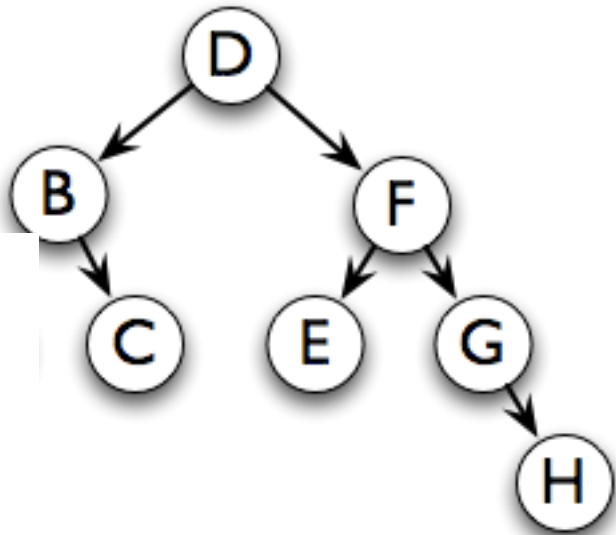


CSSE 230 Day 13

Height-Balanced Trees



Announcements

- ▶ Doublets Milestone 1 due next Tuesday night
- ▶ Exams redux now and Tuesday

Today's Agenda (a lot of it may spill over into Monday)

- ▶ Exam 1 review?
- ▶ Doublets: what's it all about?
- ▶ Finding k -th smallest in BST
- ▶ Meet your Doublets partner
- ▶ Another induction example
- ▶ Recap: The need for balanced trees
- ▶ Analysis of worst case for height-balanced (AVL) trees

Doublets: What's it all about?

Welcome to Doublets, a game of "verbal torture."

Enter starting word: *flour*

Enter ending word: *bread*

Enter chain manager (s: stack, q: queue, x: exit): *s*

Chain: [flour, floor, flood, blood, bloom, gloom, groom, broom, brood, broad, bread]

Length: 11

Candidates: 16

Max size: 6

Enter starting word: *wet*

Enter ending word: *dry*

Enter chain manager (s: stack, q: queue, x: exit): *q*

Chain: [wet, set, sat, say, day, dry]

Length: 6

Candidates: 82651

Max size: 847047

Enter starting word: *who*

Enter ending word: *rye*

The word "oat" is not valid. Please try again.

Enter starting word: *owner*

Enter ending word: *bribe*

Enter chain manager (s: stack, q: queue, x: exit): *s*

No doublet chain exists from owner to bribe.

Enter starting word: *C*

Enter chain manager (s: stack, q: queue, x: exit): *x*

Goodbye!

StackChainManager: depth-first search

QueueChainManager: breadth-first search

PriorityQueueChainManager: First extend the chain that ends with a word that is closest to the ending word.

Answers will vary from these!

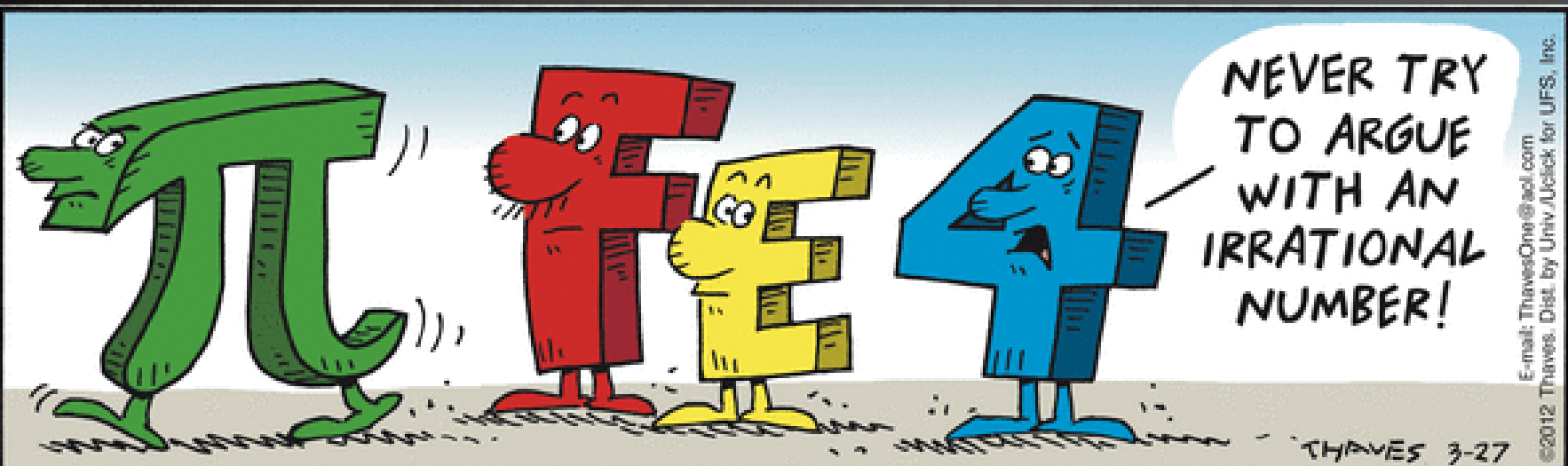
A **Link** is the collection of all words that can be reached from a given word in one step. I.e. all words that can be made from the given word by substituting a single letter.

