## DTTF/NB479: Dszquphsbqiz

- Announcements:
. Homework 2 due now
- Computer quiz Friday on chapter 2

Questions?

- Today:
- Wrap up congruences
- Fermat's little theorem
- Euler's theorem
- Both really important for RSA - pay careful attention!
- A single congruence mod a composite number is equivalent to a system of congruences mod its factors
- Two-factor form
- Given $\operatorname{gcd}(m, n)=1$. For integers a and $b$, there exists exactly 1 solution (mod mn ) to the system:


## $x \equiv a(\bmod m)$ $x \equiv b(\bmod n)$

# CRT Equivalences let us use systems of congruences to solve problems 

- Solve the system:

$$
\begin{aligned}
& x \equiv 3(\bmod 7) \\
& x \equiv 5(\bmod 15)
\end{aligned}
$$

- How many solutions?
- Find them.

Chinese Remainder Theorem

- n-factor form
- Let $m_{i}, m_{2}, \ldots m_{k}$ be integers such that $\operatorname{gcd}\left(m_{i}, m_{j}\right)=1$ when $\mathrm{i} \neq \mathrm{j}$. For integers $\mathrm{a}_{1}, \ldots \mathrm{a}_{\mathrm{k}}$, there exists exactly 1 solution ( $\bmod m_{1} m_{2} \ldots m_{k}$ ) to the system:

$$
\begin{aligned}
& x \equiv a_{1}\left(\bmod m_{1}\right) \\
& x \equiv a_{2}\left(\bmod m_{2}\right) \\
& \cdots \\
& x \equiv a_{k}\left(\bmod m_{k}\right)
\end{aligned}
$$

Modular Exponentiation is extiremely efficient since the partial resulis are always small

- Compute the last digit of $3^{\wedge} 2000$
- Compute 3^2000 (mod 19) Idea:
- Get the powers of 3 by repeatedly squaring 3, BUT taking mod at each step.


## Modular Exponentiation Technique and Example

(All congruences are mod 19)

- Compute 3^2000 (mod 19)
- Technique:
- Repeatedly square 3, but take mod at each step.
- Then multiply the terms you need to get the desired power.
- Book's powermod()


## Modular Exponentiation Example

- Compute 3^2000 (mod 152)

$$
\begin{aligned}
& 3^{2} \equiv 9 \\
& 3^{4}=9^{2} \equiv 81 \\
& 3^{8}=81^{2} \equiv 6561 \equiv 25 \\
& 3^{16}=25^{2} \equiv 625 \equiv 17 \\
& 3^{32}=17^{2} \equiv 289 \equiv 137 \\
& 3^{64}=137^{2} \equiv 18769 \equiv 73 \\
& 3^{128} \equiv 9 \\
& 3^{256} \equiv 81 \\
& 3^{512} \equiv 25 \\
& 3^{1024} \equiv 17 \\
& 3^{2000} \equiv\left(3^{1024}\right)\left(3^{512}\right)\left(3^{256}\right)\left(3^{128}\right)\left(3^{64}\right)\left(3^{16}\right) \\
& 3^{2000} \equiv(17)(25)(81)(9)(73)(17) \\
& 3^{2000} \equiv(384492875) \\
& 3^{2000} \equiv 9(\bmod 152)
\end{aligned}
$$

Fermat's Little Theorem:
1-2
If $p$ is prime and $\operatorname{gcd}(a, p) \neq 1$, then $a^{(p-1)} \equiv 1(\bmod p)$

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Example: $a=2, p=7$

Examples:

- $2^{2}=1(\bmod 3)$
- $6^{4}=1(\bmod ? ? ?)$
- $\left(3^{2000}\right)(\bmod 19)$


## The converse when $a=2$ usually holds

- Fermat:

If $\rho$ is prime and doesn't divide $\mathrm{a}, a^{p-1} \equiv 1(\bmod p)$

- Converse:
- If $a^{p-1} \equiv 1(\bmod p)$, then p is prime and doesn't divide a.
- This is almost always true when $\mathrm{a}=2$. Rare counterexamples:
- $n=561=3^{*} 11^{*} 17$, but


## $2^{560} \equiv 1(\bmod 561)$

- $n=1729=7^{*} 13^{*} 19$
- Can do first one by hand if use Fermat and combine results with

Chinese Remainder Theorem $\qquad$

## Primality testing schemes typically use Fermat



## Primality testing schemes typically use Fermat

Use Fermat as a filter since it's faster than factoring (if calculated using the powermod method).

Fermat: p prime $\rightarrow 2^{p-1} \equiv 1(\bmod p)$ Contrapositive?

Why can't we just compute $2^{n-1}(\bmod n)$ using Fermat if it's so much faster?


Euler's Theorem is like Fermat's, but for composite modulf

If $p$ is prime and $\operatorname{gcd}(a, p) \neq 1$, then

$$
a^{\phi(n)} \equiv 1(\bmod n)
$$

So what's $\phi(n) ?$

# $\phi(n)$ is the number of integers a, 

 such that $1 \leq a \leq n$ and $\operatorname{gcd}(a, n)=1$.Example: $\phi(10)=4$.

When $p$ is prime, $\phi(p)=$

When $n=p q$ (product of 2 primes), $\phi(n)=$

## The general formula for $\phi(n)$



Example: $\phi(60)=16$

Euler's Theorem can lead to computations that are 7-10 more efficient even than modular exponentiation

## $a^{\phi(n)} \equiv 1(\bmod n)$

as long as $\operatorname{gcd}(a, n)=1$

Principle: when working mod $n$, view the exponents $\bmod \phi(n)$.

## Examples:

1. Find last 3 digits of $7^{803}$
2. Find $3^{2007}(\bmod 12)$
3. Find $2^{6004}(\bmod 99)$
4. Find $2^{6004}(\bmod 101)$
