

### ● Announcements:

- Please pass in Assignment 1 now.
- Assignment 2 posted (when due?)

### ● Questions?

### ● Roll Call

### ● Today: Vigenere ciphers

- Invented in 1553 by Bellaso, not Vigenere

# Vigenere Ciphers

- Idea: the key is a *vector* of shifts
  - The key and its length are unknown to Eve
- Encryption:
  - Repeat the vector as many times as needed to get the same length as the plaintext
  - Add this repeated vector to the plaintext.

## ● Example:

- Key = *hidden* (7 8 3 3 4 13).

Key	● The recent development of various methods of
	7 8 3    3 4 13 7 8 3    3 4 13 7 8 3 3 4 13 7 8    3 3    4 13 7 8 3 3 4    13 7 8 3 3 4 13    7 8 0 15 7    20 8 15 11 2 122    6 8 8 11 19 17 18 16 17 20 1    17 8    25 13 24 16 17 23 22    25 11 11 0 17 7 5    2 11 3
	● aph uiplvw giiltrsqrub ri znyqrxw zlbkrhf vn

- Demo

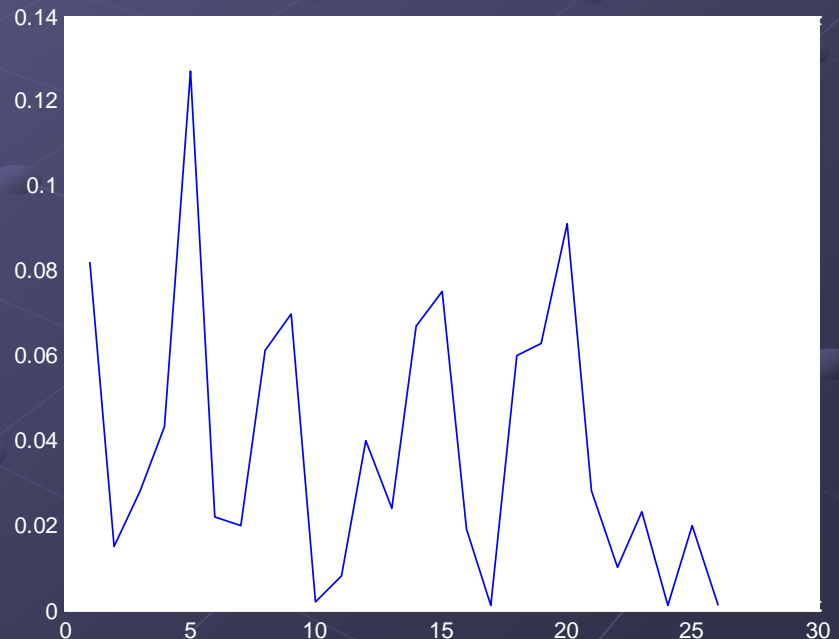
# Security

- The shift vector isn't known (of course)
- 1. It's length isn't even known!
- 2. With shift ciphers, the most frequent cipher letter is probably e.
  - But here, e maps to H, I, L, ... (**spread out!**)
- Consider 4 attacks:
  - Known plaintext?
  - Chosen plaintext?
  - Chosen ciphertext?
  - Ciphertext only?

# English letter frequencies

A 0.082	H 0.061	O 0.075	U 0.028
B 0.015	I 0.070	P 0.019	V 0.010
C 0.028	J 0.002	Q 0.001	W 0.023
D 0.043	K 0.008	R 0.060	X 0.001
E 0.127	L 0.040	S 0.063	Y 0.020
F 0.022	M 0.024	T 0.091	Z 0.001
G 0.020	N 0.067		

Graph:



# Ciphertext-only attack

- Assume you know the key length,  $L$ .
- Make any other assumptions you need.
- Take 3-4 min with a partner and devise a method to break Vigenere.



# Perhaps yours looks something like this?

- Assume we know the key length,  $L$ , ...

- We'll see how to find it shortly

- Method 1:

- Parse out the characters at positions  $p = i \pmod{L}$

- These have all been shifted the same amount

- Do a frequency analysis to find shift

- The most frequent letter should be e, given enough text. Can verify to see how shift affects other letters

- This gives the first letter of the key

- Repeat for positions  $p = 2, p = 3, \dots p = L$

- Problem: involves some trial and error.

- For brute force to work, would need to brute force all letters of key simultaneously: \_\_\_\_\_ possibilities

# Dot products

$$A \cdot B = A.*B = \sum_i A_i B_i$$

Consider  $A =$  (0.082 0.015 0.028 0.043 0.127 0.022 0.020 0.061 0.070 0.002  
0.008 0.040 0.024 0.067 0.075 0.019 0.001 0.060 0.063 0.091 .  
0.028 0.010 0.023 0.001 0.020 0.001);

$A_i = A$  displaced  $i$  positions to the right

$A_0 =$  (0.082 0.015 0.028 ... 0.001 0.020 0.001)

$A_1 =$  (0.001 0.082 0.015 0.028 ... 0.023 0.001 0.020)