A **Chain** is a sequence of words (no duplicates) such that each word can be made from the one before it by a single letter substitution.

A **ChainManager** stores a collection of chains, and tries to extend one at a time, with a goal of extending to the ending word.

BST with Rank

- » Explore the concept
- How do Find and Insert work?



BSTs are an efficient way to represent ordered lists

- ▶ What's the performance of
 - insertion? $O(h(T))$
 - deletion? $O(h(T))$
 - find? $O(h(T))$
 - iteration? $O(n)$ to iterate through all
- ▶ What about finding the k^{th} smallest element?

We can find the k th smallest element easily if we add a *rank* field to BinaryNode

- ▶ Gives the in-order position of this node within its own subtree
 - i.e., the size of its left subtree
- ▶ How would we do *findK_{th}*?
- ▶ *Insert* and *delete* start similarly



0-based
indexing

Another induction example (we'll use this result)

- ▶ Recall our definition of the Fibonacci numbers:
 - $F_0 = 0, F_1 = 1, F_{n+2} = F_{n+1} + F_n$
- ▶ An exercise from the textbook

7.8 Prove by induction the formula

$$F_N = \frac{1}{\sqrt{5}} \left(\left(\frac{1 + \sqrt{5}}{2} \right)^N - \left(\frac{1 - \sqrt{5}}{2} \right)^N \right)$$

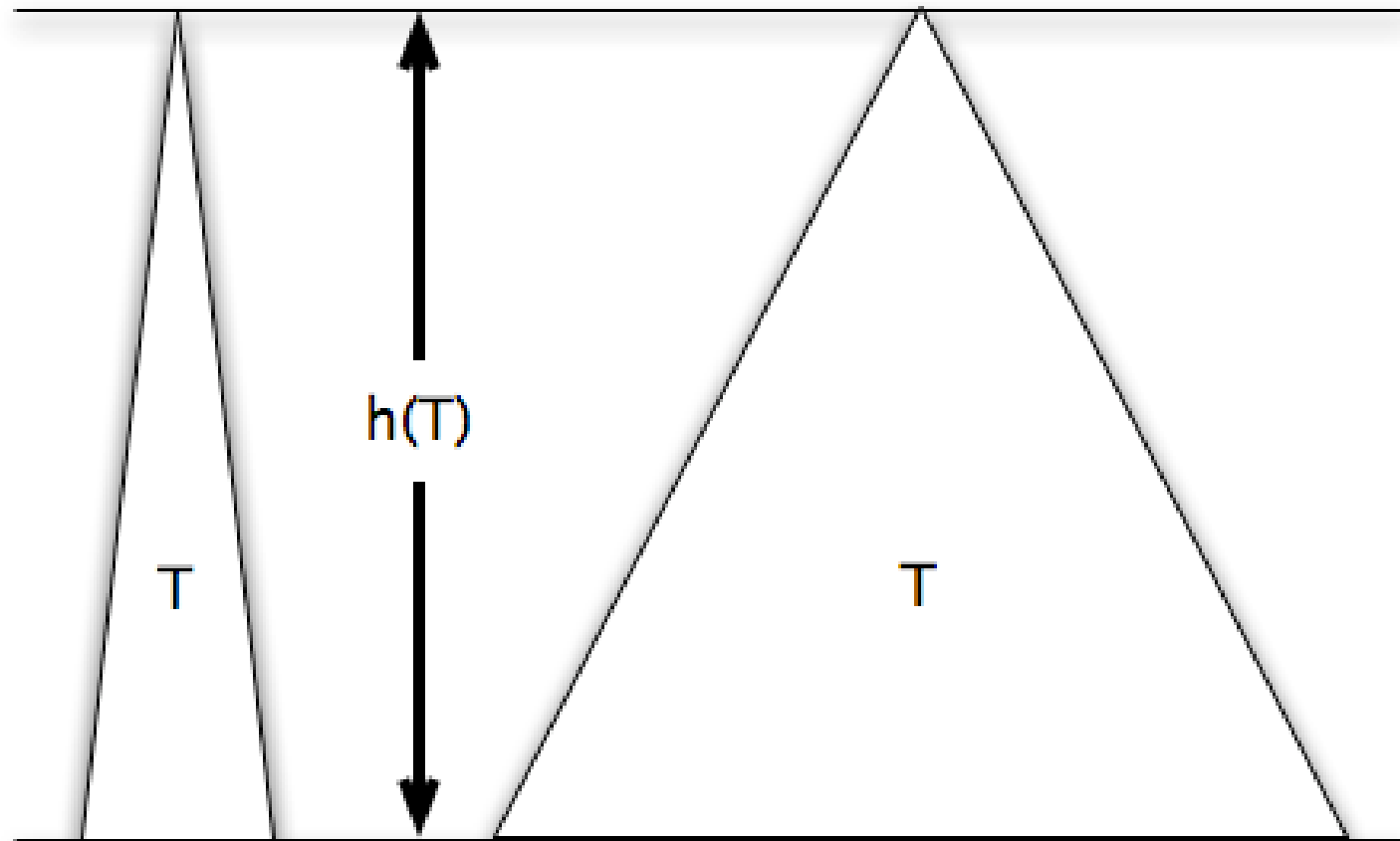
Recall: How to show that property $P(n)$ is true for all $n \geq n_0$:

