CSSE 230 Day 14

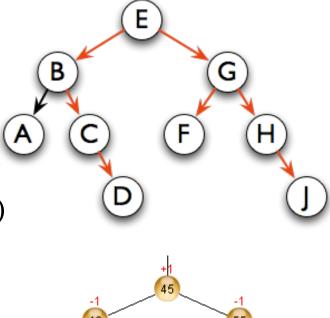
AVL trees and rotations

Pay careful attention to the ACM Code of Ethics essay

- Part of written assignment 5
 - Examine the Code of Ethics of the ACM
 - Focus on property rights
 - Write a short reaction (up to 1 page single-spaced)
 - Details are in the assignment

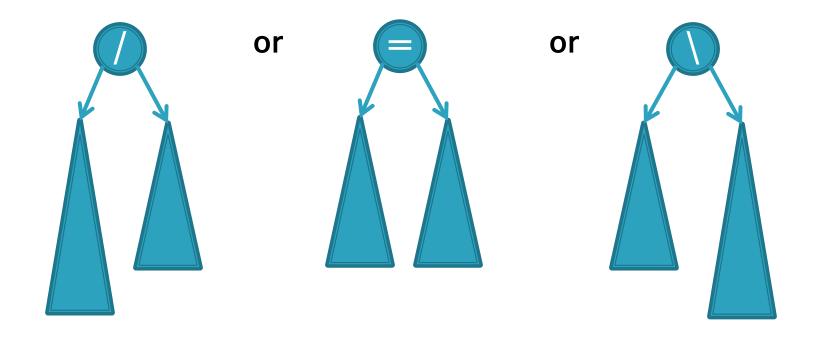
Summary: for tree operations to be fast, we must keep the tree somewhat balanced in O(log n) time

- Operations (insert, delete, search) are O(height)
- Tree height is O(log n) if perfectly balanced
 - But maintaining perfect balance is O(n)
- Height-balanced trees are still O(log n)
 - For T with height h, $N(T) \le Fib(h+3) 1$
 - \circ So H < 1.44 log (N+2) 1.328 *
- AVL (Adelson-Velskii and Landis) trees maintain height-balance using rotations



<sup>Are rotations O(log n)? We'll see...
* Don't use this formula to do problem 4 on WA5. Instead draw the trees like we did on quiz #5 yesterday.</sup>

AVL nodes are just like BinaryNodes, but also have an extra "balance code"

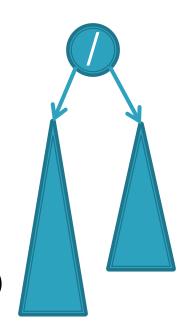


Different representations for $/ = \setminus$:

- Just two bits in a low-level language
- Enum in a higher-level language

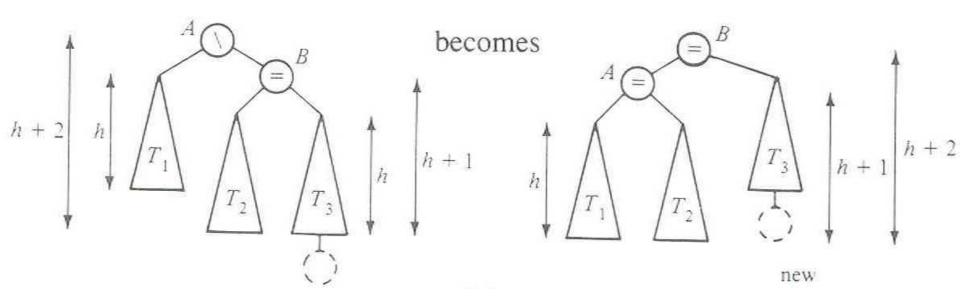
AVL Tree (Re)balancing Act

- Assume tree is height-balanced before insertion
- Insert as usual for a BST
- Move up from the newly inserted node to the lowest "unbalanced" node (if any)
 - Use the balance code to detect unbalance how?
- Do an appropriate rotation to balance the sub-tree rooted at this unbalanced node



