

CSSE 230 Day 24

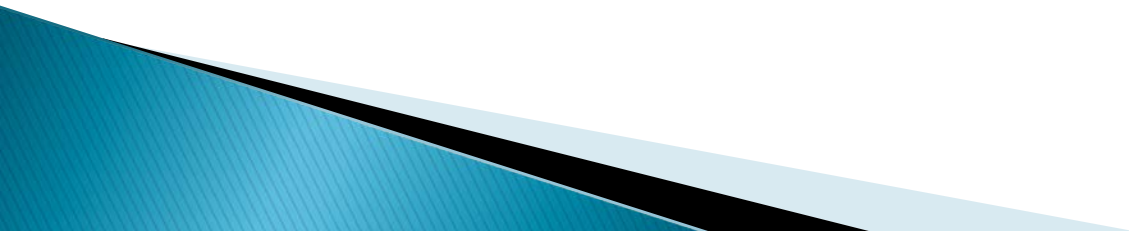
Sorting Lower Bound
Radix Sort

A Lower-Bound on Sorting Time

- » We can't do much better than what we already know how to do.

What's the best best case?

- ▶ Lower bound for best case?
- ▶ A particular algorithm that achieves this?



What's the best worst case?

- ▶ Want a function $f(N)$ such that the worst case running time for **all sorting algorithms** is $\Omega(f(N))$
- ▶ How do we get a handle on “all sorting algorithms”?

Tricky!

What are “all sorting algorithms”?

- ▶ We can't list all sorting algorithms and analyze all of them
 - Why not?
- ▶ But we can find a **uniform representation** of any sorting algorithm that is based on **comparing** elements of the array to each other

This "uniform representation" idea is exploited in a big way in Theory of Computation, e.g., to demonstrate the unsolvability of the "Halting Problem"

First of all...

- ▶ The problem of sorting N elements is at least as hard as determining their ordering
 - e.g., determining that $a_3 < a_4 < a_1 < a_5 < a_2$
- ▶ So any lower bound on all "order-determination" algorithms is also a lower bound on "all sorting algorithms"

Sort Decision Trees

- ▶ Let A be any **comparison-based algorithm** for sorting an array of distinct elements
- ▶ Note: sorting is asymptotically equivalent to determining the correct order of the originals
- ▶ We can draw an EBT that corresponds to the comparisons that will be used by A to sort an array of N elements
 - This is called a **sort decision tree**
 - Just a pen-and-paper concept, not actually a data structure
 - Different algorithms will have different trees

So what?

- ▶ Minimum number of external nodes in a sort decision tree? (As a function of N)
- ▶ Is this number dependent on the algorithm?
- ▶ What's the height of the shortest EBT with that many external nodes?

$$\lceil \log N! \rceil \approx N \log N - 1.44N = \Omega(N \log N)$$

No comparison-based sorting algorithm, known or not yet discovered, can ever do better than this!

Can we do better than $N \log N$?

- ▶ $\Omega(N \log N)$ is the best we can do if we compare items
- ▶ Can we sort without comparing items?

Yes, we can! We can sort if we avoid comparing items

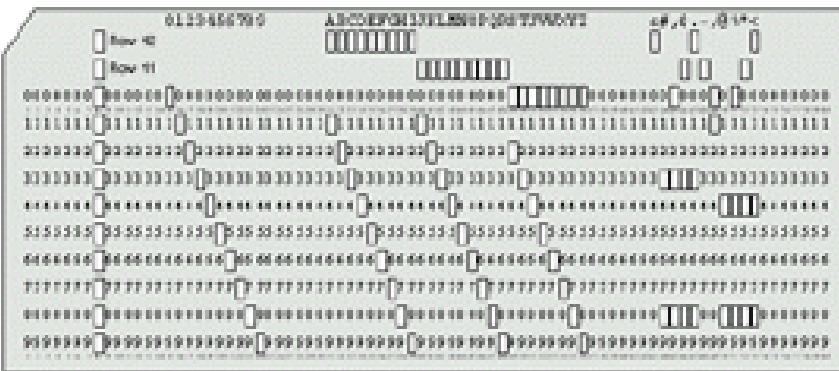
Q5

- ▶ $O(N)$ sort: Bucket sort
 - Works if possible values come from limited range
 - Example: Exam grades histogram
- ▶ A variation: Radix sort

Radix sort

- ▶ A picture is worth 10^3 words, but an animation is worth 2^{10} pictures, so we will look at one.
- ▶ <http://www.cs.auckland.ac.nz/software/AlgAnim/radixsort.html>

Radix sort example: card sorter



Used an appropriate
combo of
mechanical, digital,
and human effort to
get the job done.

Type 82 Electric Punched Card Sorting Machine