- (1) Show the base case(s) directly
- (2) Show that if $P(j)$ is true for all j with $n_0 \leq j < k$, then $P(k)$ is true also

Details of step 2:

- a. Write down the induction assumption for this specific problem
- b. Write down what you need to show
- c. Show it, using the induction assumption

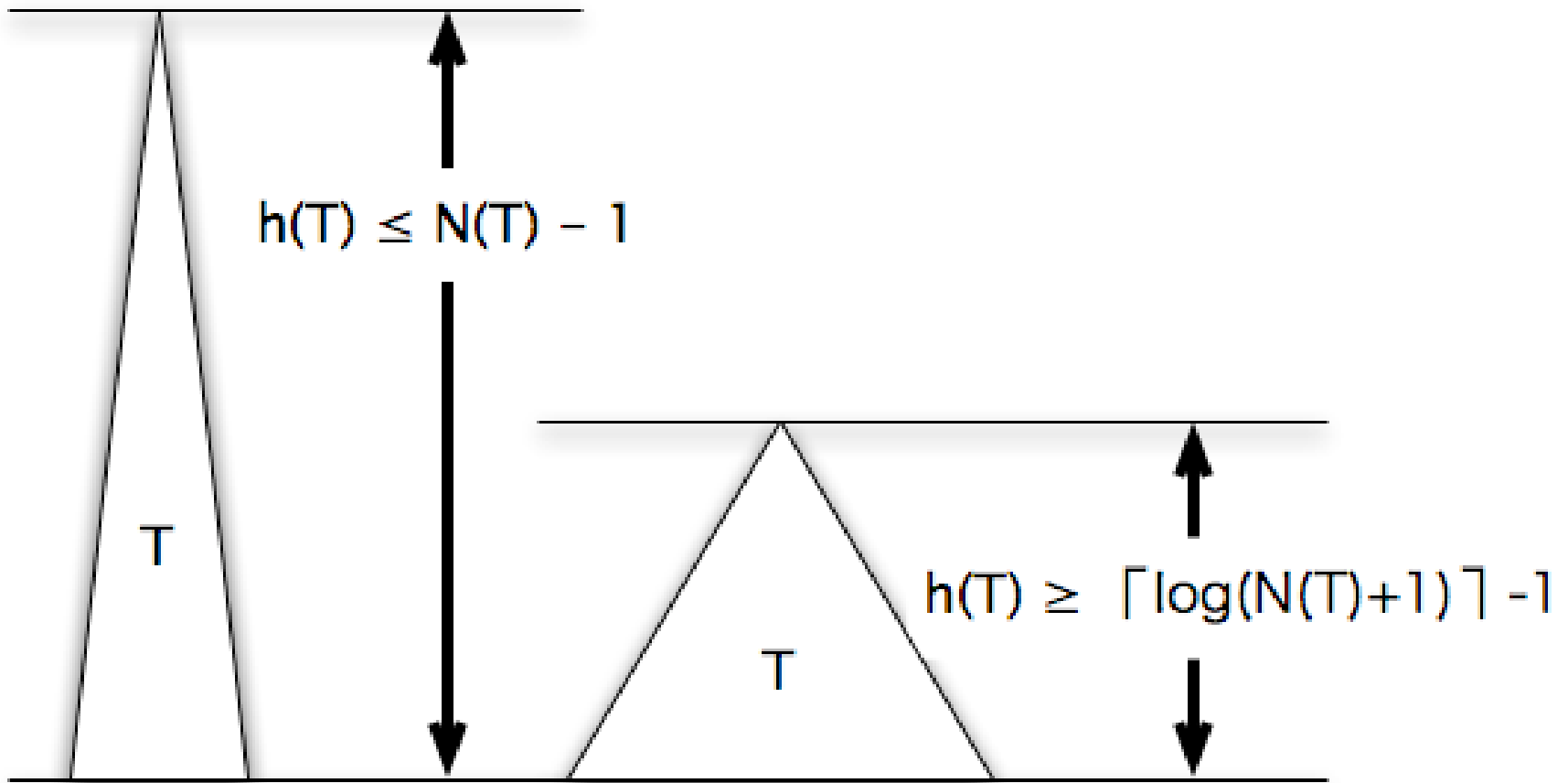
Review: The number of nodes in a tree with height $h(T)$ is bounded



