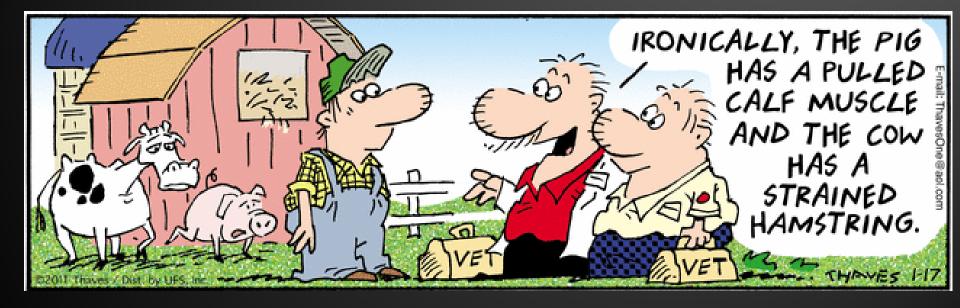


CSSE 230 Day 16

Exhaustive search, backtracking, object-oriented Queens

Check out from SVN: Queens

Questions



Announcements and Agenda

 Exhaustive search, backtracking, and object-oriented queens

Exhaustive Search and Backtracking

A taste of artificial intelligence

Exhaustive search

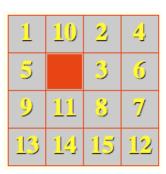
Given: a (large) set of possible solutions to a problem

The "search space"

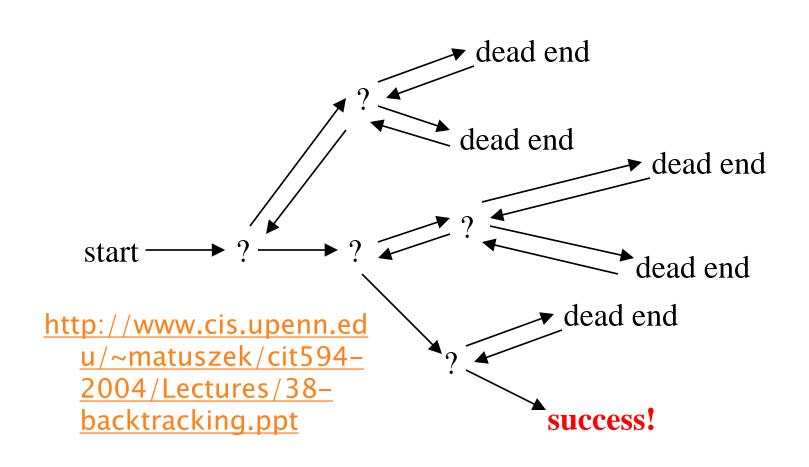
- Goal: Find all solutions (or an optimal solution) from that set
- Questions we ask:
 - How do we represent the possible solutions?
 - How do we organize the search?
 - Can we avoid checking some obvious nonsolutions?

In backtracking, we always try to extend a partial solution

- Examples: Doublets, solving a maze, the "15" puzzle.
- Taken from:
 - http://www.cis.upenn.edu/~matuszek/cit594– 2004/Lectures/38-backtracking.ppt



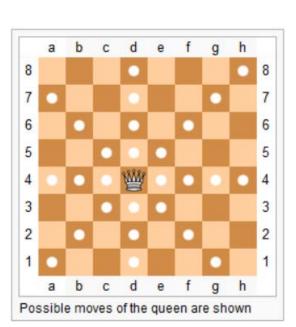
In backtracking, we picture the search as a tree and explore it using a pre-order traversal



The non-attacking chess queens problem is a famous example

- In how many ways can N chess queens be placed on an NxN grid, so that none of the queens can attack any other queen?
- I.e. there are not two queens on the same row, same column, or same diagonal.
- There is no "formula" for generating a solution.
- The famous computer scientist Niklaus Wirth described his approach to the problem in 1971: Program Development by Stepwise Refinement

http://sunnyday.mit.edu/16.355/wirth-refinement.html#3



With a partner, discuss "possible solution" search strategies

- In how many ways can N chess queens be placed on an NxN grid, so that none of the queens can attack any other queen?
 - I.e. no two queens on the same row, same column, or same diagonal.

Two minutes No Peeking!

Search Space Possibilities 1/5

Very naive approach. Perhaps stupid is a better word!

There are N queens, N² squares.

- For each queen, try every possible square, allowing the possibility of multiple queens in the same square.
 - Represent each potential solution as an N-item array of pairs of integers (a row and a column for each queen).
 - Generate all such arrays (you should be able to write code that would do this) and check to see which ones are solutions.
 - Number of possibilities to try in the NxN case:
 - Specific number for N=8:

281,474,976,710,656

Search Space Possibilities 2/5

Slight improvement. There are N queens, N² squares. For each queen, try every possible square, notice that we can't have multiple queens on the same square.

- Represent each potential solution as an N-item array of pairs of integers (a row and a column for each queen).
- Generate all such arrays and check to see which ones are solutions.
- Number of possibilities to try in NxN case:
- Specific number for N=8:

178,462,987,637,760 (vs. 281,474,976,710,656)

Search Space Possibilities 3/5

- Slightly better approach. There are N queens, N columns. If two queens are in the same column, they will attack each other. Thus there must be exactly one queen per column.
- Represent a potential solution as an N-item array of integers.
 - Each array position represents the queen in one column.
 - The number stored in an array position represents the row of that column's queen.
 - Show array for 4x4 solution.
 - Generate all such arrays and check to see which ones are solutions.
 - Number of possibilities to try in NxN case:
 - Specific number for N=8:

Search Space Possibilities 4/5

- Still better approach There must also be exactly one queen per row.
- Represent the data just as before, but notice that the data in the array is a set!
 - Generate each of these and check to see which ones are solutions.
 - How to generate? A good thing to think about.
 - Number of possibilities to try in NxN case:
 - Specific number for N=8:

Search Space Possibilities 5/5

Backtracking solution

- Instead of generating all permutations of N queens and checking to see if each is a solution, we generate "partial placements" by placing one queen at a time on the board
- Once we have successfully placed k<N queens, we try to extend the partial solution by placing a queen in the next column.
- When we extend to N queens, we have a solution.

