GC Design Choices

What about other storage reclamation schemes?

Memory management options

- Manual /explicit memory management
 - Strengths?
 - Challenges?
- Automated memory management (garbage collection)
 - Strengths?
 - Challenges?
- Any others?

Real-time garbage collection (RTGC)

- Real-time system
 - A system that meets real-time requirements.
- Real-time requirements
 - As expected, operations must be logically correct
 - Additionally, operations must be completed within deadline
- RTGC
 - Bounded-time allocation
 - Predictable deallocation
 - Must be incremental

Real-time garbage collection (RTGC)

```
public void f() {
    startLaser();
    Obj o = new Obj();
    stopLaser();
}
public static void main(...) {
    f();
}
```



RTGC strengths and challenges

- Need extra storage
 - Store state of application when collector runs
- Application can allocate memory during garbage collection
- Space-time trade-off

RTSJ scoped-memory

- RTSJ Real-time specification for Java proposed by the Real-time for Java expert group (RTJEG).
- Semi-manual with scopes
 - Scopes: regions of memory
 - Scopes: limited life times
 - Threads allocate from current scope
 - Predictable allocation
 - Predictable deallocation
 - No dangling pointers

RTSJ scoped-memory

```
ScopedMemory
scope = new ScopedMemory(1024);
scope.enter(new Runnable() {
  public void run(){
    // do some stuff
    someObj o = new someObj();
    // do some more stuff
    someObj s = new someObj();
  }
});
// scope is collected (no threads)
```

RTSJ scoped-memory challenges

- Restrictive memory model
- Difficult to use
- Can leak memory

Memory management options

- Manual/explicit memory management
- Automated memory management (GC)
- Real-time garbage collection
- RTSJ scoped-memory

Garbage collection design choices

- Stop-the-world
- Incrementality
- Hybrid
- Concurrency
- Parallelism

Stop-the-world collectors

- Typically used on uniprocessor systems
- Suspend application
- Run collector from start to finish
- Resume application

Stop-the-world collectors

• Execution costs?

- Pause time
- Discovery of live objects (how long does it take?)
- Instruction overhead (per instruction)
- Delay between object death and collection
- Number of collectible objects collected
- Overall execution time
- Worst-case vs average case performance
- frequency

Incremental collection

- Interleave GC with application
- Note: for full heap tracing
 - Pause time increases with heap size
- Incremental tracing
 - Bounded tracing time
 - Conservative assumption
 - All other objects in heap are live
 - Remember pointers from objects in heap
 - Add such pointers to root set for tracing

Hybrid collection

- Generational collectors
 - Collect young objects frequently
 - Young objects die quickly
- Example
 - Copy collection for young objects
 - Non-copy collection for older objects
- Partitioning
 - Copy intra-partition incrementally
 - Reference count inter-partition

Concurrent collection

- Application is called a *mutator*
- GC regards application as such because it is mutating the heap
- Mutator and GC function at the same time except when GC needs info from mutator
 - Synchronization

Parallel collection

- Concurrency among multiple GC threads
 - Load balancing
 - Synchronization
 - Race condition when tracing