Four types of rotations are required to remove different cases of tree imbalances

For example, a *single left rotation*:

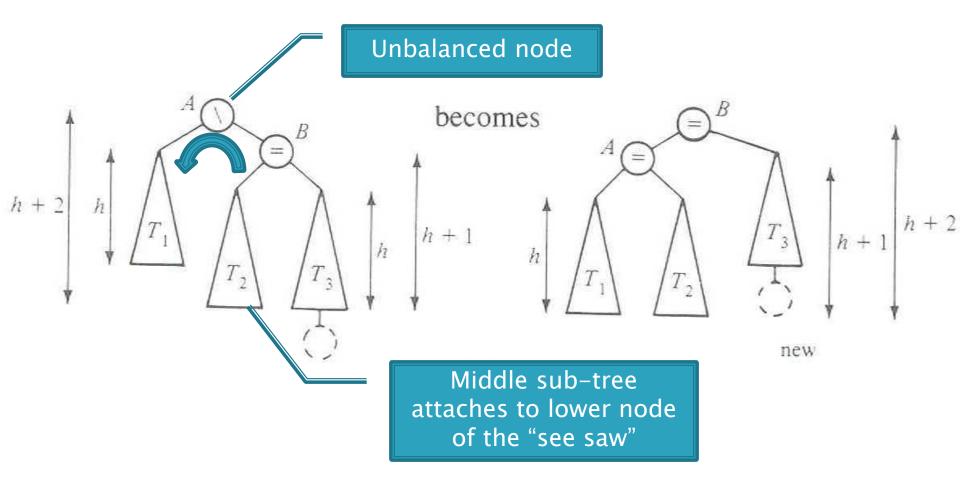


We rotate by pulling the "too tall" sub-tree up and pushing the "too short" sub-tree down

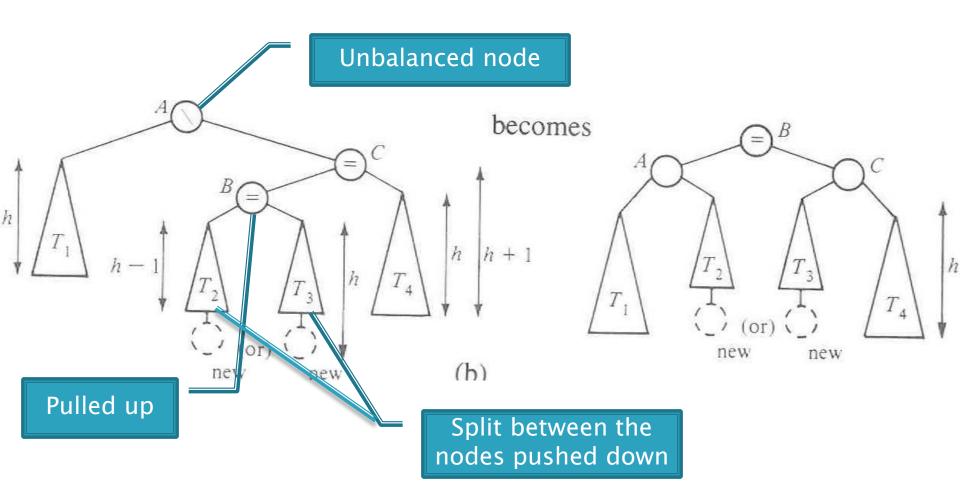
Two basic cases

- "See saw" case:
 - Too-tall sub-tree is on the outside
 - So tip the see saw so it's level
- "Suck in your gut" case:
 - Too-tall sub-tree is in the middle
 - Pull its root up a level