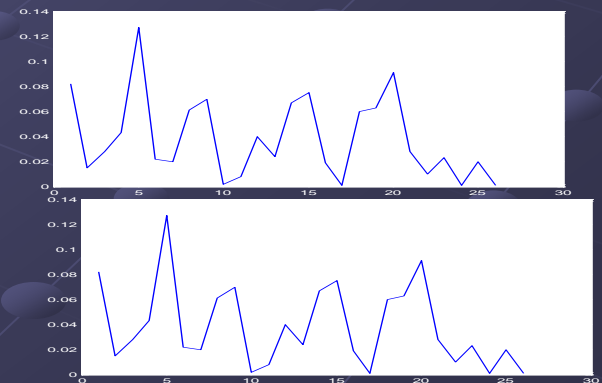
$A_2 =$  (0.020 0.001 0.082 0.015 0.028 ... 0.023 0.001)

$A_0 .* A_1 = 0.039$

$A_0 .* A_0 = 0.066$

$A_i .* A_j$  depends on \_\_\_\_\_ only.

Max occurs when \_\_\_\_\_.  
3 reasons why:



# Towards another method

## ● Method 1

- Parse out the characters at positions  $p = 1 \pmod{L}$ 
  - These have all been shifted the same amount
  - Do a frequency analysis to find shift
    - The most frequent letter should be e, given enough text.  
Can verify to see how shift affects other letters.
- This gives the first letter of the key
- Repeat for positions  $p = 2, p = 3, \dots p = L$



# Another method

## ● Method 2

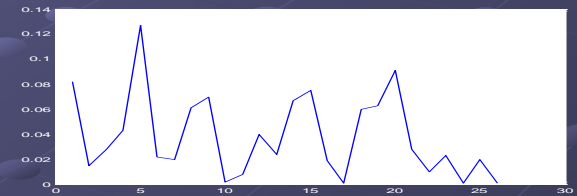
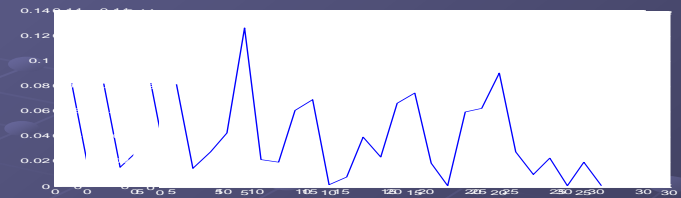
- Parse out the characters at positions  $p = 1 \pmod{L}$ 
  - These have all been shifted the same amount
  - Get the whole freq. distribution  $W = (0.05, 0.002, \dots)$ 
    - $W$  approximates  $A$ . Calculate  $W \cdot A_i$  for  $0 \leq i \leq 25$
    - Max occurs when we got the shift correct.
- This gives the first letter of the key
- Repeat for positions  $p = 2, p = 3, \dots p = L$
- Demo

# Method 2 is more robust since it uses the whole letter distribution

● Find dot product of  $A_i$ :  
and  $W$ :

More robust than just using 1 letter ('e')...

...but harder to compute by hand.



# Finding the key length

- What if the frequency of letters in the plaintext approximates A?
- Then for each  $k$ , the frequency of each group of letters in position  $p = k \pmod{L}$  in the ciphertext approximates A.
- Then loop, displacing the ciphertext by  $i$ , and counting the number of matches.
  - Get max when displace by correct key length
  - So just look for the max number of matches!

	displacement
APHUIPLVWGIILTRSQRUBRIZNYQRXWZLBKRHFVN	(0)
NAPHUIPLVWG <b>I</b> ILTRSQRUBRIZNYQRXWZLBKRHFV	(1) 1 match
VNAPHUIPLVWGIILTRSQRUBRIZNYQRXWZLBKRHF	(2) 0 matches
...	
KR <b>H</b> FVNAPHUI <b>I</b> PL <b>V</b> WGIIL <b>T</b> RSQRUB <b>R</b> IZNYQRXWZLB	(6) 5 matches
...	

# Key length: an example

Take any random pair in the ciphertext:

The letter in the top row is shifted by  $i$  (say 0)

The letter in the bottom row is shifted by  $j$  (say 2)

$$\text{Prob(both 'A')} = P('a') * P('y') = 0.082 * 0.020$$

$$\text{Prob(both 'B')} = P('b') * P('z') = 0.015 * 0.001$$

Prob both same (any letter) is \_\_\_\_ or generally \_\_\_\_

Recall, this is maximum when \_\_\_\_\_

When are each letter in the top and bottom rows shifted by same amount?

$$\begin{aligned} A_0 &= (0.082 \quad 0.015 \quad 0.028 \quad \dots & 0.001 \quad 0.020 \quad 0.001) \\ A_2 &= (0.020 \quad 0.001 \quad 0.082 \quad 0.015 \quad 0.028 \quad \dots & 0.023 \quad 0.001) \end{aligned}$$

# Still a bit fuzzy?

- Nothing like implementation to aid understanding!
  - Homework 2: Program it
  - Third week programming quiz: use your program to decrypt a message