$$N(T) \geq h(T) + 1$$

$$N(T) \leq 2^{h(T)+1} - 1$$

Review: Therefore the height of a tree with $N(T)$ nodes is also bounded

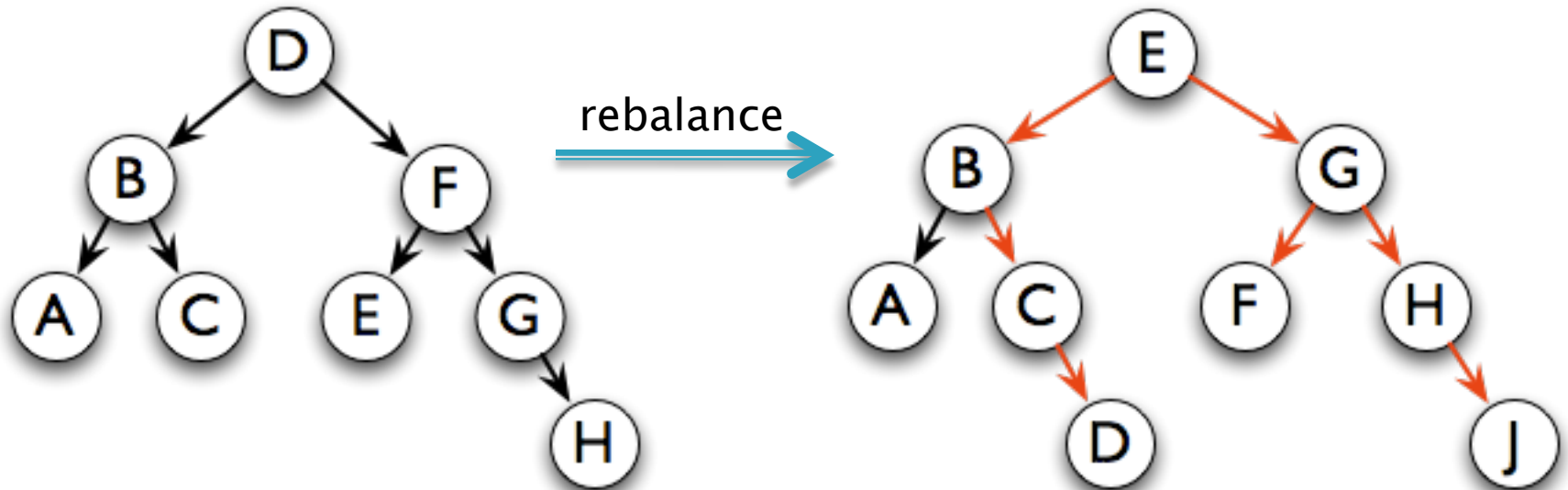


We want to keep trees balanced so that the run time of BST algorithms is minimized

- ▶ BST algorithms are $O(h(T))$
- ▶ Minimum value of $h(T)$ is $\lceil \log(N(T)+1) \rceil - 1$
- ▶ Can we rearrange the tree after an insertion to guarantee that $h(T)$ is always minimized?

