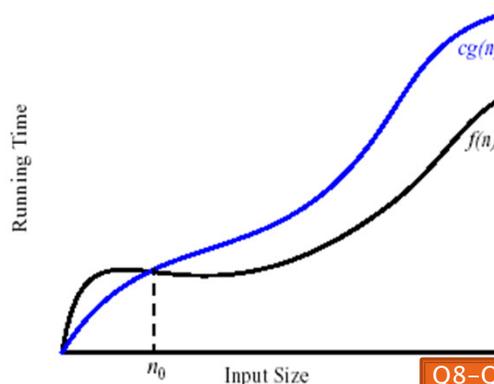
Q2-Q7

Big-Oh Notation

- ▶ In analysis of algorithms we care about differences between algorithms on very large inputs
- ▶ We say, “selection sort takes on the order of n^2 steps”
- ▶ Big-Oh gives a formal definition for “on the order of”

Formally

- ▶ We write $f(n) = O(g(n))$, and say “ f is big-Oh of g ”
- ▶ if there exists positive constants c and n_0 such that
- ▶ $0 \leq f(n) \leq c g(n)$ for all $n > n_0$
- ▶ g is a **ceiling** on f



Q8-Q9



Another Interesting Comic on Sorting ... follow link
<http://www.smbc-comics.com/?db=comics&id=1989>

Perhaps it's time for a break.

Insertion Sort

- ▶ Basic idea:
 - Think of the list as having a **sorted part** (at the beginning) and an **unsorted part** (the rest)
 - Get the **first** value in the unsorted part
 - Insert it into the **correct** location in the sorted part, moving larger values up to make room

Repeat until
unsorted
part is
empty

Insertion Sort Exercise, Q10–19

- ▶ **Profile** insertion sort
- ▶ **Analyze** insertion sort assuming the inner while loop runs the maximum number of times
- ▶ What input causes the worst case behavior? The best case?
- ▶ Does the input affect selection sort?

Ask for help if you're stuck!

Q10-Q19

Searching

- ▶ Consider:
 - Find Cary Laxer's number in the phone book
 - Find who has the number 232-2527
- ▶ Is one task harder than the other? Why?
- ▶ For searching unsorted data, what's the worst case number of comparisons we would have to make?

Binary Search of Sorted Data

- ▶ A **divide and conquer** strategy
- ▶ Basic idea:
 - Divide the list in half
 - Decide whether result should be in upper or lower half
 - Recursively search that half

Analyzing Binary Search

- ▶ What's the best case?
- ▶ What's the worst case?
- ▶ We use **recurrence relations** to analyze recursive algorithms:
 - Let $T(n)$ count the number of comparisons to search an array of size n
 - Examine code to find recursive formula of $T(n)$
 - Solve for n

Q20-Q21

Work Time

»» Review Homework.

Q22-Q23