Experimenting with 8 x 8 Case

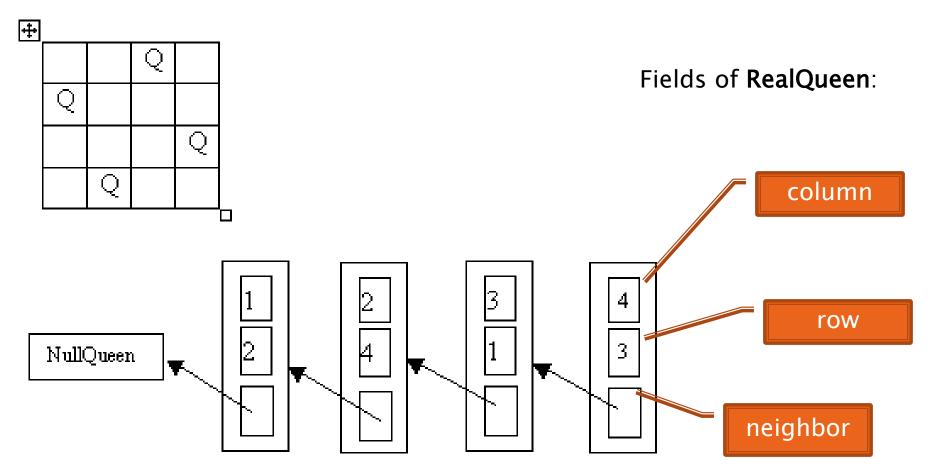
- Play the game:
 - http://homepage.tinet.ie/~pdpals/8queens.htm
- See the solutions:
 - http://www.dcs.ed.ac.uk/home/mlj/demos/queens

Program output:

```
>java RealQueen 5
SOLUTION:
         1 3 5 2 4
         1 4 2 5 3
SOLUTION:
         2 4 1 3 5
SOLUTION:
SOLUTION:
         2 5 3 1 4
SOLUTION:
         3 1 4 2 5
SOLUTION:
         3 5 2 4 1
           4 1 3 5 2
SOLUTION:
           4 2 5 3 1
SOLUTION:
         5 2 4 1 3
SOLUTION:
SOLUTION:
         5 3 1 4 2
```

The Linked List of Queen Objects

Board configuration represented by a linked list of *Queen* objects



Designed by Timothy Budd

http://web.engr.oregonstate.edu/~budd/Books/oopintro3e/info/slides/chap06/java.htm

Outline of the algorithm

- Each queen sends messages directly to its immediate neighbor to the left (and recursively to all of its left neighbors)
- Return value provides information concerning all of the left neighbors:
- Example: neighbor.canAttack(currentRow, col)
 - Message goes to the immediate neighbor, but the real question to be answered by this call is
 - "Hey, neighbors, can any of you attack me if I place myself on this square of the board?"

- findFirst()
- findNext()
- canAttack(int row, int col)

```
Your job (part of WA6):

Understand the job of each of these methods.

Javadoc from the Queen interface can help

Fill in the (recursive) details in the RealQueen class

Debug

More details on next slide
```

More algorithm outline

- Queen asks its neighbors to find the first position in which none of them attack each other
 - Found? Then queen tries to position itself so that it cannot be attacked.
- 2. If the rightmost queen is successful, then a solution has been found! The queens cooperate in recording it.
- 3. Otherwise, the queen asks its neighbors to find the next position in which they do not attack each other
- 4. When the queens get to the point where there is no next non-attacking position, all solutions have been found and the algorithm terminates

Worktime for Queens problem now

I will introduce the term project at the end of class

EditorTrees project

Brief description Meet your team

Reminder: Grading of Team Projects

- I'll assign an overall grade to the project
- Grades of individuals will be adjusted up or down based on team members' assessments
- At the end of the project each of you will:
 - Rate each member of the team, including yourself
 - Write a short Performance Evaluation of each team member with evidence that backs up the rating
 - Positives
 - Key negatives

How to Write a Performance Evaluation

- Document dates and actions:
 - Jan. 1, stayed after mtg. to help Bob with hashing
 - Jan. 19, failed to complete UML diagram as agreed
 - [this means take notes so you don't forget]
- List positives:
 - The only way to help people improve!
- List key negatives:
 - Not all negatives
 - Egos are too fragile for long lists, can't fix everything at once anyway

EditorTrees project

- In general, *implementation* of a Data Structure is separate from *application*.
- Most CSSE 230 projects have used existing data structures to create an application
- In this project you will create an efficient data structure that *could* be used in a text editor.
- But you will not create an application that uses it.
- EditTree: A height-balanced (but not AVL) binary tree with rank. Insertion and deletion are by position, not by natural ordering of the inserted elements.
 - So is it a BST?
- Lots of O(log N) operations
- Uses balance code and rank.

EditorTrees Teams

See schedule page

Meet your partners to plan when you will meet to begin work.

Suggestion: Meet before next class to discuss the project requirements. Formulate a list questions to ask during next class.

Whether or not you meet before next class, read the EditorTrees requirements and come with questions.

Next class will be a workday.