But keeping complete balance is too expensive!

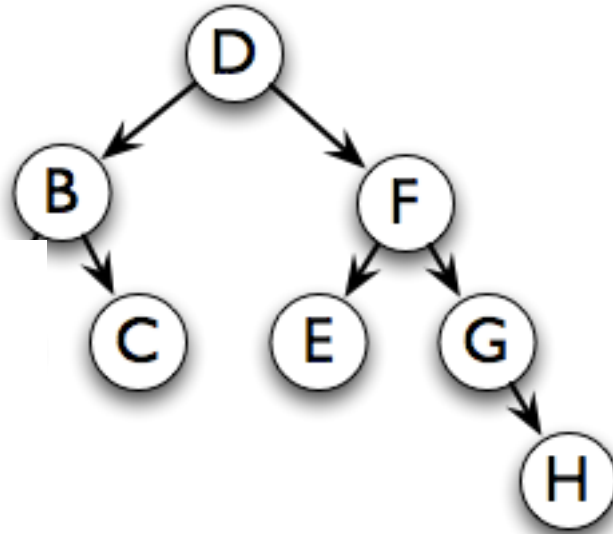
- ▶ Height of the tree can vary from $\log N$ to N
- ▶ Where would J go in this tree?
- ▶ What if we keep the tree perfectly balanced?
 - so height is always proportional to $\log N$
- ▶ What does it take to balance that tree?
- ▶ Keeping completely balanced is too expensive:
 - $O(N)$ to rebalance after insertion or deletion



Solution: Height Balanced Trees (less is more)

Height-Balanced Trees have subtrees whose heights differ by at most 1

Still height-balanced?



More precisely , a binary tree T is height balanced if

T is empty, or if

$| \text{height}(T_L) - \text{height}(T_R) | \leq 1$, and

T_L and T_R are both height balanced.

What is the tallest height-balanced tree with N nodes?

Is it taller than a completely balanced tree?

- Consider the dual concept: find the minimum number of nodes for height h .

A binary search tree T is height balanced if

T is empty, or if

$| \text{height}(T_L) - \text{height}(T_R) | \leq 1$, and T_L and T_R are both height balanced.

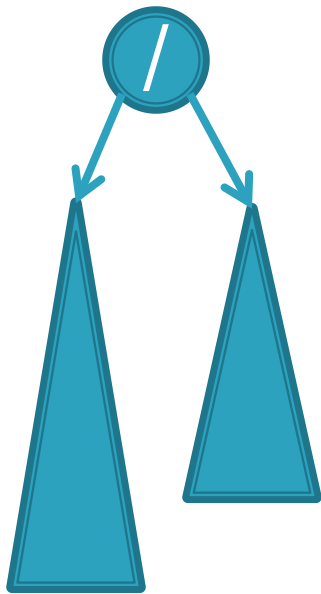
An AVL tree is a height-balanced BST that maintains balance using “rotations”

- ▶ Named for authors of original paper, **A**delson-**V**elskii and **L**andis (1962).
- ▶ Max. height of an AVL tree with N nodes is:
 $H < 1.44 \log(N+2) - 1.328 = O(\log N)$

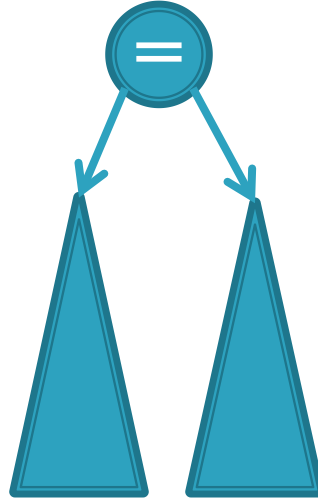
Our goal is to rebalance an AVL tree after insert/delete in $O(\log n)$ time

- ▶ Why?
- ▶ Worst cases for BST operations are $O(h(T))$
 - **find**, **insert**, and **delete**
- ▶ $h(T)$ can vary from $O(\log N)$ to $O(N)$
- ▶ Height of a height-balanced tree is $O(\log N)$
- ▶ So if we can rebalance after insert or delete in $O(\log N)$, then **all operations are $O(\log N)$**