Single Left Rotation

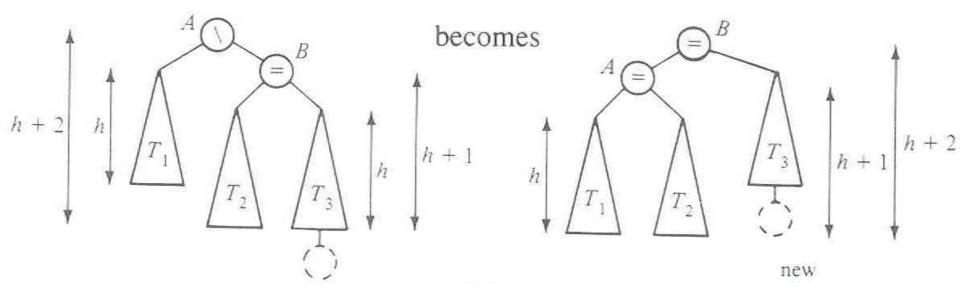


Diagrams are from Data Structures by E.M. Reingold and W.J. Hansen



Weiss calls this "right-left double rotation"

Your turn — work with a partner



- Write the method:
- Returns a reference to the new root of this subtree.
- Don't forget to set the balanceCode fields of the nodes.

More practice—(sometime after class)

- Write the method:

- Returns a reference to the new root of this subtree.
- Rotation is mirror image of double rotation from an earlier slide

- Both kinds of rotation leave height the same as before the insertion!
- Is insertion plus rotation cost really O(log N)?

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Insertion/deletion
in AVL Tree:

Find the imbalance point (if any):

Single or double rotation:

in deletion case, may have
to do O(log N) rotations

Total work:

O(log n)

O(1)

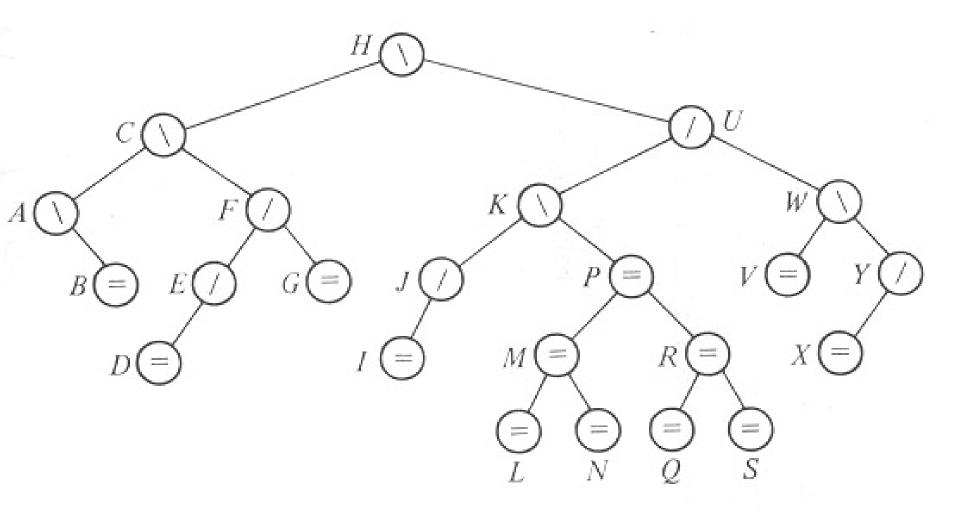
O(log n)
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Which kind of rotation to do after an insertion?

Depends on the first two links in the path from the lowest node that has the imbalance (A) down to the newly-inserted node.

First link (down from A)	Second link (down from A's child)	Rotation type (rotate "around A's position")
Left	Left	Single right
Left	Right	Double right
Right	Right	Single left
Right	Left	Double left

A sample AVL tree



Insert **HA** into the tree, then **DA**, then **O**. Delete **G** from the original tree, then **I**, **J**, **V**.

Your turn again

- Start with an empty AVL tree.
- Add elements in the following order; do the appropriate rotations when needed.
 - 1 2 3 4 5 6 11 13 12 10 9 8 7
- How should we rebalance if each of the following sequences is deleted from the above tree?
 - (10 9 7 8) (13) (1 5)
 - For each of the three sequences, start with the original 13-element tree. E.g. when deleting 13, assume 10 9 8 7 are still in the tree.

Work with your Doublets partner. When you finish, work on Doublets or WA5. Or write the rotateDoubleRight code from a previous slide