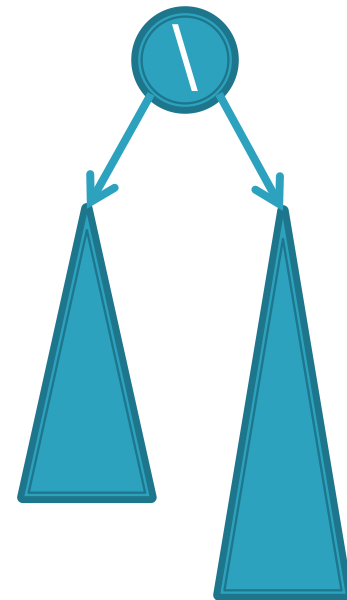
AVL nodes are just like BinaryNodes,
but also have an extra “balance code”



or



or

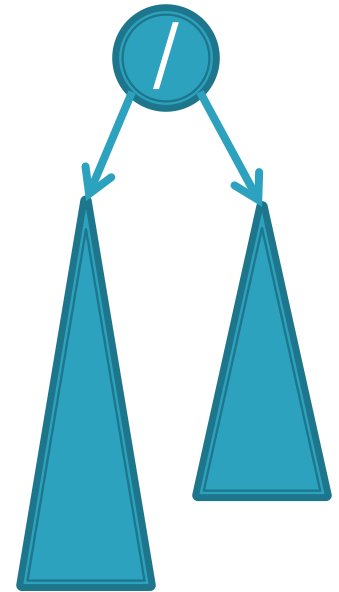


Different representations for / = \ :

- Just two bits in a low-level language
- Enum in a higher-level language

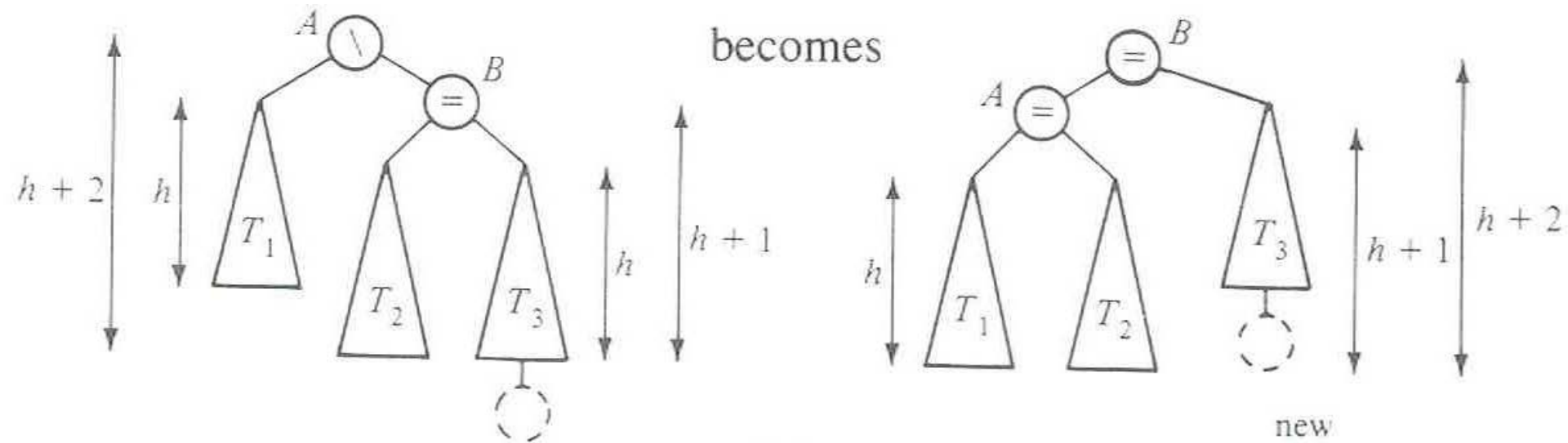
AVL Tree (Re)balancing Act

- ▶ Assume tree is height-balanced before insertion
- ▶ Insert as usual for a BST
- ▶ Move up from the newly inserted node to the lowest “unbalanced” node (if any)
 - Use the **balance code** to detect unbalance – how?
- ▶ Do appropriate rotation to balance the sub-tree rooted at this unbalanced node



Four types of rotations are required to remove different cases of tree imbalances

- ▶ For example, a *single left rotation*:



We'll pick up here next class...